

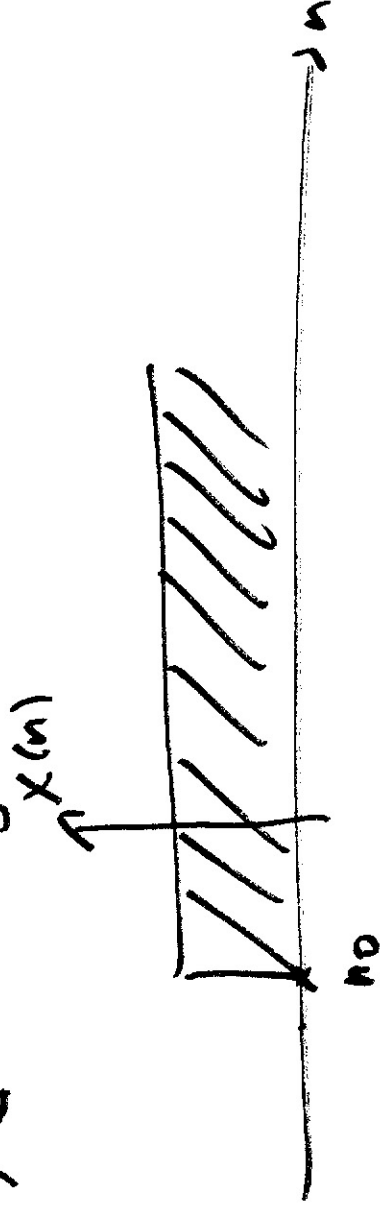
# How to choose B.C. for 2D L.C.C. D.E

Feb 10<sup>1</sup>/<sub>06</sub>

Approach: Generalize the approach 1.D.

causal

$$y(n) \leftarrow y(n-1) + x(n)$$



To obtain LTI system  $y(n_0-1) = 0$

① Compute Region of support of output given the D.E.  $\rightarrow R_y$

② set B.C. = 0 outside of  $R_y$ .

$\psi$

## Detailed

(1) ROS of input  $R_x$

(2) Given Comp. provides for  $D E$

→ Know impulse response  $h(n)$

→ Compute  $R_h =$  Region of support for impulse response.

(3) Use  $R_x$  and  $R_h$  to compute

$R_y =$  Region of support of output

$H_n \notin R_y$

(4) B.C. = 0

Ex  $y(n) \leftarrow y(n-1) + x(n)$  *causal.*

$$x(n) = \delta(n)$$

$R_y =$  what is ROS for output

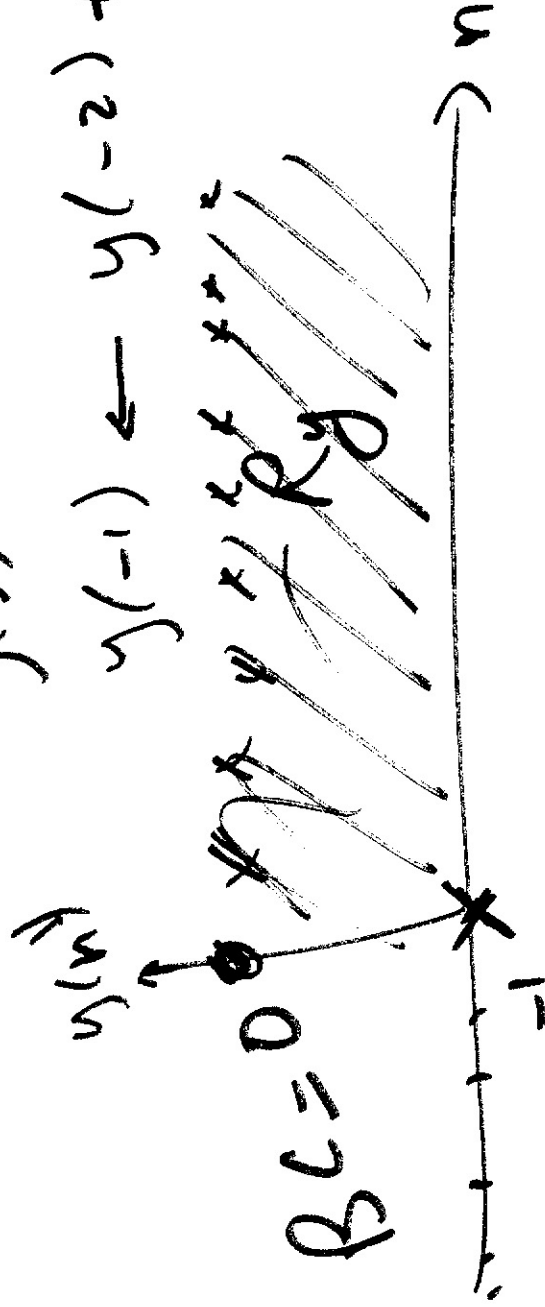
$$\begin{array}{c} \xrightarrow{y(0)} \leftarrow y(0) \leftarrow y(-1) + \delta(0) \end{array}$$

