EECS150 - Digital Design

Lecture 23 - Arithmetic Blocks, Part 2 + Shifters

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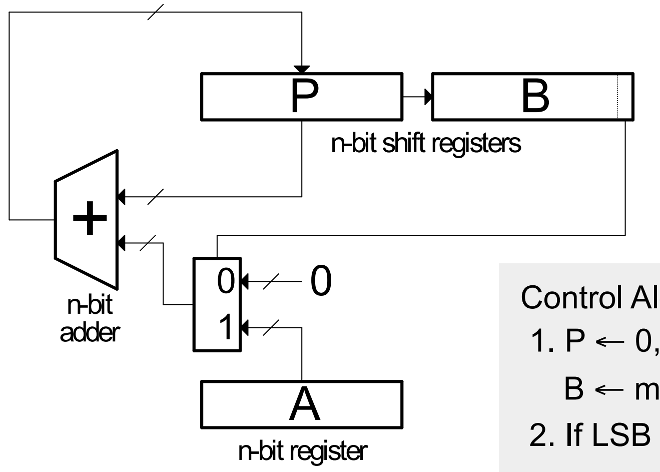
Multiplication

$$\begin{array}{ccccccc l}
 & & & a_3 & a_2 & a_1 & a_0 & \leftarrow \textit{Multiplicand}\\
 & & & b_3 & b_2 & b_1 & b_0 & \leftarrow \textit{Multiplier}\\ \cline{4-7}
 & & X & a_3b_0 & a_2b_0 & a_1b_0 & a_0b_0 & \\
 & & a_3b_1 & a_2b_1 & a_1b_1 & a_0b_1 & & \Bigg\} \textit{Partial products}\\
 & a_3b_2 & a_2b_2 & a_1b_2 & a_0b_2 & & & \\
 a_3b_3 & a_2b_3 & a_1b_3 & a_0b_3 & & & & \\ \hline
 & & & \ldots & & a_1b_0+a_0b_1 & a_0b_0 & \leftarrow \textit{Product}
\end{array}$$

Many different circuits exist for multiplication.
Each one has a different balance between
speed (performance) and amount of logic (cost).

“Shift and Add” Multiplier



- Sums each partial product, one at a time.
- In binary, each partial product is shifted versions of A or 0.

Control Algorithm:

1. P ← 0, A ← multiplicand,
   B ← multiplier
2. If LSB of B==1 then add A to P
   else add 0
3. Shift [P][B] right 1
4. Repeat steps 2 and 3 n-1 times.
5. [P][B] has product.

- Cost α n, T = n clock cycles.
- What is the critical path for determining the min clock period?

“Shift and Add” Multiplier

Signed Multiplication:

Remember for 2’s complement numbers MSB has negative weight:

$$X = \sum_{i=0}^{N-2} x_i 2^i - x_{n-1} 2^{n-1}$$

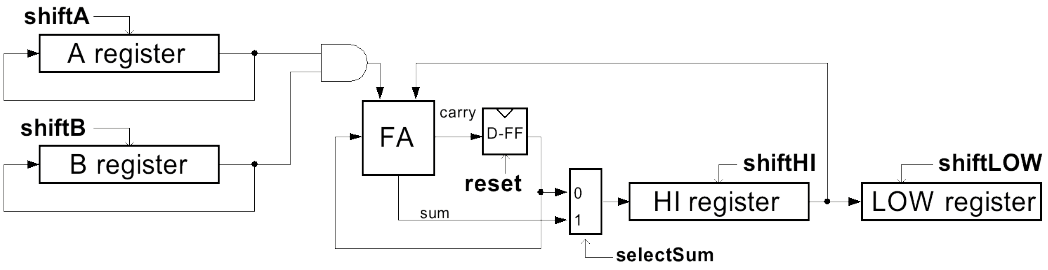
ex: -6 = 11010_2 = $0\bullet2^0 + 1\bullet2^1 + 0\bullet2^2 + 1\bullet2^3 - 1\bullet2^4$

= 0 + 2 + 0 + 8 - 16 = -6

- Therefore for multiplication:
  - a) subtract final partial product
  - b) sign-extend partial products
- Modifications to shift & add circuit:
  - a) adder/subtractor
  - b) sign-extender on P shifter register

