Lecture 15: Practical Bison: Error Handling, etc.	Identifying Errors		
<ul> <li>One purpose of the parser is to filter out errors that show up in parsing</li> </ul>	<ul> <li>All of the valid parsers we've seen identify syntax errors as soon as possible.</li> </ul>		
<ul> <li>Later stages should not have to deal with possibility of malformed constructs</li> </ul>	• Valid prefix property: all the input that is shifted or scanned is the beginning of some valid program		
<ul> <li>Parser must identify error so programmer knows what to correct</li> </ul>	<ul> <li>But the rest of the input might not be.</li> </ul>		
<ul> <li>Parser should recover so that processing can continue (and other errors found).</li> </ul>	<ul> <li>So in principle, deleting the lookahead (and subsequent symbols) and inserting others will give a valid program.</li> </ul>		
<ul> <li>Parser might even correct error (e.g., PL/C compiler could "correct" some Fortran programs into equivalent PL/1 programs!)</li> </ul>			
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Automating Recovery	Bison's Technique		
<ul> <li>Unfortunately, best results require using semantic knowledge and hand tuning.</li> </ul>	<ul> <li>The special terminal symbol error is never actually returned by the lexer.</li> </ul>		
<ul> <li>E.g., a(i].y = 5 might be turned to a[i].y = 5 if a is statically known to be a list, or a(i).y = 5 if a function.</li> </ul>	<ul> <li>Gets inserted by parser in place of erroneous tokens.</li> <li>Parsing then proceeds normally.</li> </ul>		
<ul> <li>Some automatic methods can do an OK job that at least allows parser to catch more than one error.</li> </ul>			

Example of Bison's Error Rules	Response to Error		
Suppose we want to throw away bad statements and carry on	• Consider erroneous text like		
stmt : whileStmt	if x y:		
ifStmt	<ul> <li>When parser gets to the y, will detect error.</li> </ul>		
   error NEWLINE ;	<ul> <li>Then pops items off parsing stack until it finds a state that allows a shift or reduction on 'error' terminal</li> </ul>		
	<ul> <li>Does reductions, then shifts 'error'.</li> </ul>		
	• Finally, throws away input until it finds a symbol it can shift after 'error', according to the grammar.		
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Error Response, contd.	Of Course, It's Not Perfect		
• So with our example: stmt : whileStmt	<ul> <li>"Throw away and punt" is sometimes called "panic-mode error recov- ery"</li> </ul>		
ifStmt 	• Results are often annoying.		
error NEWLINE ;	<ul> <li>For example, in our example, there could be an INDENT after the NEWLINE, which doesn't fit the grammar and causes another error.</li> </ul>		
We see 'y', throw away the 'if x', so as to be back to where a stmt can start.	<ul> <li>Bison compensates in this case by not reporting errors that are too close together</li> </ul>		
<ul> <li>Shift 'error' and throw away more symbols to NEWLINE. Then carry</li> </ul>	<ul> <li>But in general, can get cascade of errors.</li> </ul>		
on.	<ul> <li>Doing it right takes a lot of work.</li> </ul>		

## Bison Examples

[See lecture15 directory.]

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