

CS70: Jean Walrand: Lecture 25.

Markov Chains: Distributions

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Markov Chains: Distributions

1. Review
2. Distribution
3. Irreducibility
4. Convergence

Review

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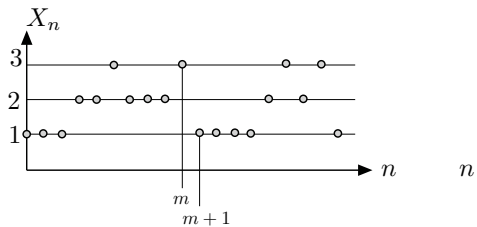
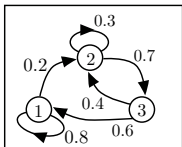
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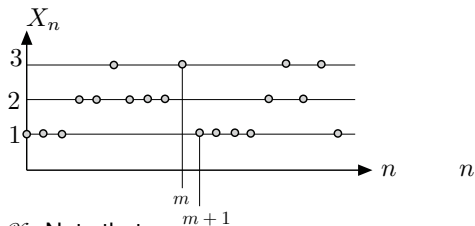
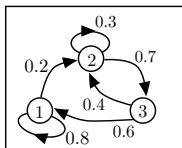
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Distribution of X_n

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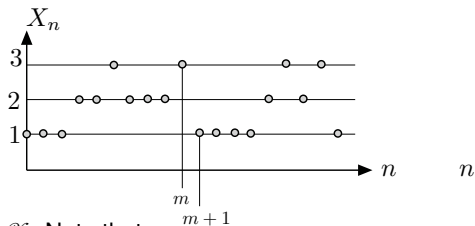
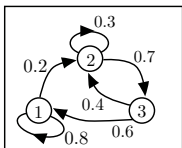


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Let $\pi_m(i) = Pr[X_m = i], i \in \mathcal{X}$. Note that

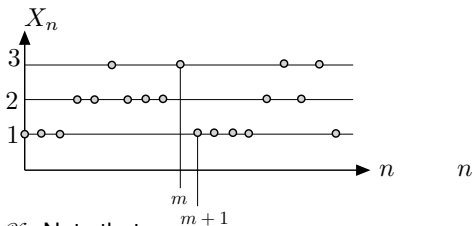
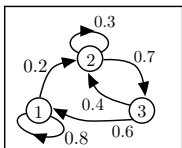
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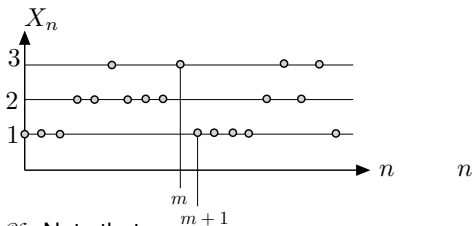
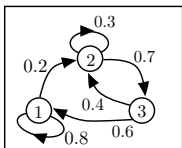
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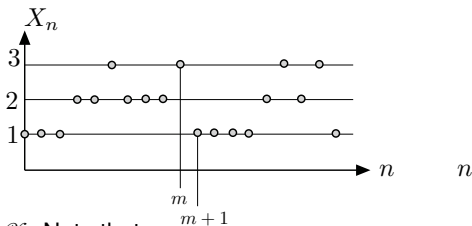
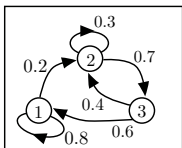
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Hence,

$$\pi_{m+1}(j) = \sum_i \pi_m(i) P(i, j), \forall j \in \mathcal{X}.$$

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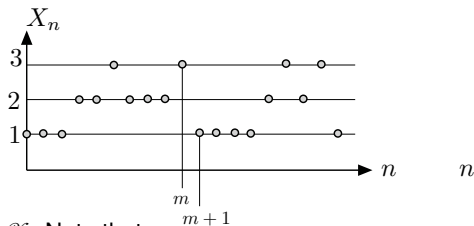
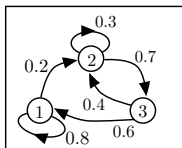
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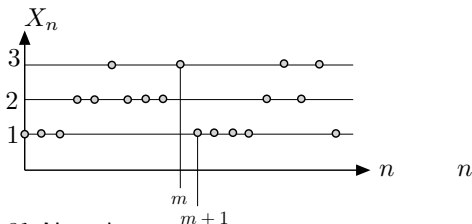
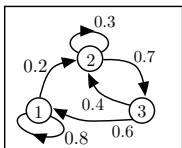
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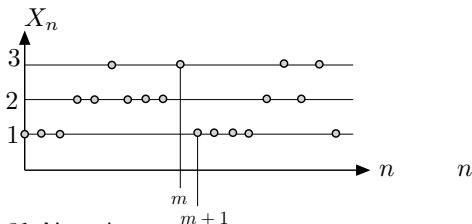
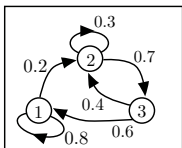
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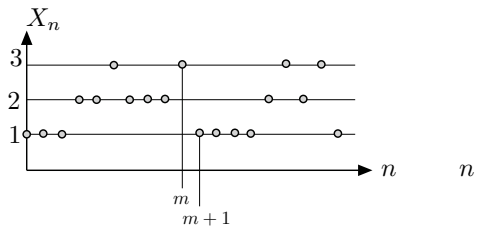
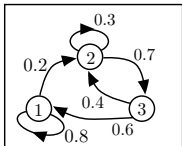
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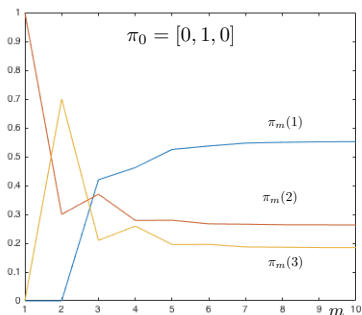
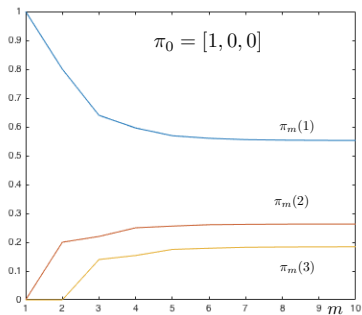
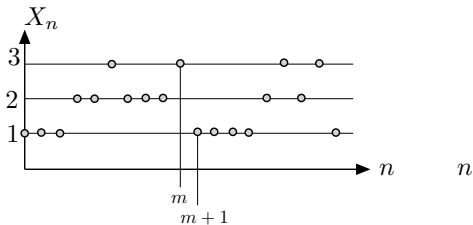
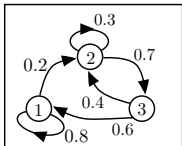
Thus, $\pi_1 = \pi_0 P, \pi_2 = \pi_1 P = \pi_0 P^2, \dots$ Hence,

