Some Questions & Answers

- Q: Will there be reserved Lab times?
  A: We asked for these, but the scheduling folks are resisting.

- Q: Can we do the projects in groups?
  A: The first two projects are to be done individually. We haven’t yet decided about the third project.

- Q: What about the scheduling conflict with the finals for CS 162 and CS 194?
  A: The instructors are aware of these and are trying to work out a general solution. If need be, we will have an alternative time for those students.
Goals for Today’s Class

- Type of Networks
  - And the key concept of **multiplexing**
- What’s a **Protocol**?
- Clients & Servers & Peer-to-Peer (time permitting)

Taxonomy of Communication Networks

- Communication networks can be classified based on the way in which the nodes exchange information:
  
  Communication Network
Communication networks can be classified based on the way in which the nodes exchange information:

Information transmitted by any node is received by every other node in the network
- Examples?
  - Usually in LANs (*Local Area Networks*)
    - E.g., Ethernet, WiFi
    - E.g., lecture!

What problems does this raise?
Problem #1: limited range.
Problem #2: coordinating access to the shared communication medium (*Multiple Access Problem*)
Communication networks can be classified based on the way in which the nodes exchange information:

- Switched Communication Network
- Broadcast Communication Network

Switched Communication Networks

- Information transmitted along a path of intermediary nodes (“switches” or “routers”)
- Basic issue: how the switches figure out the next hop along the path

- Example and properties of a global switched communication network that you and millions of others use every day?
- Another example?
Communication networks can be classified based on the way in which the nodes exchange information:

- **Switched Communication Network**
- **Broadcast Communication Network**
- **Circuit-Switched Communication Network**

**Circuit Switching (e.g., Phone Network)**

- **Establish**: source creates circuit to destination
  - Node along the path store connection info
  - Nodes generally reserve resources for the connection
  - If circuit not available: “Busy signal”

- **Transfer**: source sends data over the circuit
  - No destination address, since nodes know path

- **Teardown**: source tears down circuit when done
Circuit Switching

- Node (switch) in a circuit switching network

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Telephone Network

- Alexander Graham Bell
  - 1876: Demonstrates the telephone at US Centenary Exhibition in Philadelphia
Circuit Switching With Human Operator

Telephone Network

• Almon Brown Strowger (1839 - 1902)
  – 1889: Invents the "girl-less, cuss-less" telephone system
    -- the mechanical switching system
Timing in Circuit Switching

Circuit Establishment

Transmission delay

propagation delay between Host 1 and Switch1

time
Timing in Circuit Switching

Circuit Establishment

Transmission delay

propagation delay between Host 1 and Switch 1

propagation delay between Host 1 and Host 2

Transfer

Information

time
Timing in Circuit Switching

Circuit Switching

- Node (switch) in a circuit switching network

How do the black and orange circuits share the outgoing link?
Circuit Switching: *Multiplexing a Link*

- **Time-division**
  - Each circuit allocated certain time slots

- **Frequency-division**
  - Each circuit allocated certain frequencies

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**Time-Division Multiplexing/Demultiplexing**

- Time divided into frames; frames into slots
- Relative slot position inside a frame *determines* to which conversation data belongs
  - E.g., slot 0 belongs to orange conversation
- Requires synchronization between sender and receiver—surprisingly difficult!
- In case of non-permanent conversations
  - Need to dynamically bind a slot to a conversation
  - How to do this?
- If a conversation does not use its circuit *the capacity is lost!*
Communication networks can be classified based on the way in which the nodes exchange information:

- **Packet Switching**
  - Data sent as chunks of formatted bit-sequences (Packets)
  - Packets have following structure:
    - Header and Trailer carry control information (e.g., destination address, checksum)
    - Each packet traverses the network from node to node along some path (Routing)
    - Once a node receives the entire packet, it stores it (briefly) and then forwards it to the next node (Store-and-Forward Networks)
    - Typically no capacity is allocated for packets
Packet Switching

- Node in a packet switching network

Box diagram showing incoming and outgoing links, and memory.

Packet Switching: Multiplexing/Demultiplexing

- Data from any conversation can be transmitted at any given time
  - Single conversation can use the entire link capacity if it is alone
- How to tell them apart?
  - Use meta-data (header) to describe data
5 Minute Break

Questions Before We Proceed?