Bit-serial Multiplier

- Bit-serial multiplier (n^2 cycles, one bit of result per n cycles):



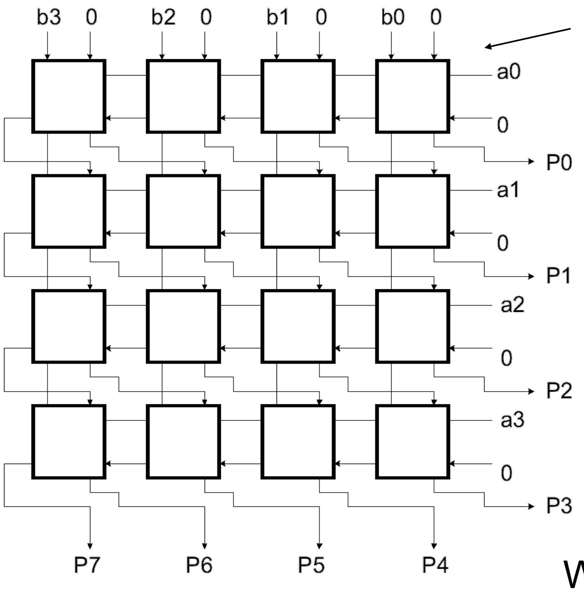
- Control Algorithm:

```
repeat n cycles {   // outer (i) loop
        repeat n cycles{    // inner (j) loop
                shiftA, selectSum, shiftHI
        }
        shiftB, shiftHI, shiftLOW, reset
}
```

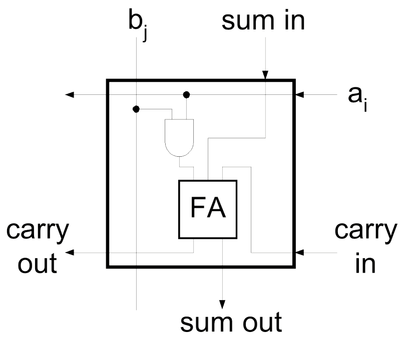
Note: The occurrence of a control signal x means x=1. The absence of x means x=0.

Array Multiplier

Single cycle multiply: Generates all n partial products simultaneously.



Each row: n-bit adder with AND gates



What is the critical path?

Carry-Save Addition

- Speeding up multiplication is a matter of speeding up the summing of the partial products.
- “Carry-save” addition can help.
- Carry-save addition passes (saves) the carries to the output, rather than propagating them.

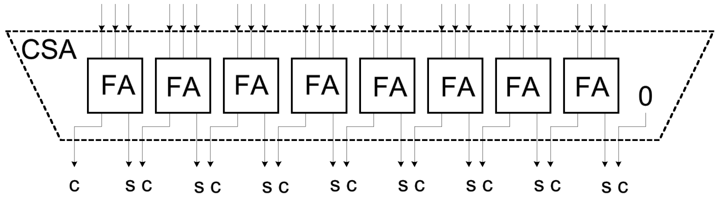
- Example: sum three numbers, 3_{10} = 0011, 2_{10} = 0010, 3_{10} = 0011

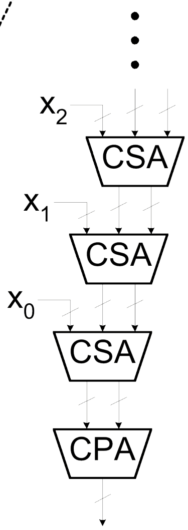
3_{10} 0011
+ 2_{10} 0010
c 0100 = 4_{10}
s 0001 = 1_{10}
carry-save add

carry-save add
3_{10} 0011
c 0010 = 2_{10}
s 0110 = 6_{10}
1000 = 8_{10}
carry-propagate add

- In general, *carry-save* addition takes in 3 numbers and produces 2.
- Whereas, *carry-propagate* takes 2 and produces 1.
- With this technique, we can avoid carry propagation until final addition