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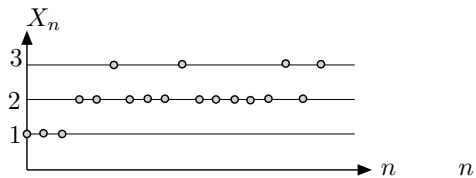
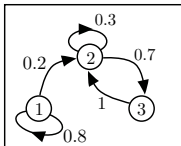
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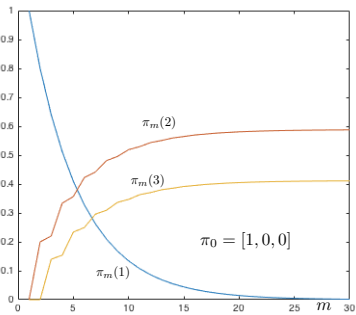
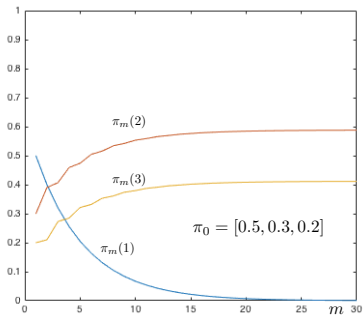
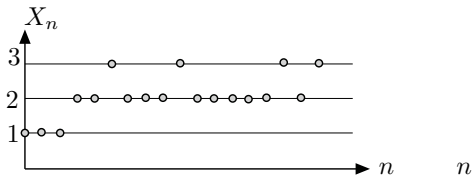
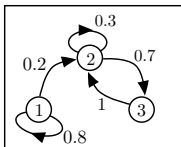
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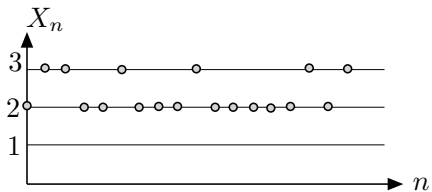
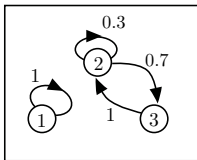
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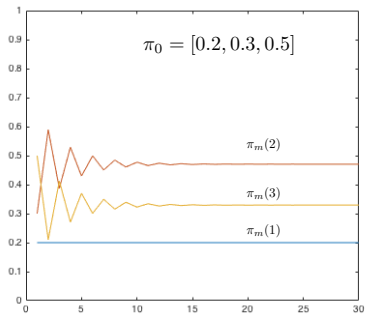
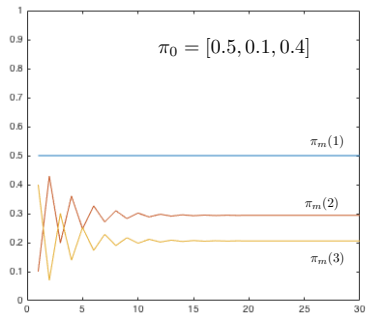
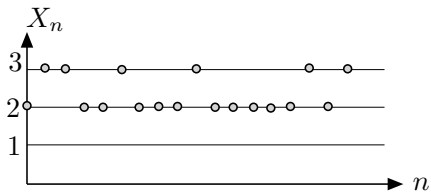
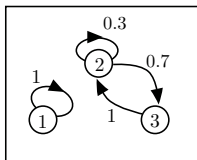
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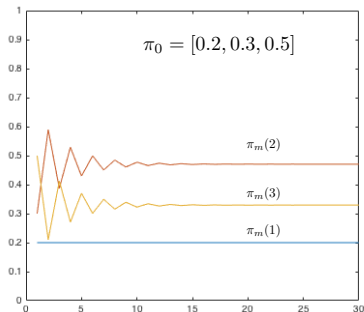
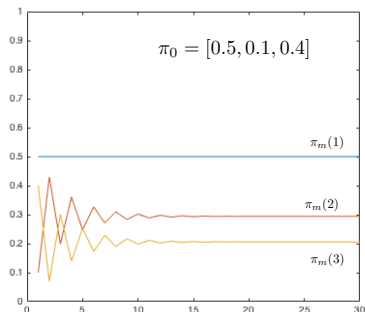
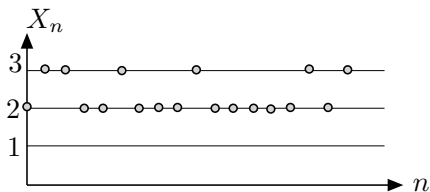
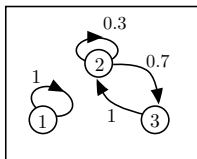
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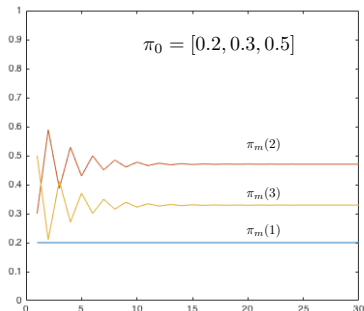
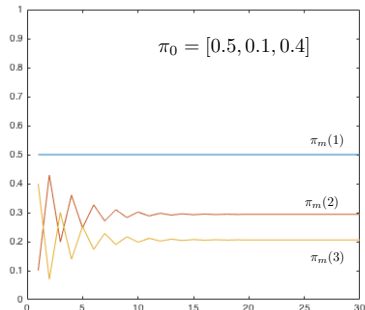
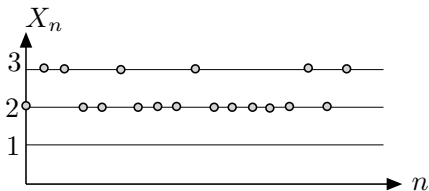
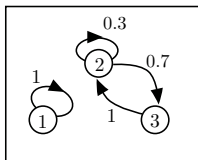