Taxonomy of Communication Networks

- Communication networks can be classified based on the way in which the nodes exchange information:

  - Communication Network
    - Switched Communication Network
    - Circuit-Switched Communication Network
    - Packet-Switched Communication Network
    - Broadcast Communication Network
    - Datagram Network
Datagram Packet Switching

- Each packet is **independently switched**
  - Each packet header contains full destination address
- No resources are pre-allocated (reserved) in advance
- Leverages “statistical multiplexing” (or *stat-muxing*)
  - Essentially: “chances are good that packets from different conversations won’t all arrive at the same time, so we can get by without enough capacity for all of them at their peak transmission rate”
- Example: IP networks; postal system

Timing of Datagram Packet Switching

- Diagram showing propagation delay between Host 1 and Node 1
- Conceptual representation of packet delivery and node processing in a network scenario.
Timing of Datagram Packet Switching

Transmission time of Packet 1 at Host 1

Propagation delay between Host 1 and Node 1

Processing delay of Packet 1 at Node 2

Packet 1

Transmission time of Packet 1 at Host 1

Propagation delay between Host 1 and Node 1

Packet 1

Packet 1
Timing of Datagram Packet Switching

Datagram Packet Switching
Communication networks can be classified based on the way in which the nodes exchange information:

- **Switched Communication Network**
- **Packet-Switched Communication Network**
- **Dataagram Network**
- **Virtual Circuit Network**

A hybrid of circuits and packets; we will visit these later for “quality of service”

**Advantages of Circuit Switching**

- **Guaranteed bandwidth**
  - Predictable communication performance
  - Not “best-effort” delivery with no real guarantees

- **Simple abstraction**
  - Reliable communication channel between hosts
  - No worries about lost or out-of-order packets

- **Simple forwarding**
  - Forwarding based on time slot or frequency
  - No need to inspect a packet header

- **Low per-packet overhead**
  - Forwarding based on time slot or frequency
  - No IP (and TCP/UDP) header on each packet
Disadvantages of Circuit Switching

- **Wasted bandwidth**
  - Bursty traffic leads to idle connection during silent period
  - Unable to achieve gains from “statistical multiplexing”

- **Blocked connections**
  - Connection refused when resources are not sufficient
  - Unable to offer “okay” service to everybody

- **Connection set-up delay**
  - No communication until the connection is set up
  - Unable to avoid extra latency for small data transfers

- **Network state**
  - Network nodes must store per-connection information
  - Unable to avoid per-connection storage and state
  - This makes failures more disruptive!

Packet-Switching vs. Circuit-Switching

- **Critical advantage of packet-switching over circuit switching:** *Exploitation of statistical multiplexing*

- **Another:** since routers don’t know about individual conversations, when a router or link fails, it’s **easy to fail over to a different path**

- A third: easier for different parties to link their networks together because they’re **not** promising to reserve resources for one another

- **However, packet-switching must handle congestion:**
  - More complex routers
  - Harder to provide good network services (e.g., delay and bandwidth guarantees)

- In practice, sometimes combined, e.g., **IP over SONET**
Questions?

What Is A Protocol?

• A protocol is an agreement on how to communicate
• Includes syntax and semantics
  – How a communication is specified & structured
  – What a communication means
Examples of Protocols in Human Interactions

• Telephone
  1. Dialing.
  2. Ringing …
  3. Callee: “Hello?”
  4. Caller: “Hi, it’s Alice …”
     Or: “Hi, it’s me” (← what’s that about?)
  5. Caller: “Hey, do you think … blah blah blah …” pause
  6. Callee: “Yeah, blah blah blah …” pause

Examples of Protocols in Human Interactions

• Asking a question
  1. Raise your hand.
  2. Wait to be called on.
  3. Or: wait for speaker to pause and vocalize
Examples of Protocols in Human Interactions

• Expressing approval
  1. Applause.

    Who is this woman?

    Why don’t her fans applaud her performances?

