1. Modular Arithmetic.

1. Modular Arithmetic. Clock Math!!!

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor.

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!
- 3. Euclid's GCD Algorithm.

- 1. Modular Arithmetic. Clock Math!!!
- 2. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!
- 3. Euclid's GCD Algorithm. A little tricky here!

If it is 1:00 now.

If it is 1:00 now. What time is it in 2 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

Custom is only to use the representative in $\{12, 1, \dots, 11\}$

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

Custom is only to use the representative in $\{12, 1, ..., 11\}$ (Almost remainder, except for 12 and 0 are equivalent.)

Today is Monday.

Today is Monday. What day is it a year from now?

Today is Monday.

What day is it a year from now? on February 9, 2016?

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2. 5 days from now.

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.

25 days from now.

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27 or day 6.

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27 or day 6.
two days are equivalent up to addition/subtraction of multiple of 7.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27 or day 6.
two days are equivalent up to addition/subtraction of multiple of 7.
11 days from now is day 6

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27 or day 6.
two days are equivalent up to addition/subtraction of multiple of 7.
11 days from now is day 6 which is Saturday!

What day is it a year from now?

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27 or day 6.
two days are equivalent up to addition/subtraction of multiple of 7.
11 days from now is day 6 which is Saturday!

What day is it a year from now? This year is leap year.

Today is Monday.
What day is it a year from now? on February 9, 2016?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 2.
5 days from now. day 7 or day 0 or Sunday.
25 days from now. day 27 or day 6.
two days are equivalent up to addition/subtraction of multiple of 7.
11 days from now is day 6 which is Saturday!

What day is it a year from now? This year is leap year. So 366 days from now.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now? This year is leap year. So 366 days from now. Day 2+366 or day 368.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now? This year is leap year. So 366 days from now. Day 2+366 or day 368. Smallest representation:

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now? This year is leap year. So 366 days from now. Day 2+366 or day 368.

Smallest representation:

subtract 7 until smaller than 7.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now? This year is leap year. So 366 days from now. Day 2+366 or day 368. Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now?

This year is leap year. So 366 days from now.

Day 2+366 or day 368.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

368/7

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now?

This year is leap year. So 366 days from now.

Day 2+366 or day 368.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

368/7 leaves quotient of 52 and remainder 4.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now?

This year is leap year. So 366 days from now.

Day 2+366 or day 368.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

368/7 leaves quotient of 52 and remainder 4.

or February 9, 2017 is a Thursday.

Today is Monday. What day is it a year from now? on February 9, 2016? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 2.

5 days from now. day 7 or day 0 or Sunday.

25 days from now. day 27 or day 6.

two days are equivalent up to addition/subtraction of multiple of 7.

11 days from now is day 6 which is Saturday!

What day is it a year from now?

This year is leap year. So 366 days from now.

Day 2+366 or day 368.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

368/7 leaves quotient of 52 and remainder 4.

or February 9, 2017 is a Thursday.

80 years from now?

80 years from now? 20 leap years.

80 years from now? 20 leap years. 366×20 days

80 years from now? 20 leap years. 366×20 days 60 regular years.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7?

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7?
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80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2. What is remainder of 365 when dividing by 7? 1
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80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2. What is remainder of 365 when dividing by 7? 1
```

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
Today is day 2.
```

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

```
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
What is remainder of 365 when dividing by 7? 1
Today is day 2.
```

Get Day: $2 + 2 \times 20 + 1 \times 60$

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 2. Get Day: $2+2 \times 20+1 \times 60 = 102$

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 2.

```
Get Day: 2 + 2 \times 20 + 1 \times 60 = 102
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```
Remainder when dividing by 7?
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80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 2. Get Day: $2+2 \times 20+1 \times 60 = 102$ Remainder when dividing by 7? $102 = 14 \times 7$

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 2. Get Day: $2+2 \times 20+1 \times 60 = 102$ Remainder when dividing by 7? $102 = 14 \times 7 + 4$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 2. Get Day: $2+2 \times 20+1 \times 60 = 102$ Remainder when dividing by 7? $102 = 14 \times 7 + 4$. Or February 9, 2096 is Thursday!