Distribution of X_n



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Distribution of X_n



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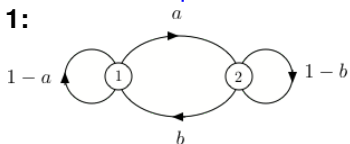
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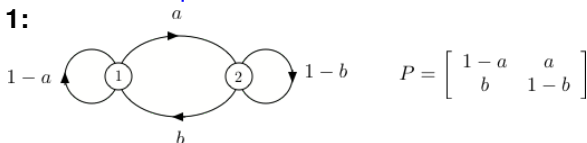


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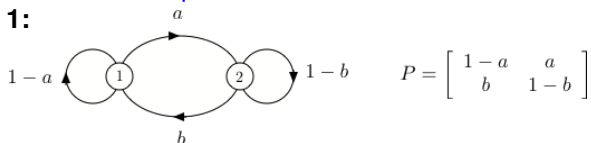
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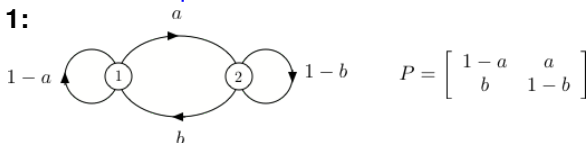
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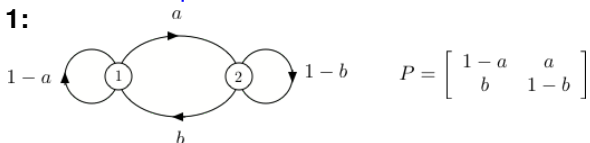


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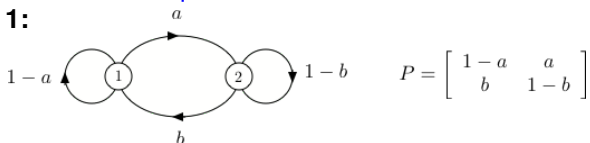
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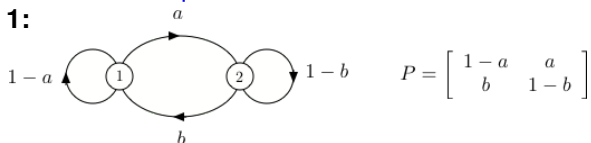
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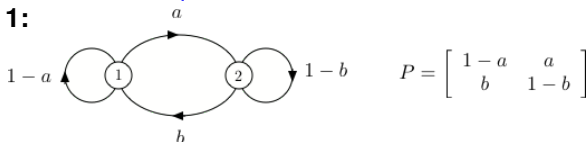
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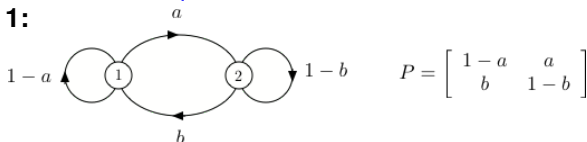
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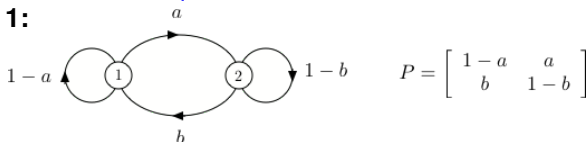
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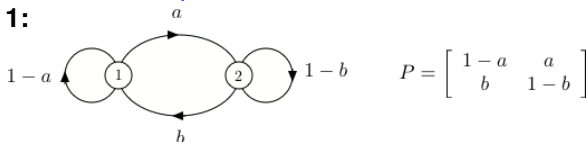
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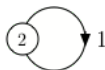


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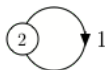
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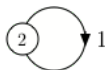
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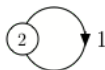
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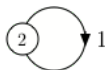
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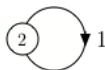
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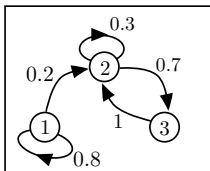
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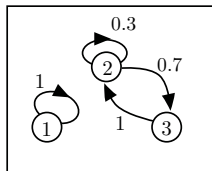
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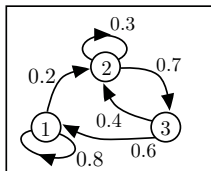
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[A]



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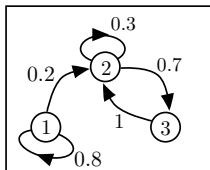


[C]

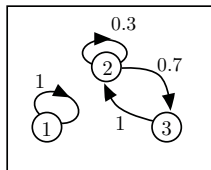
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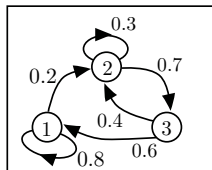
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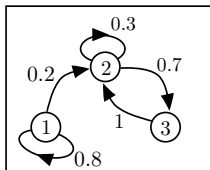
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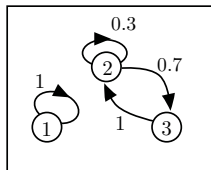
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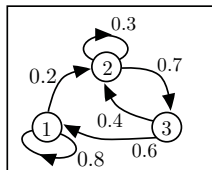
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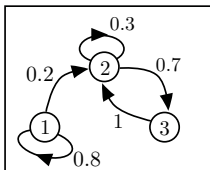
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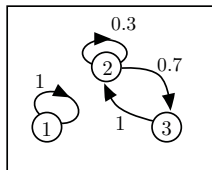
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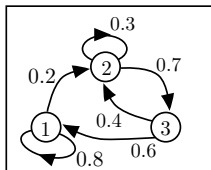
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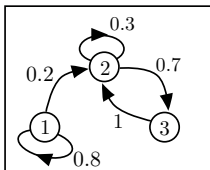
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[A] is **not irreducible**. It cannot go from (2) to (1).