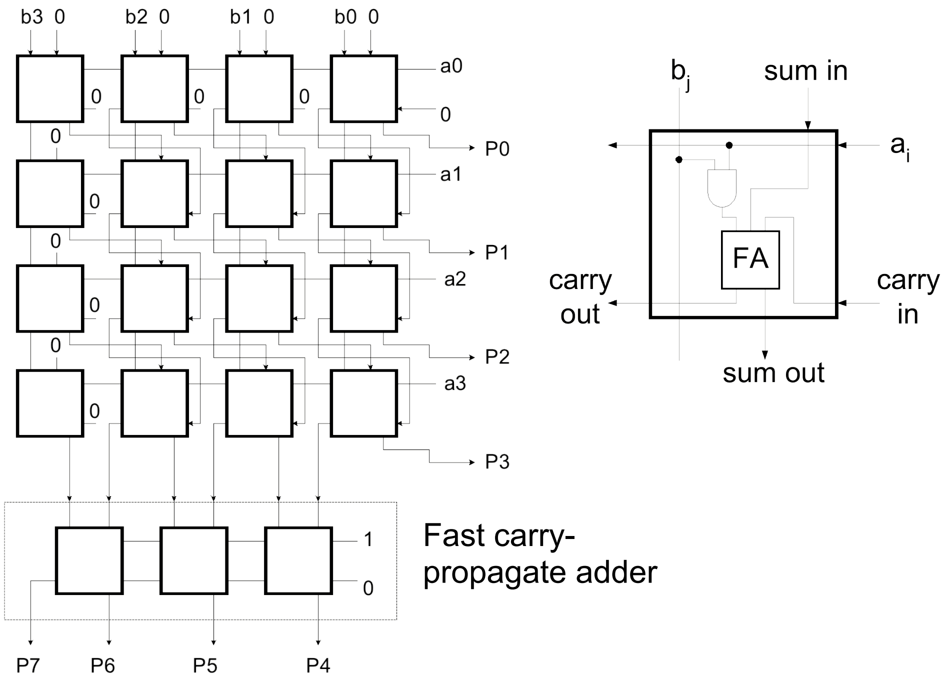
Carry-save Circuits





- When adding sets of numbers, carry-save can be used on all but the final sum.
- Standard adder (carry propagate) is used for final sum.
- Carry-save is fast (no carry propagation) and cheap (same cost as ripple adder)

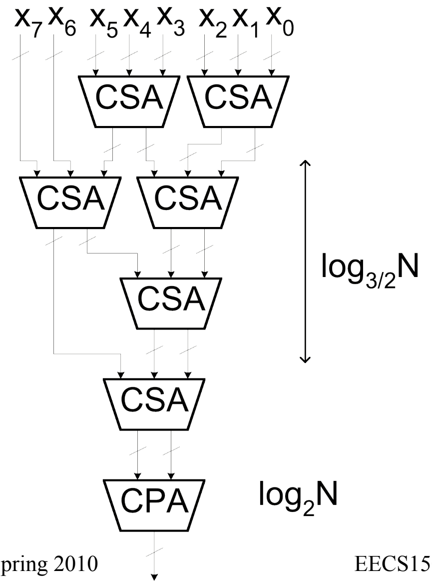
Array Multiplier using Carry-save Addition



Carry-save Addition

CSA is associative and communitive. For example:

$$(((X_0 + X_1) + X_2) + X_3) = ((X_0 + X_1) + (X_2 + X_3))$$



- A balanced tree can be used to reduce the logic delay.
- This structure is the basis of the ***Wallace Tree Multiplier***.
- Partial products are summed with the CSA tree. Fast CPA (ex: CLA) is used for final sum.
- Multiplier delay $\alpha \log_{3/2}N + \log_2 N$

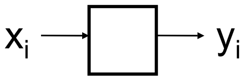
Constant Multiplication

- *Our discussion so far has assumed both the multiplicand (A) and the multiplier (B) can vary at runtime.*
- What if one of the two is a constant?

$$Y = C * X$$

- "Constant Coefficient" multiplication comes up often in signal processing and other hardware. Ex:

$$y_i = \alpha y_{i-1} + x_i$$

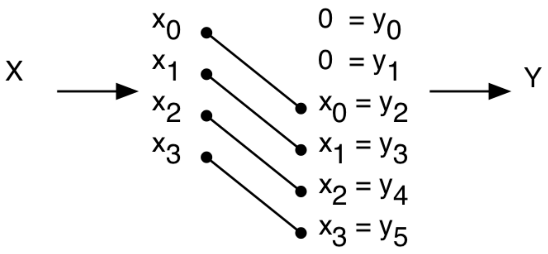


where α is an application dependent constant that is hard-wired into the circuit.