```
80 years from now? 20 leap years. 366 \times 20 days
60 regular years. 365 \times 60 days
Today is day 2.
It is day 2 + 366 \times 20 + 365 \times 60. Equivalent to?
Hmm.
What is remainder of 366 when dividing by 7? 52 \times 7 + 2.
```

```
What is remainder of 365 when dividing by 7? 1
```

```
Today is day 2.
```

```
Get Day: 2 + 2 \times 20 + 1 \times 60 = 102
```

```
Remainder when dividing by 7? 102 = 14 \times 7 + 4.
```

```
Or February 9, 2096 is Thursday!
```

Further Simplify Calculation:

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2+366 \times 20+365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 2.

Get Day: $2 + 2 \times 20 + 1 \times 60 = 102$

Remainder when dividing by 7? $102 = 14 \times 7 + 4$.

Or February 9, 2096 is Thursday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 2.

Get Day: $2 + 2 \times 20 + 1 \times 60 = 102$

Remainder when dividing by 7? $102 = 14 \times 7 + 4$.

Or February 9, 2096 is Thursday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 2.

Get Day: $2+2 \times 20 + 1 \times 60 = 102$ Remainder when dividing by 7? $102 = 14 \times 7 + 4$. Or February 9, 2096 is Thursday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

What is remainder of 365 when dividing by 7? 1

Today is day 2.

Get Day: $2 + 2 \times 20 + 1 \times 60 = 102$

Remainder when dividing by 7? $102 = 14 \times 7 + 4$.

Or February 9, 2096 is Thursday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7. 60 has remainder 4 when divided by 7. Get Day: $2+2\times 6+1\times 4=18$.

Or Day 4.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

What is remainder of 365 when dividing by 7? 1

Today is day 2.

Get Day: $2 + 2 \times 20 + 1 \times 60 = 102$

Remainder when dividing by 7? $102 = 14 \times 7 + 4$.

Or February 9, 2096 is Thursday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

Or Day 4. February 9, 2095 is Thursday.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 2. It is day $2 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

What is remainder of 365 when dividing by 7? 1

Today is day 2.

Get Day: $2 + 2 \times 20 + 1 \times 60 = 102$

Remainder when dividing by 7? $102 = 14 \times 7 + 4$.

Or February 9, 2096 is Thursday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

Or Day 4. February 9, 2095 is Thursday.

"Reduce" at any time in calculation!

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

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Can calculate with representative in $\{0, \ldots, m-1\}$.

x (mod m) or mod(x,m)

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x (mod *m*) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29, 12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12

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 $x \pmod{m} \operatorname{or} \mod{(x,m)}$ $\operatorname{remainder of} x \operatorname{divided by} m \operatorname{in} \{0, \dots, m-1\}.$ $\operatorname{mod} (x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor \text{ is quotient.}$ $\operatorname{mod} (29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = \cancel{x} = 5$ Work in this system. $a \equiv b \pmod{m}.$

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Division: multiply by multiplicative inverse.

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Can solve $4x = 5 \pmod{7}$.

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Check! 4(3) = 12 = 5 \pmod{7}.
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"Common factor of 4" \implies 8k - 12l is a multiple of four for any l and k \implies

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For 8 modulo 12: no multiplicative inverse!

"Common factor of 4" \implies 8k - 12l is a multiple of four for any l and k \implies 8k \neq 1 (mod 12) for any k.

Greatest Common Divisor and Inverses.

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 $5x = 3 \pmod{6}$ What is x?

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For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6).

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd. $4x = 2 \pmod{6}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4...

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Very different for elements with inverses.

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How to find if x has an inverse modulo m?

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Next up.

Next up.

Next up. Euclid's Algorithm.

Next up.

Euclid's Algorithm. Runtime.

Next up.

Euclid's Algorithm. Runtime. Euclid's Extended Algorithm.

Does 2 have an inverse mod 8?

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Euclid's algorithm.

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(define (euclid x y)
  (if (= y 0)
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Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

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Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest." **Induction Step:** mod $(x, y) < y \le x$ when $x \ge y$ call in line (***) meets conditions plus arguments "smaller"

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

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Before discussing running time of gcd procedure...

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What is the value of 1,000,000?

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one million or 1,000,000!

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Euclid procedure is fast.

Theorem: (euclid x y) uses 2n "divisions" where $n = b(x) \approx \log_2 x$.

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 2^{n-1} divisions! Exponential dependence on size!

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101 bit number.

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Trying everything

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Check 2, check 3, check 4, check 5 ..., check y/2.

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```
euclid(700,568)
```

```
Trying everything
Check 2, check 3, check 4, check 5..., check y/2.
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euclid (700, 568)
```

```
euclid(568, 132)
```

```
Trying everything
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```

```
euclid(700,568)
euclid(568, 132)
euclid(132, 40)
```

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4
```

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Notice: The first argument decreases rapidly.

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Notice: The first argument decreases rapidly. At least a factor of 2 in two recursive calls.

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(The second is less than the first.)

Break.

Theorem: (euclid x y) uses O(n) "divisions" where n = b(x).