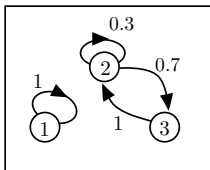
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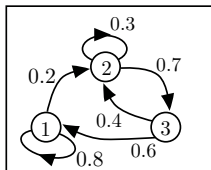
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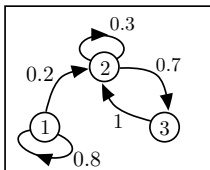
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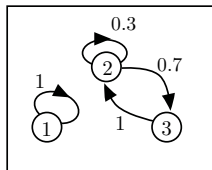
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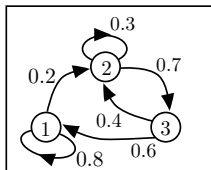
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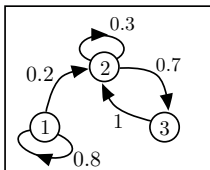
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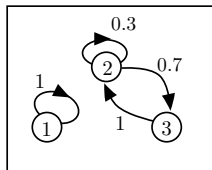
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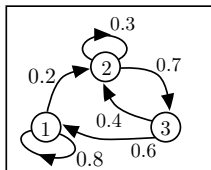
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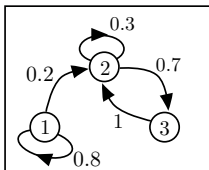
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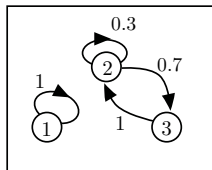
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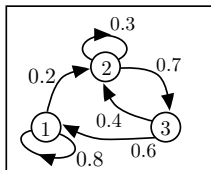
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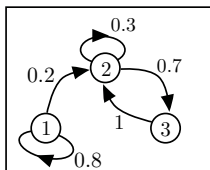
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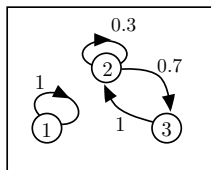
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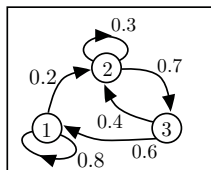
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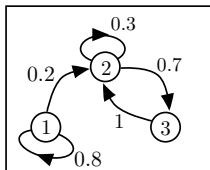
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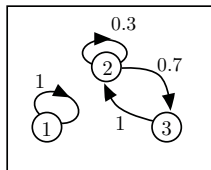
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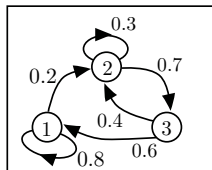
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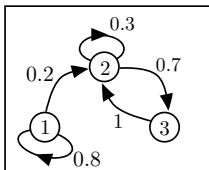
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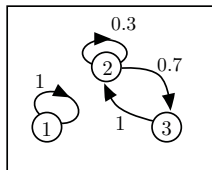
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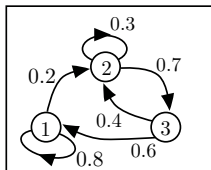
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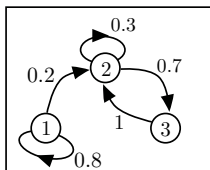
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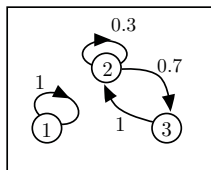
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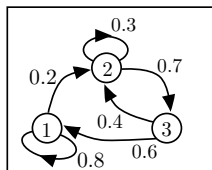
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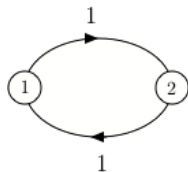
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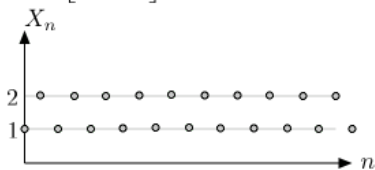
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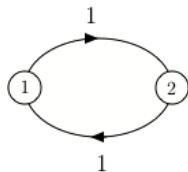
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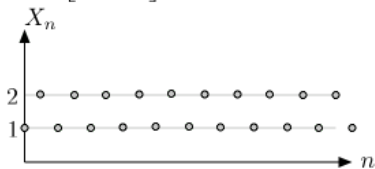
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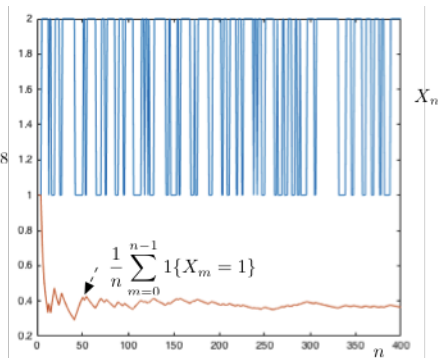
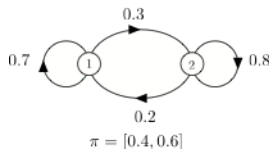
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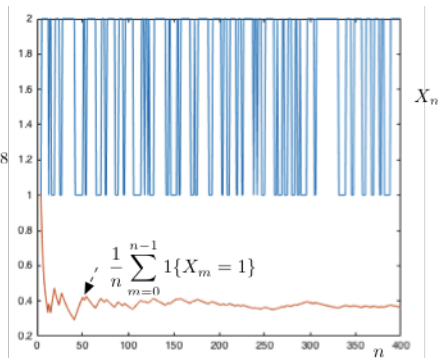
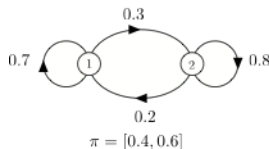
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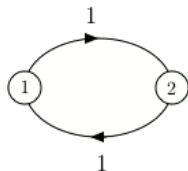
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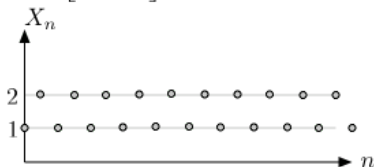
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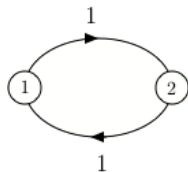
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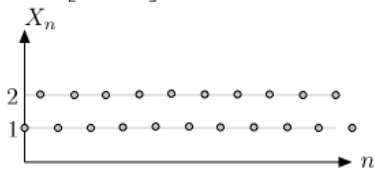
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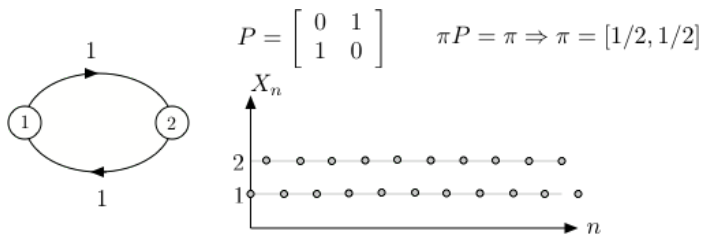