- How do we build and array style (combinational) multiplier that takes advantage of the constancy of one of the operands?

Multiplication by a Constant

- If the constant C in C*X is a power of 2, then the multiplication is simply a shift of X.
- Ex: 4*X

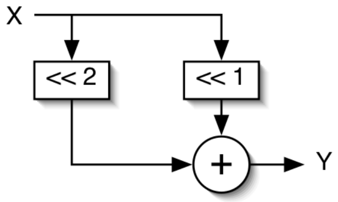


- What about division?

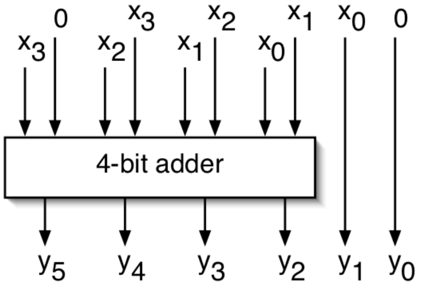
- What about multiplication by non- powers of 2?

Multiplication by a Constant

- In general, a combination of fixed shifts and addition:
  - Ex: $6 * X = 0110 * X = (2^2 + 2^1) * X$

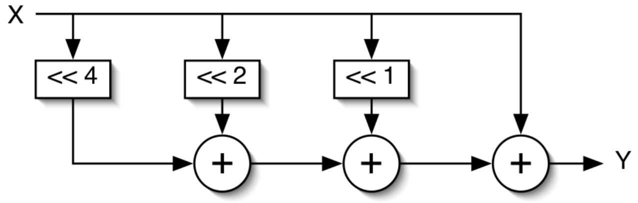


  - Details:



Multiplication by a Constant

- Another example: $C = 23_{10} = 010111$



- *In general, the number of additions equals the number of 1's in the constant minus one.*
- Using carry-save adders (for all but one of these) helps reduce the delay and cost, but the number of adders is still the number of 1's in C minus 2.

- Is there a way to further reduce the number of adders (and thus the cost and delay)?

Multiplication using Subtraction

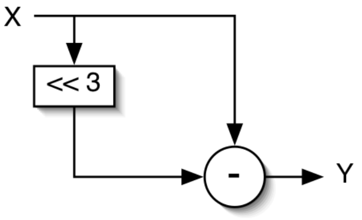
- *Subtraction is ~ the same cost and delay as addition.*
- Consider C*X where C is the constant value $15_{10} = 01111$.

  C*X requires 3 additions.
- We can "recode" 15

  from $01111 = (2^3 + 2^2 + 2^1 + 2^0)$

  to $1000\bar{1} = (2^4 - 2^0)$

  where $\bar{1}$ means negative weight.
- Therefore, 15*X can be implemented with only one subtractor.



Canonic Signed Digit Representation

- CSD represents numbers using 1, $\bar{1}$, & 0 with the least possible number of non-zero digits.
  - Strings of 2 or more non-zero digits are replaced.
  - Leads to a unique representation.
- To form CSD representation might take 2 passes:
  - First pass: replace all occurrences of 2 or more 1's:

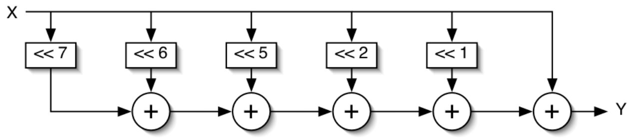
    $01..10$ by $10..\bar{1}0$
  - Second pass: same as a above, plus replace $0\bar{1}10$ by $00\bar{1}0$
- Examples:

011101 = 29	0010111 = 23	0110110 = 54
$100\bar{1}01$ = 32 - 4 + 1	$001100\bar{1}$	$10\bar{1}10\bar{1}0$
	$010\bar{1}00\bar{1}$ = 32 - 8 - 1	$100\bar{1}0\bar{1}0$ = 64 - 8 - 2

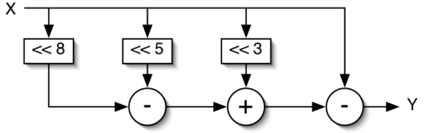
- Can we further simplify the multiplier circuits?