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Proof:

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First arg decreases by at least factor of two in two recursive calls. After $2\log_2 x = O(n)$ recursive calls, argument *x* is 1 bit number.

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1 division per recursive call.

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O(n) divisions.

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Proof of Fact: Recall that first argument decreases every call.

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Case 1: y < x/2, first argument is y

 \implies true in one recursive call;

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Case 2: Will show " $y \ge x/2$ " \implies "mod $(x, y) \le x/2$."

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mod(x, y) is second argument in next recursive call, and becomes the first argument in the next one.

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Proof:

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Proof of Fact: Recall that first argument decreases every call.

Case 2: Will show " $y \ge x/2$ " \implies "mod $(x, y) \le x/2$." When $y \ge x/2$, then

$$\lfloor \frac{x}{y} \rfloor = 1,$$

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$$\lfloor \frac{x}{y}
floor = 1$$

mod $(x,y) = x - y \lfloor \frac{x}{y}
floor =$

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Finding an inverse?

We showed how to efficiently tell if there is an inverse.

Finding an inverse?

We showed how to efficiently tell if there is an inverse. Extend euclid to find inverse.

Euclid's GCD algorithm.

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Computes the gcd(x, y) in O(n) divisions.

Euclid's GCD algorithm.

Computes the gcd(x, y) in O(n) divisions.

For x and m, if gcd(x, m) = 1 then x has an inverse modulo m.

Multiplicative Inverse.

GCD algorithm used to tell if there is a multiplicative inverse.

Multiplicative Inverse.

GCD algorithm used to tell **if** there is a multiplicative inverse. How do we **find** a multiplicative inverse?



Euclid's Extended GCD Theorem: For any *x*, *y* there are integers *a*, *b* such that

ax + by

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What is multiplicative inverse of x modulo m?

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By extended GCD theorem, when gcd(x, m) = 1.

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$$ax + bm = 1$$

 $ax \equiv 1 - bm \equiv 1 \pmod{m}$.

Euclid's Extended GCD Theorem: For any *x*, *y* there are integers *a*, *b* such that

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So a multiplicative inverse of $x \pmod{m}$!!

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(3)12 + (-1)35 = 1.

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(3)12 + (-1)35 = 1.a = 3 and b = -1.

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So *a* multiplicative inverse of $x \pmod{m}$!! Example: For x = 12 and y = 35, gcd(12,35) = 1.

(3)12 + (-1)35 = 1.

a = 3 and b = -1. The multiplicative inverse of 12 (mod 35) is 3.

gcd(35,12)

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
```

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
```

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

How did gcd get 11 from 35 and 12?

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

How did gcd get 11 from 35 and 12? $35 - \lfloor \frac{35}{12} \rfloor 12 = 35 - (2)12 = 11$

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
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gcd(1,0)
1
```

How did gcd get 11 from 35 and 12? $35 - \lfloor \frac{35}{12} \rfloor 12 = 35 - (2)12 = 11$

How does gcd get 1 from 12 and 11?

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

```
How did gcd get 11 from 35 and 12?

35 - \lfloor \frac{35}{12} \rfloor 12 = 35 - (2)12 = 11

How does gcd get 1 from 12 and 11?

12 - \lfloor \frac{12}{11} \rfloor 11 = 12 - (1)11 = 1
```

```
gcd(35,12)
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How did gcd get 11 from 35 and 12?

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How does gcd get 1 from 12 and 11?

12 - \lfloor \frac{12}{11} \rfloor 11 = 12 - (1)11 = 1
```

Algorithm finally returns 1.

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

```
How did gcd get 11 from 35 and 12?

35 - \lfloor \frac{35}{12} \rfloor 12 = 35 - (2)12 = 11

How does gcd get 1 from 12 and 11?

12 - \lfloor \frac{12}{11} \rfloor 11 = 12 - (1)11 = 1
```

Algorithm finally returns 1.

But we want 1 from sum of multiples of 35 and 12?

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

```
How did gcd get 11 from 35 and 12?

35 - \lfloor \frac{35}{12} \rfloor 12 = 35 - (2)12 = 11

How does gcd get 1 from 12 and 11?

12 - \lfloor \frac{12}{11} \rfloor 11 = 12 - (1)11 = 1
```

Algorithm finally returns 1.

But we want 1 from sum of multiples of 35 and 12?

Get 1 from 12 and 11.

```
gcd(35,12)
gcd(12, 11) ;; gcd(12, 35%12)
gcd(11, 1) ;; gcd(11, 12%11)
gcd(1,0)
1
```

```
How did gcd get 11 from 35 and 12?