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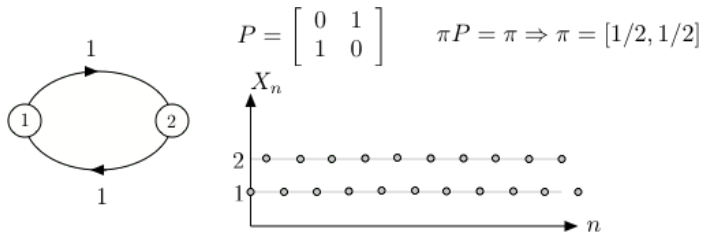


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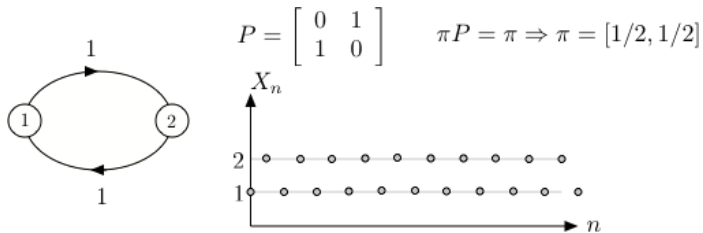


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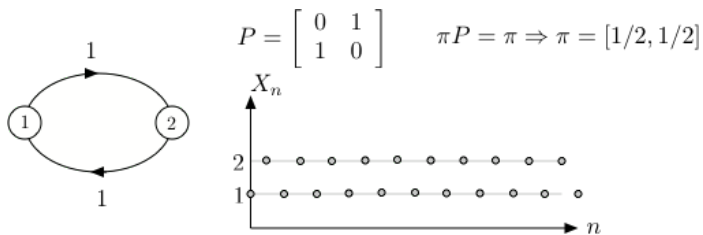


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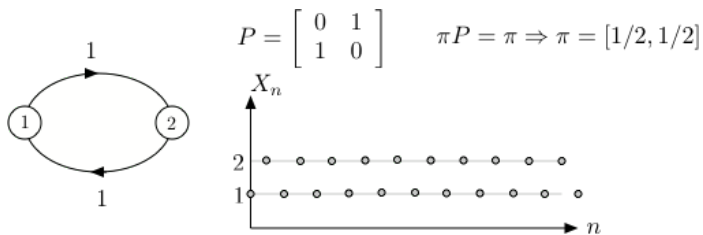
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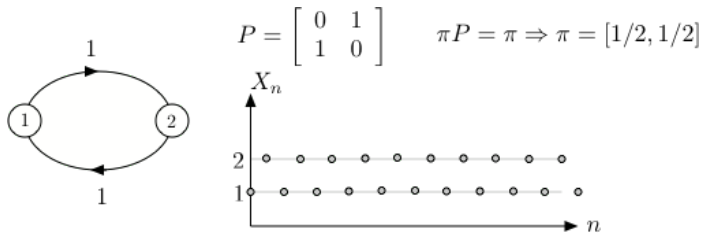
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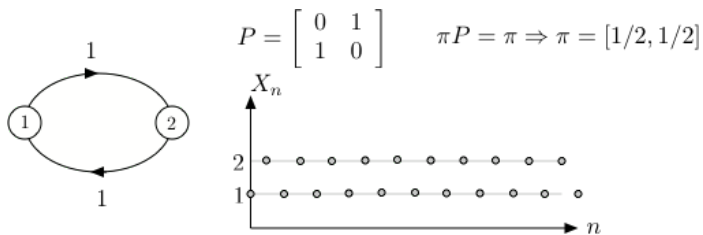
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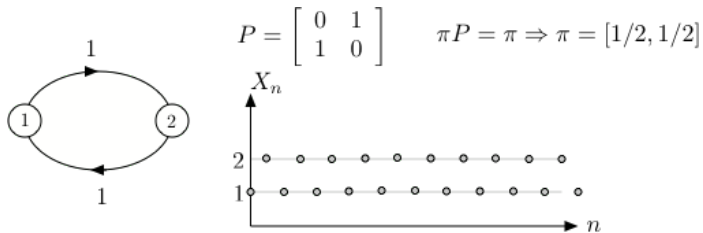
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Assume $X_0 = 1$. Then $X_1 = 2, X_2 = 1, X_3 = 2, \dots$

Thus, if $\pi_0 = [1, 0], \pi_1 = [0, 1], \pi_2 = [1, 0], \pi_3 = [0, 1], \dots$, etc.

Hence, π_n does not converge to $\pi = [1/2, 1/2]$.