  2. Foot-stomping

  3. The key: a protocol should effectively enable communication.

Example: HyperText Transfer Protocol

Request

GET /courses/archive/spring06/cos461/ HTTP/1.1
Host: www.cs.princeton.edu
User-Agent: Mozilla/4.03

Response

HTTP/1.1 200 OK
Date: Mon, 6 Feb 2006 13:09:03 GMT
Server: Netscape-Enterprise/3.5.1
Last-Modified: Mon, 6 Feb 2006 11:12:23 GMT
Content-Length: 21

Site under construction
Example: IP Packet

<table>
<thead>
<tr>
<th>4-bit Version</th>
<th>4-bit Header Length</th>
<th>8-bit Type of Service (TOS)</th>
<th>16-bit Total Length (Bytes)</th>
</tr>
</thead>
<tbody>
<tr>
<td>16-bit Identification</td>
<td>3-bit Flags</td>
<td>13-bit Fragment Offset</td>
<td></td>
</tr>
<tr>
<td>8-bit Time to Live (TTL)</td>
<td>8-bit Protocol</td>
<td>16-bit Header Checksum</td>
<td></td>
</tr>
</tbody>
</table>

32-bit Source IP Address

32-bit Destination IP Address

Options (if any)

Payload

IP: Best-Effort Packet Delivery

- **Packet switching**
  - Send data in packets
  - Header with source & destination address

- **“Best effort” delivery**
  - Packets may be lost
  - Packets may be corrupted
  - Packets may be delivered out of order
Example: Transmission Control Protocol

- Communication service
  - Ordered, reliable byte stream
  - Simultaneous transmission in both directions

- Key mechanisms at end hosts
  - Retransmit lost and corrupted packets
  - Discard duplicate packets and put packets in order
  - Flow control to avoid overloading the receiver buffer
  - Congestion control to adapt sending rate to network load

Protocol Standardization

- Ensure communicating hosts speak the same protocol
  - Standardization to enable multiple implementations
  - Or, the same folks have to write all the software

- Standardization: Internet Engineering Task Force
  - Based on working groups that focus on specific issues
  - Produces “Request For Comments” (RFCs)
    - Promoted to standards via rough consensus and running code
  - IETF Web site is [http://www.ietf.org](http://www.ietf.org)
  - RFCs archived at [http://www.rfc-editor.org](http://www.rfc-editor.org) (per Homework #1)

- De facto standards: same folks writing the code
  - P2P file sharing, Skype, <your protocol here>…
End System: Computer on the ‘Net

Also known as a “host”...

Clients and Servers

- Client program
  - Running on end host
  - Requests service
  - E.g., Web browser

GET /index.html
Clients and Servers

- **Client program**
  - Running on end host
  - Requests service
  - E.g., Web browser

- **Server program**
  - Running on end host
  - Provides service
  - E.g., Web server

```
GET /index.html
```

“Site under construction”

Clients Are Not Necessarily Human

- **Example: Web crawler (or spider)**
  - Automated client program
  - Tries to discover & download many Web pages
  - Forms the basis of search engines like Google

- **Spider client**
  - Start with a base list of popular Web sites
  - Download the Web pages
  - Parse the HTML files to extract hypertext links
  - Download these Web pages, too
  - And repeat, and repeat, and repeat...
  - (Per Project #2)
Client-Server Communication

- Client "sometimes on"
  - Initiates a request to the server when interested
  - E.g., Web browser on your laptop or cell phone
  - Doesn’t communicate directly with other clients
  - Needs to know the server’s address

- Server is “always on”
  - Services requests from many client hosts
  - E.g., Web server for the www.cnn.com Web site
  - Doesn’t initiate contact with the clients
  - Needs a fixed, well-known address

Peer-to-Peer Communication

- No always-on server at the center of it all
  - Hosts can come and go, and change addresses
  - Hosts may have a different address each time

- Example: peer-to-peer file sharing
  - Any host can request files, send files, query to find where a file is located, respond to queries, and forward queries
  - Scalability by harnessing millions of peers
  - Each peer acting as both a client and server
Client and Server Processes

• Program vs. process
  – Program: collection of code
  – Process: a running program on a host

• Communication between processes
  – Same end host: inter-process communication
    • Governed by the operating system on the end host
  – Different end hosts: exchanging messages
    • Governed by the network protocols

• Client and server processes
  – Client process: process that initiates communication
  – Server process: process that waits to be contacted

Summary

• Types of communication networks
  – Broadcast, circuit-switched, packet-switched
  – Advantages & disadvantages of each

• Protocols: agreement on how to communicate
  – For networks, often requires formal standardization

• Roles of endpoints in communication:
  – Clients (initiate communication)
  – Servers (receive requests)
  – Peer-to-peer (hosts play both roles, linked to one another)

• Next lecture: “Architecture” and the “End-to-End Principle”
  – Note: on Wednesday Sept. 6 - Monday is a holiday!
  – Read through 1.3 of the Peterson/Davie book
  – Take the survey (http://tinyurl.com/fbc7u) ASAP and join the mailing list if you haven’t already