“Constant Coefficient Multiplication” (KCM)

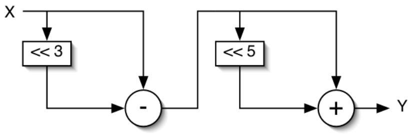
Binary multiplier: $Y = 231*X = (2^7 + 2^6 + 2^5 + 2^2 + 2^1+2^0)*X$



- CSD helps, but the multipliers are limited to shifts followed by adds.
  - CSD multiplier: $Y = 231*X = (2^8 - 2^5 + 2^3 - 2^0)*X$



- How about shift/add/shift/add …?
  - KCM multiplier: $Y = 231*X = 7*33*X = (2^3 - 2^0)*(2^5 + 2^0)*X$

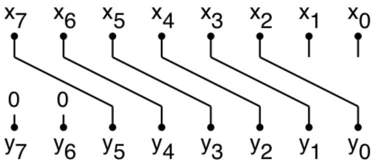


- No simple algorithm exists to determine the optimal KCM representation.
- Most use exhaustive search method.

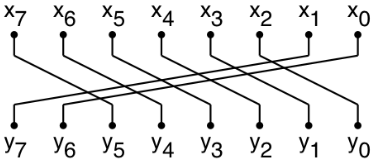
Fixed Shifters / Rotators

- “fixed” shifters “hardwire” the shift amount into the circuit.
- Ex: verilog: X >> 2
  - (right shift X by 2 places)
- Fixed shift/rotator is nothing but wires!

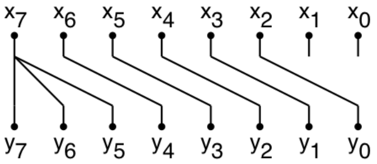
  So what?



Logical Shift



Rotate

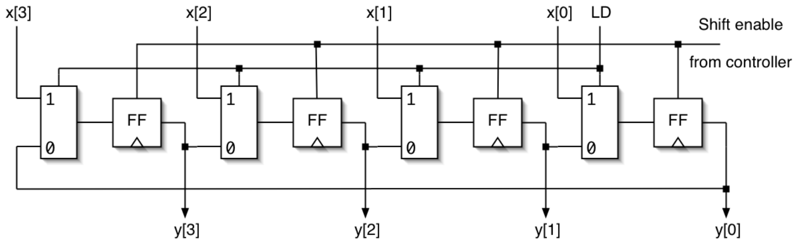


Arithmetic Shift

Variable Shifters / Rotators

- Example: X >> S, where S is unknown when we synthesize the circuit.
- Uses: shift instruction in processors (ARM includes a shift on every instruction), floating-point arithmetic, division/multiplication by powers of 2, etc.
- One way to build this is a simple shift-register:

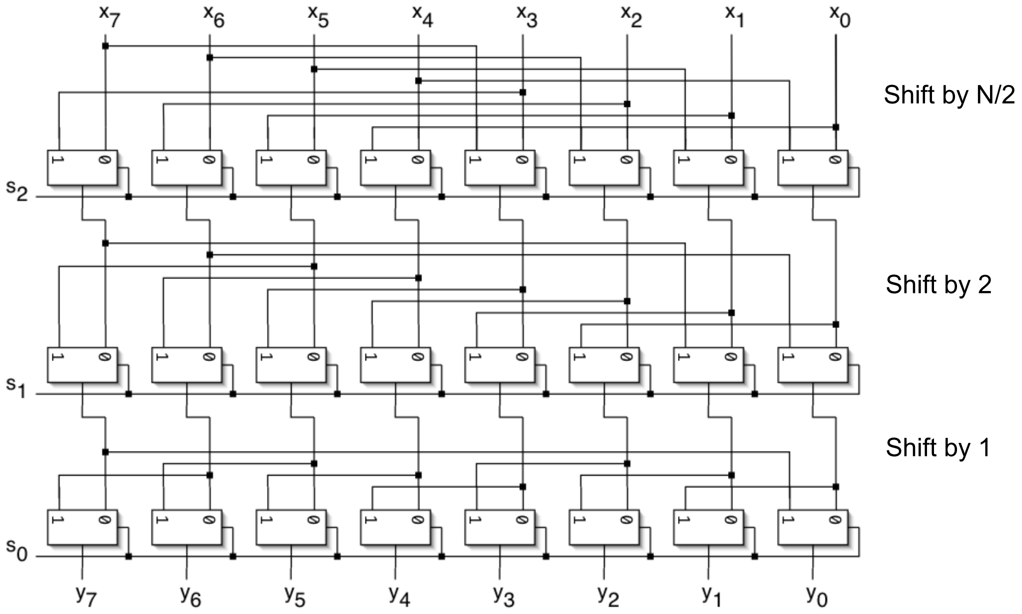
  a) Load word, b) shift enable for S cycles, c) read word.