Periodicity

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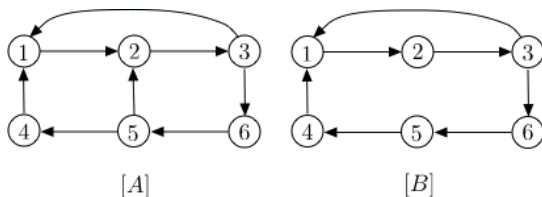
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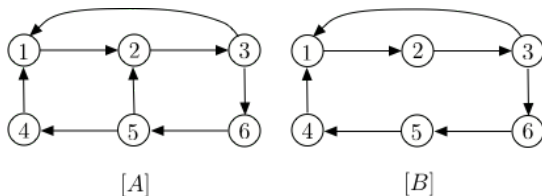
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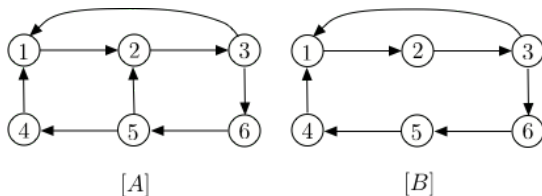
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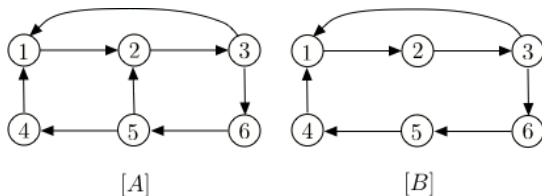
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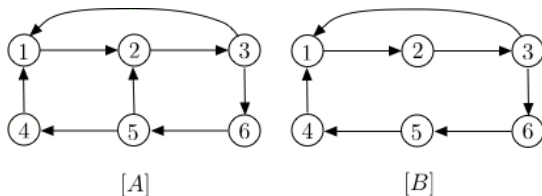
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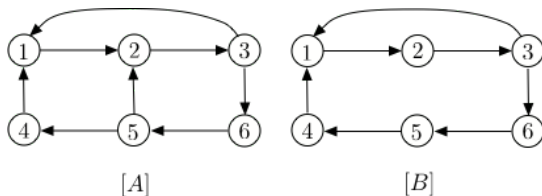
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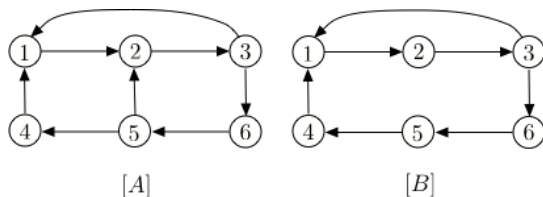
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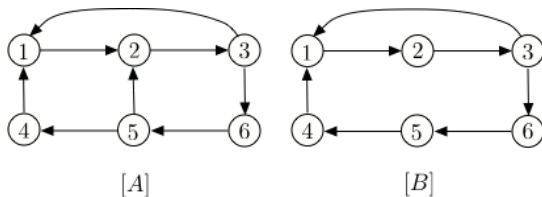
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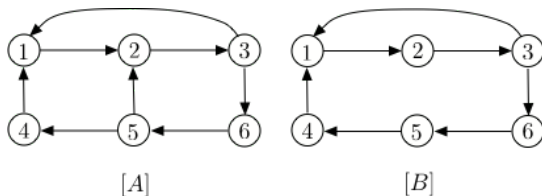
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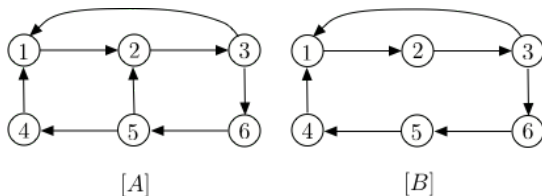
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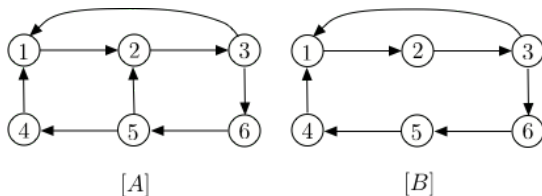
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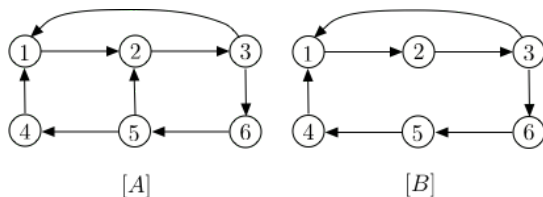
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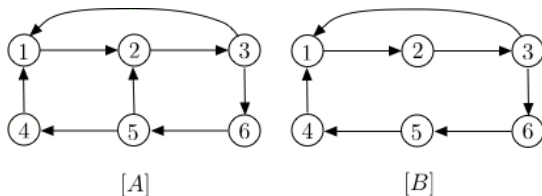
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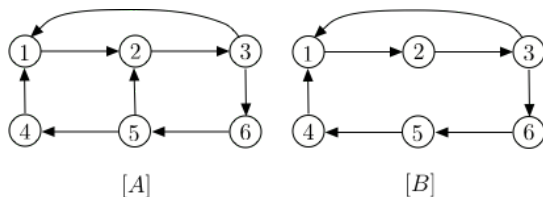
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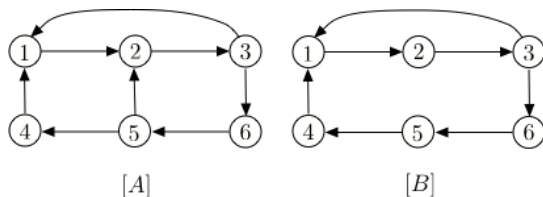
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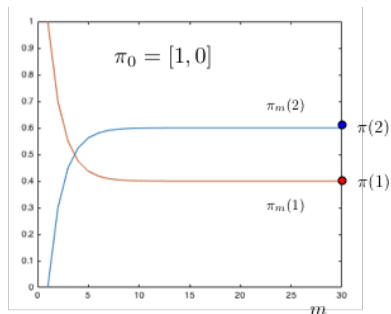
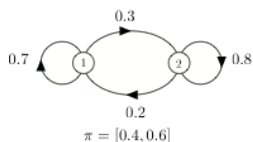
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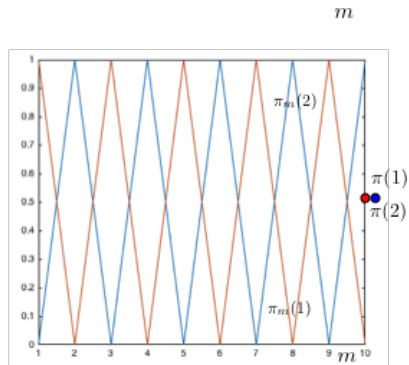
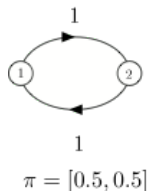
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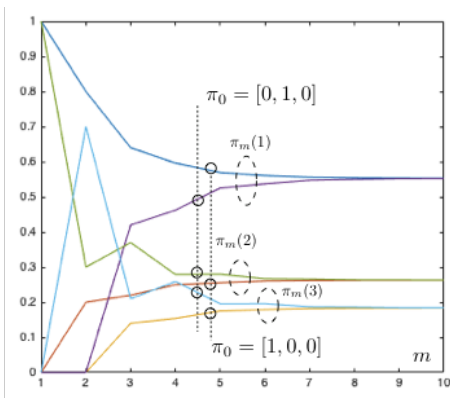
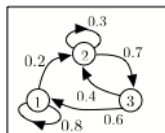
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Let P be irreducible. How do we find π ?

