

**CS61A SUMMER 2010**  
**FINAL REVIEW SESSION 2**

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Derived from the notes of Chung Wu, Justin Chen and Carolen,  
and past CS61A review sessions (Spring 2007)

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**QUESTION 1.**

What are the possible values of `x` after the following is executed:

```
(define x 10)
(parallel-execute (lambda () (set! x (+ 5 x)) (set! x (* x 3)))
                  (lambda () (if (> x 16)
                                (set! x 100)
                                (set! x (- x 20)))))
```

**QUESTION 2.** (Question 13 of the Final Exam, Spring 2003)

Given the following definitions:

```
(define s (make-serializer))
(define t (make-serializer))
(define x 10)
(define (f) (set! x (+ x 3)))
(define (g) (set! x (* x 2)))
```

Can the following expressions produce an incorrect result, a deadlock, or neither? (By "incorrect result" we mean a result that is not consistent with some sequential ordering of the processes.)

- (a) `(parallel-execute (s f) (t g))`
- (b) `(parallel-execute (s f) (s g))`
- (c) `(parallel-execute (s (t f)) (t g))`
- (d) `(parallel-execute (s (t f)) (s g))`
- (e) `(parallel-execute (s (t f)) (t (s g)))`

**QUESTION 3.**

Which of the following interactions will execute faster or the same in the analyzing evaluator than in the original metacircular evaluator? Circle FASTER or SAME for each.

```
> (define (gauss-recur n) ;; sum of #s from 1 to n
      (if (= n 1)
          1
          (+ n (gauss-recur (- n 1)))))
> (gauss-recur 1000)
```

Analyzing will be:                      FASTER                      SAME

```
> (define (gauss n)
      (/ (* (+ n 1) n) 2))
> (gauss 1000)
```

Analyzing will be:                      FASTER                      SAME

**QUESTION 4.**

What are the first seven terms of the following stream definition?

```
(define mystery (cons-stream 1
                              (cons-stream 2
                                             (stream-map (lambda (x y) (+ x (* 2 y)))
                                                         mystery
                                                         (stream-cdr mystery)))))
```

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### QUESTION 5.

Write code to generate **cxr-stream**, the stream of all the possible combinations of **car** and **cdr**:

```
(car cdr caar cdar cadr cddr ...)
```

Each element of the stream is a *procedure*, so that, for example, we can write statements such as

```
((stream-ref cxr-stream 2) '((4 5) (foo bar)))
```

that would work (and in this case, return 4). You may find the **compose** and **interleave** functions useful here.

### QUESTION 6.

Ben Bitdiddle has conveniently defined **stream-accumulate** for you below:

```
(define (stream-accumulate combiner null-value s)
  (if (stream-null? s)
      null-value
      (combiner (stream-car s)
                (stream-accumulate combiner null-value
                                   (stream-cdr s)))))
```

What happens when we do:

(a) (define foo (stream-accumulate + 0 integers))

(b) (define bar (cons-stream 1 (stream-accumulate + 0 integers)))

(c) (define baz (stream-accumulate (lambda (x y) (cons-stream x y))
 the-empty-stream
 integers))

### QUESTION 7.

We would like to implement a cheating detection system for tests. We start by simulating a row of test-takers by a vector, where each element in the vector is the answers to each student's test. These answers are also simulated by a vector, where each index corresponds to a problem number, and the element at the index is the student's solutions to the problem. For example:

```
##( #('cs61a 8 'none)      ;; student 0 answered 'cs61a, 8, 'none
    #('mother 6 'a)
    #('cs61a 5 'b)      )
```

(The answers are aligned here for your benefit.) Here, the test requires three answers, and student 0 answered 'cs61a, 8, and 'none to the three questions. Student 1 is the only person with two neighbors.

Two students are suspected of cheating if they have at least half of the answers identical to a neighbor. So if there are three questions on the test, if two are identical in consecutive students, the students are suspected of cheating. Write **catch-cheaters** that takes in a vector of vectors and returns the index of the first student suspected of cheating. (So if students 1 and 2 cheated off each other, return 1.)

**QUESTION 8.**

At lines A, B, C, and D, how many times have + and \* been called in lazy evaluation and in applicative order evaluation?

```
(define (foo n m)
  (if (> n 10)
      (begin (display m) n)
      (begin (display n) m)))
```

Lazy> (define z (\* 8 4))

A

Lazy> (define x ((lambda (x) x) (+ 2 2)))

B

Lazy> (define y (foo z x))

C

Lazy> y

D

**QUESTION 9.**

Consider the following interactions in the lazy evaluator:

```
(define w 100)
(define (foo x y) (x y))
(define q (foo (lambda (z) (set! w 50) z)
              (begin (set! W 10) 3))))
```

What are the values of the following statements typed at the prompt immediately after?

(a) w

(b) q

(c) w

**QUESTION 10.**

A *magic square* is an arrangement of numbers in a square matrix, where the sum of each row, each column, and each main diagonal is the same. For example, here is a 3x3 magic square:

```
2 7 6 (For the 3x3 square, each row, column, and diagonal sums to 15)
9 5 1
4 3 8
```

Write a procedure called **magic-square** that uses the nondeterministic evaluator to find 3x3 magic square configurations.

In this implementation, a magic square is a list of lists, where each list is associated with a row of the magic square. So, for example, the magic square above is represented as

```
(list (list 2 7 6) (list 9 5 1) (list 4 3 8)).
```

If it helps, you may assume the **distinct?** procedure from your homework. Also, the = primitive can take more than two arguments, and returns true if all of the arguments are equal. Thus,

```
(= 3 3 3) is true, while
(= 3 3 4) is false.
```

**QUESTION 11.**

Consider the following Scheme program:

```
(let ((a (amb 1 2 3))
      (b (amb 4 5 6)))
  (display "hello")
  (require (= b (* a 2)))
  a)
```

How many times will *hello* be printed? What is the return value?

**QUESTION 12.**

What are the results of the following statements when entered into the nondeterministic evaluator? Write down all of the results after multiple **try-again** statements, until the evaluator claims that there are no more solutions.

(a) (amb 1 2 3)

(b) (amb (list 1 2 3))

(c) (amb 1 (amb 2 (amb 3)))

(d) (amb (amb 1) (amb 2) (amb 3))

(e) (amb (amb 2 3) 1 (amb 4))

```
(f) (define (foo x)
      (cond ((not (pair? x)) (amb))
            ((word? (cdr x)) (cdr x))
            (else (amb (foo (car x))
                       (foo (cdr x))))))
```

(foo '(a (b c) (d e . f) (g (h . i) j) k))

**QUESTION 13.**

Rotating lists is fun, so let's keep doing it! For the following, assume that only the rule **append** has been defined, as in the lecture.

Implement a rule **rotate-forward** so that

```
(rotate-forward (1 2 3 4) ?what) ==> ?what = (2 3 4 1).
```

That is, the second list is the first list with the first element attached to the end instead. Assume the list is non-empty.

Let's get both sides of the story. We'd like a rule **rotate-backward** so that

```
(rotate-backward (1 2 3 4) ?what) ==> ?what = (4 1 2 3)
```

That is, the second list is the first list with the last element attached to the front instead. You may define other helper rules if you'd like.

**QUESTION 14.**

Write a rule or rules to determine if one integer is less than another. For example, the query

```
(less ?x (a a a))
```

should give the results

```
(less () (a a a)) (less (a) (a a a)) (less (a a) (a a a))
```