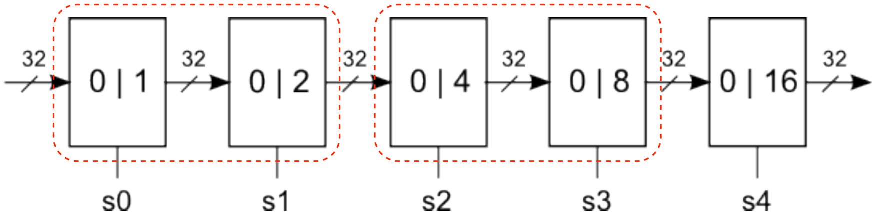
- Worst case delay O(N) , not good for processor design.
- Can we do it in O(logN) time and fit it in one cycle?

Log Shifter / Rotator

- Log(N) stages, each shifts (or not) by a power of 2 places, S=

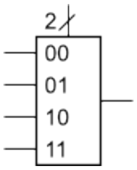


LUT Mapping of Log shifter

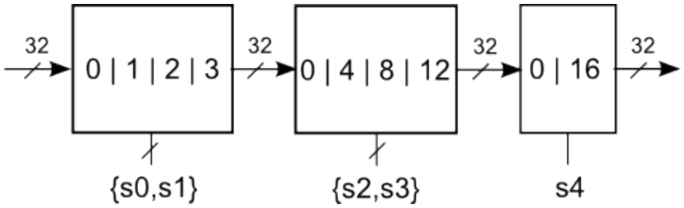


Efficient with 2to1 multiplexors, for instance, 3LUTs.

Virtex5 has 6LUTs. Naturally makes 4to1 muxes:



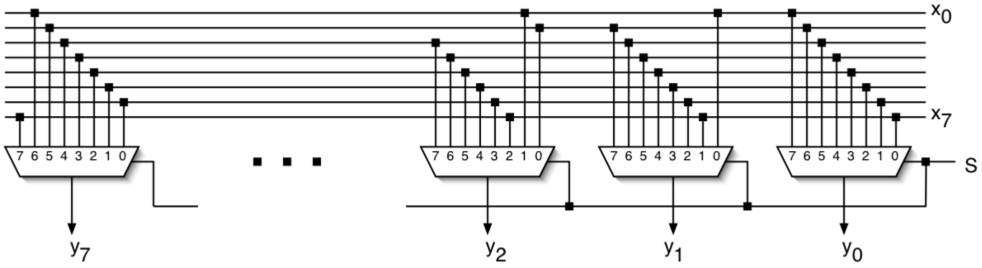
Reorganize shifter to use 4to1 muxes.