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There are relatively few problems for which one can prove such a clean result. However, there is a **systematic approach to calculate the optimal strategy** for many problems. We explain that approach next on this problem.

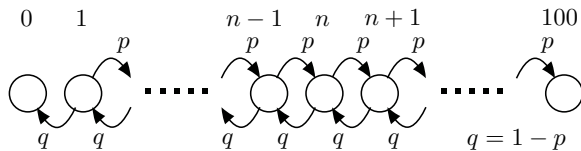
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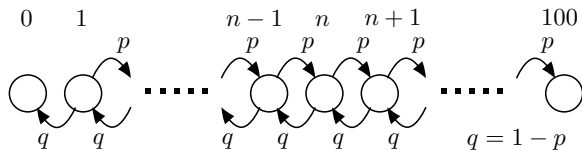
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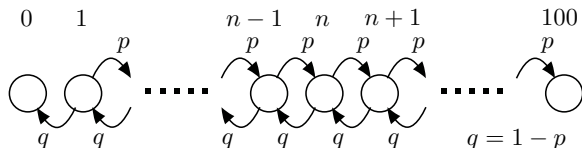
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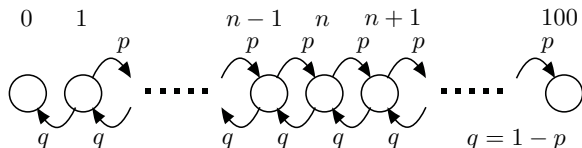


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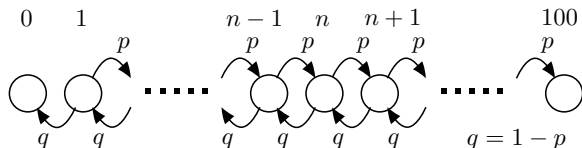


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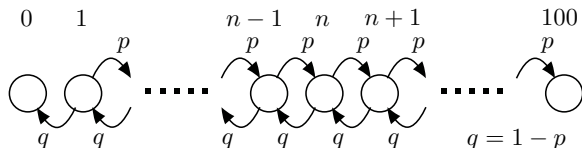


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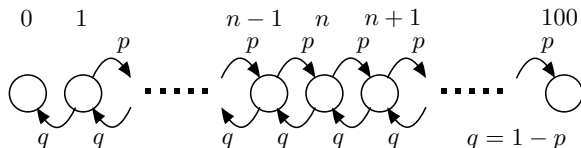


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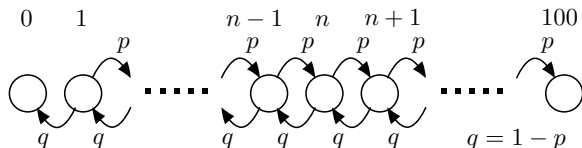
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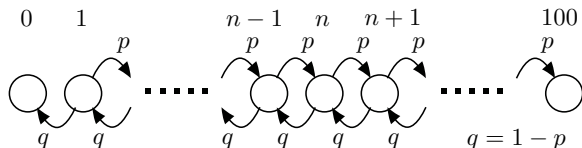
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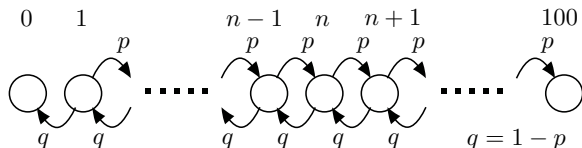
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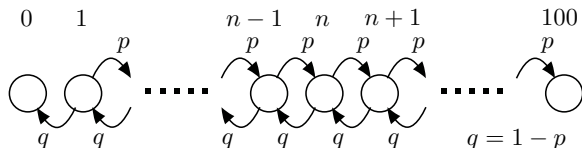
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Thus, the probability of winning the game (i.e., getting to 100 before 0) is at least 0.0448 when playing bold.

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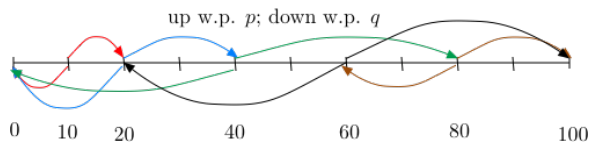
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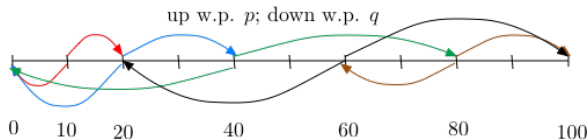
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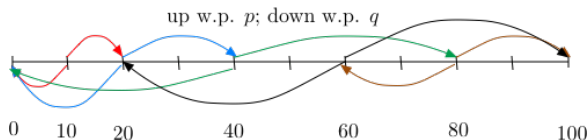
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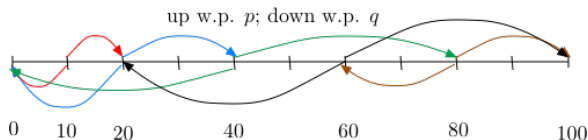


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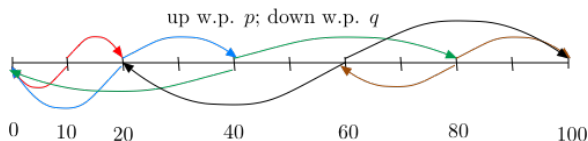


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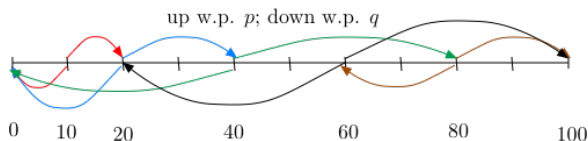


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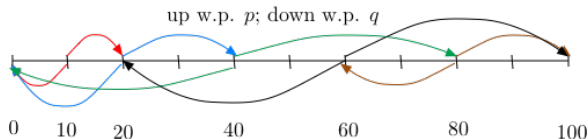