35 - \lfloor \frac{35}{12} \rfloor 12 = 35 - (2)12 = 11

How does gcd get 1 from 12 and 11?

12 - \lfloor \frac{12}{11} \rfloor 11 = 12 - (1)11 = 1
```

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1 = 12 - (1)11

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Algorithm finally returns 1.

But we want 1 from sum of multiples of 35 and 12?

Get 1 from 12 and 11. 1 = 12 - (1)11 = 12 - (1)(35 - (2)12)Get 11 from 35 and 12 and plugin....

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1 = 12 - (1)11 = 12 - (1)(35 - (2)12) = (3)12 + (-1)35Get 11 from 35 and 12 and plugin.... Simplify.

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Extended GCD Algorithm.

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Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by.

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if y = 0 then return(x, 1, 0)
    else
        (d, a, b) := ext-gcd(y, mod(x,y))
        return (d, b, a - floor(x/y) * b)
```

Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by. Example:

ext-gcd(35,12)

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```
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Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by. Example: $a - \lfloor x/y \rfloor \cdot b =$

```
ext-gcd(35,12)
ext-gcd(12, 11)
ext-gcd(11, 1)
ext-gcd(1,0)
return (1,1,0) ;; 1 = (1)1 + (0) 0
```

Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by. Example: $a - \lfloor x/y \rfloor \cdot b = 1 - \lfloor 11/1 \rfloor \cdot 0 = 1$

```
ext-gcd(35,12)
  ext-gcd(12, 11)
    ext-gcd(11, 1)
    ext-gcd(1,0)
    return (1,1,0) ;; 1 = (1)1 + (0) 0
    return (1,0,1) ;; 1 = (0)11 + (1)1
```

```
Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by.
Example: a - \lfloor x/y \rfloor \cdot b = 0 - \lfloor 12/11 \rfloor \cdot 1 = -1
```

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ext-gcd(11, 1)
ext-gcd(1,0)
return (1,1,0) ;; 1 = (1)1 + (0) 0
return (1,0,1) ;; 1 = (0)11 + (1)1
return (1,1,-1) ;; 1 = (1)12 + (-1)11
```

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```

Claim: Returns (d, a, b): d = gcd(a, b) and d = ax + by. Example: $a - \lfloor x/y \rfloor \cdot b = \lfloor 35/12 \rfloor \cdot (-1) = 3$

```
ext-gcd(35,12)
ext-gcd(12, 11)
ext-gcd(11, 1)
ext-gcd(11, 0)
return (1,1,0) ;; 1 = (1)1 + (0) 0
return (1,0,1) ;; 1 = (0)11 + (1)1
return (1,1,-1) ;; 1 = (1)12 + (-1)11
return (1,-1, 3) ;; 1 = (-1)35 + (3)12
```

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return (1,1,-1) ;; 1 = (1)12 + (-1)11
return (1,-1, 3) ;; 1 = (-1)35 + (3)12
```

Theorem: Returns (d, a, b), where d = gcd(a, b) and

$$d = ax + by$$
.

Proof: Strong Induction.¹

¹Assume *d* is gcd(x, y) by previous proof.

Proof: Strong Induction.¹ **Base:** ext-gcd(x,0) returns (d = x,1,0) with x = (1)x + (0)y.

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$$d = ay + b \cdot (\mod(x, y))$$
$$= ay + b \cdot (x - \lfloor \frac{x}{y} \rfloor y)$$

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(

$$d = ay + b \cdot (\mod(x, y))$$
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$$= bx + (a - \lfloor \frac{x}{y} \rfloor \cdot b)y$$

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ext-gcd(x, y) calls ext-gcd(y, mod(x, y)) so

$$d = ay + b \cdot (\mod (x, y))$$

= $ay + b \cdot (x - \lfloor \frac{x}{y} \rfloor y)$
= $bx + (a - \lfloor \frac{x}{y} \rfloor \cdot b)y$

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Recursively: $d = ay + b(x - \lfloor \frac{x}{y} \rfloor \cdot y)$

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Recursively: $d = ay + b(x - \lfloor \frac{x}{y} \rfloor \cdot y) \implies d = bx - (a - \lfloor \frac{x}{y} \rfloor b)y$ Returns $(d, b, (a - \lfloor \frac{x}{y} \rfloor \cdot b))$.

Conclusion: Can find multiplicative inverses in O(n) time!

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Inverse of 500,000,357 modulo 1,000,000,000,000?

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Inverse of 500,000,357 modulo 1,000,000,000,000? \leq 80 divisions.

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Internet Security.

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Internet Security. Public Key Cryptography: 512 digits.

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Next Time.