Final stage uses F7 mux

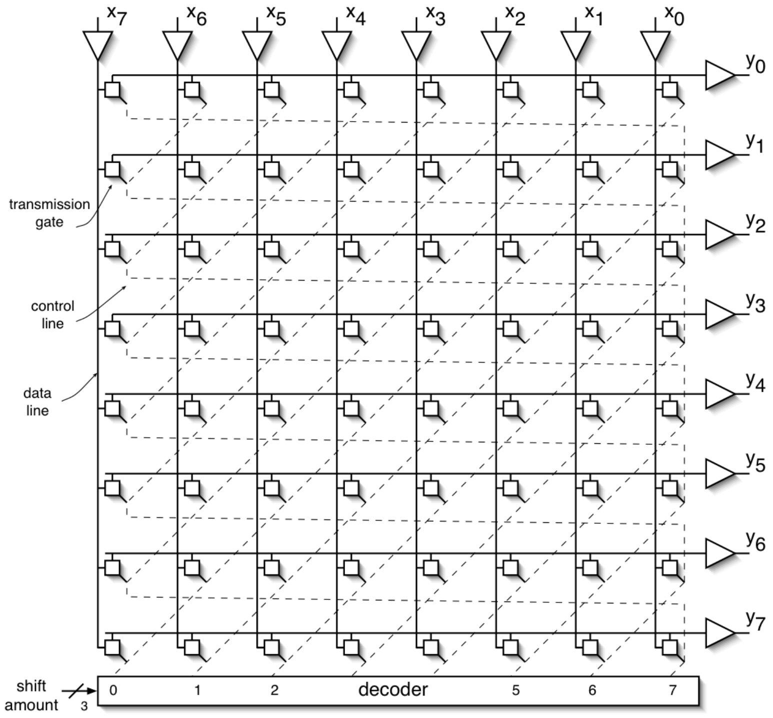
"Improved" Shifter / Rotator

- How about this approach? Could it lead to even less delay?



- What is the delay of these big muxes?
- Look a transistor-level implementation?

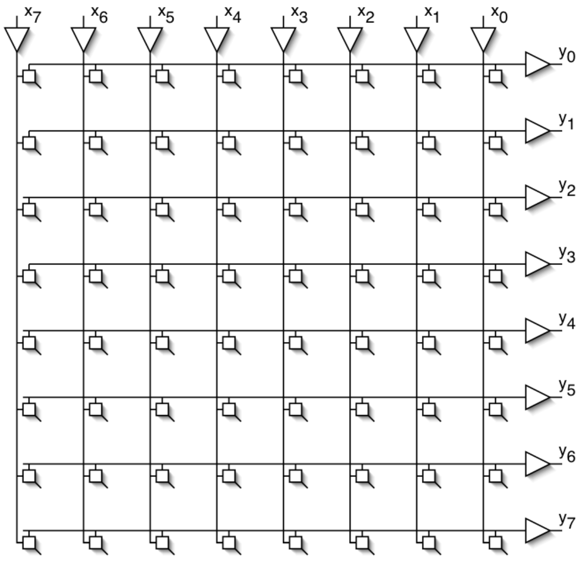
Barrel Shifter



Cost/delay?

- (don't forget the decoder)

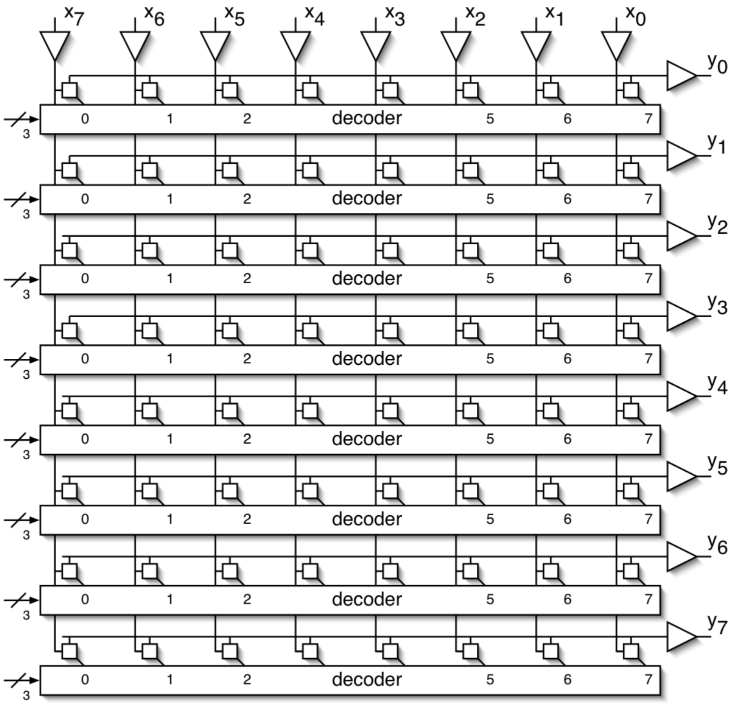
Connection Matrix



Generally useful structure:

- N^2 control points.
- What other interesting functions can it do?

Cross-bar Switch



- Nlog(N) control signals.
- Supports all interesting permutations
  - All one-to-one and one-to-many connections.
- Commonly used in communication hardware (switches, routers).