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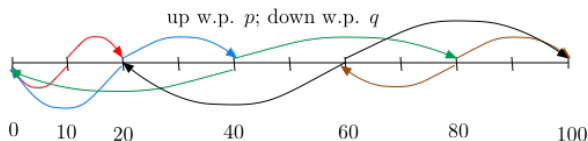
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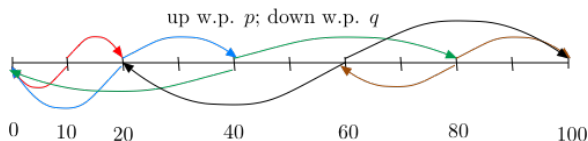
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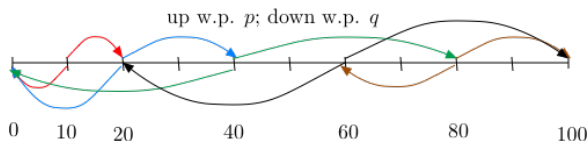
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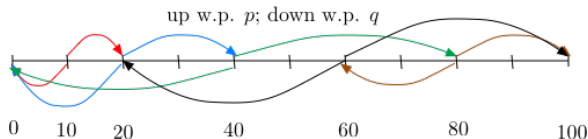
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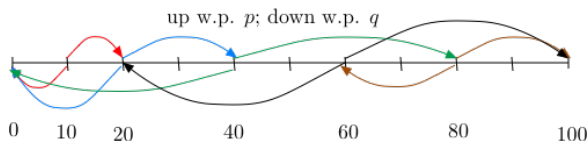
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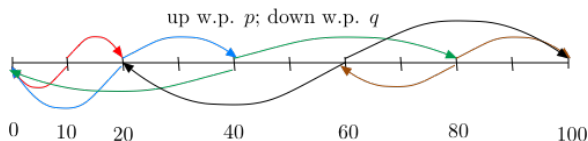
$$\alpha(10) = p\alpha(20) + q0; \alpha(20) = p\alpha(40) + q0; \alpha(40) = p\alpha(80) + q0$$
$$\alpha(80) = p1 + q\alpha(60); \alpha(60) = p1 + q\alpha(20)$$

To solve, let $\alpha(10) = x$. Then, we find

$$\alpha(20) = p^{-1}x; \alpha(40) = p^{-1}\alpha(20) = p^{-2}x$$
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Bold: Analysis

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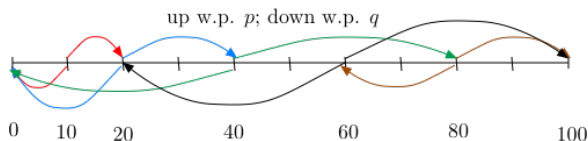
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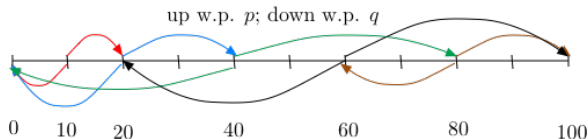
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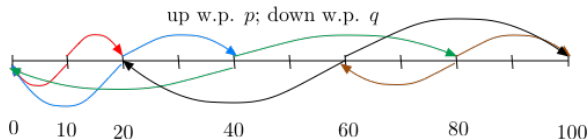
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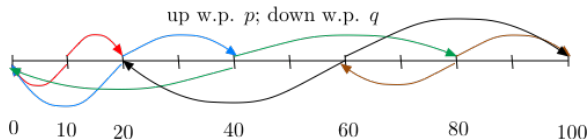
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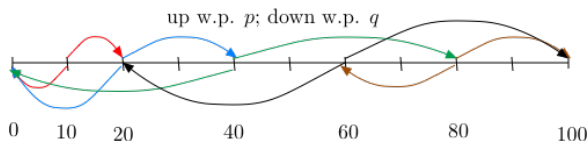
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We find $x = p^2(1 + q)/(p^2 - q^2)$

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We solve the last two equations for x .

We find $x = p^2(1 + q)/(p^{-2} - q^2) \approx 0.0735$.

Optimal Strategy

Optimal Strategy

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Here is a systematic approach. Assume you can only play 0 time.

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Let $V(k, n)$ be the maximum probability of winning the game if we can play k times

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Then,

$$V(k+1, n) = \max\{pV(k, n+m) + qV(k, n-m) \mid m \leq n \text{ and } n+m \leq 100\}.$$

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We can solve successively for $V(0, \cdot)$, $V(1, \cdot)$, $V(2, \cdot)$, \dots

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We can solve successively for $V(0, \cdot)$, $V(1, \cdot)$, $V(2, \cdot)$, \dots . In the limit, we find the best strategy. The program shows that bold is optimal when $p < 0.5$.

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We can solve successively for $V(0, \cdot)$, $V(1, \cdot)$, $V(2, \cdot)$, \dots . In the limit, we find the best strategy. The program shows that bold is optimal when $p < 0.5$. The finer result (Dubins and Savage) is to show this analytically.

Another Game

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Consider the following game.

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For $m \leq n$, let (n, m) mean that there are still n cards and m aces left in the deck.

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For $m \leq n$, let (n, m) mean that there are still n cards and m aces left in the deck. Let also $V(n, m)$ be the maximum probability of winning this game in that situation.

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The first term corresponds to stating 'the next card is an ace.'

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The first term corresponds to stating 'the next card is an ace.' The second term corresponds to not deciding yet.

Another Game

Consider the following game. One has a perfectly shuffled 52-card deck. The cards are turned over one at a time. You win if you can guess when the next card will be an ace. You can only guess once. What is the best strategy? Should you let a few cards go by, then decide that the next one will be an ace?

For $m \leq n$, let (n, m) mean that there are still n cards and m aces left in the deck. Let also $V(n, m)$ be the maximum probability of winning this game in that situation. Then,

$$V(n, m) = \max\left\{\frac{m}{n}, \frac{m}{n} V(n-1, m-1) + \frac{n-m}{n} V(n-1, m)\right\}.$$

The first term corresponds to stating 'the next card is an ace.' The second term corresponds to not deciding yet.

One boundary condition is $V(n, 0) = 0$.

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Conclusion: You might as well stop at the first card!.

Markov Decision Problems

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See EE126, CS188, EE223.