$$y(1) \leftarrow y(0) + 0$$

$$y(2) \leftarrow y(1)$$

$$y(3) \leftarrow y(2)$$

$$y(-1) \leftarrow y(-2) + 0$$

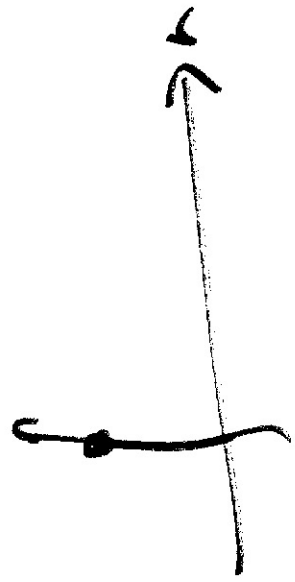


Aliter  $h(n)$ ?

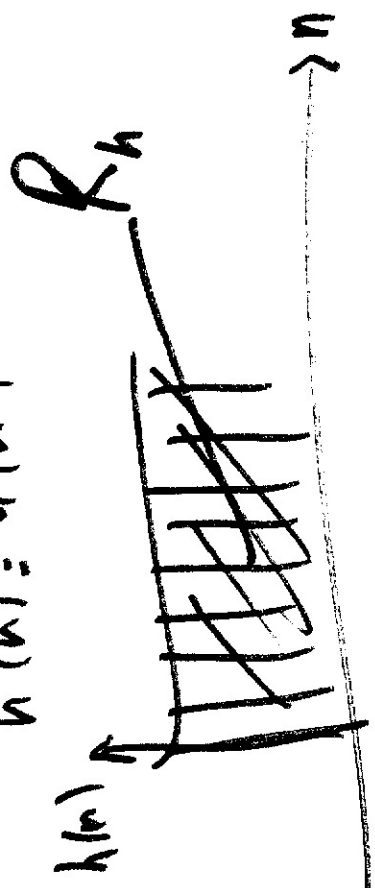
$$Y(z) = z^{-1} Y(z) + X(z)$$

$$H(z) = \frac{Y(z)}{X(z)} = \frac{z^{-1} Y(z)}{Y(z) + X(z)}$$

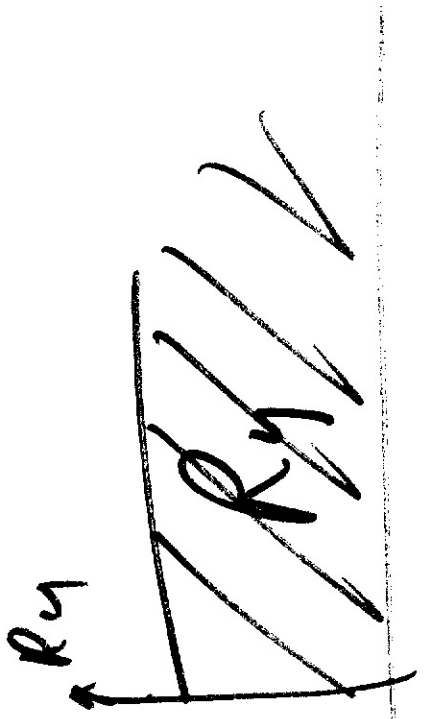
$f(n)$



$h(n) = u(n)$



$X(n) = \delta(n)$



BC=0

What happens in 2D?

Exact same steps

① Given Comp procedure  $\longrightarrow R_h$ .

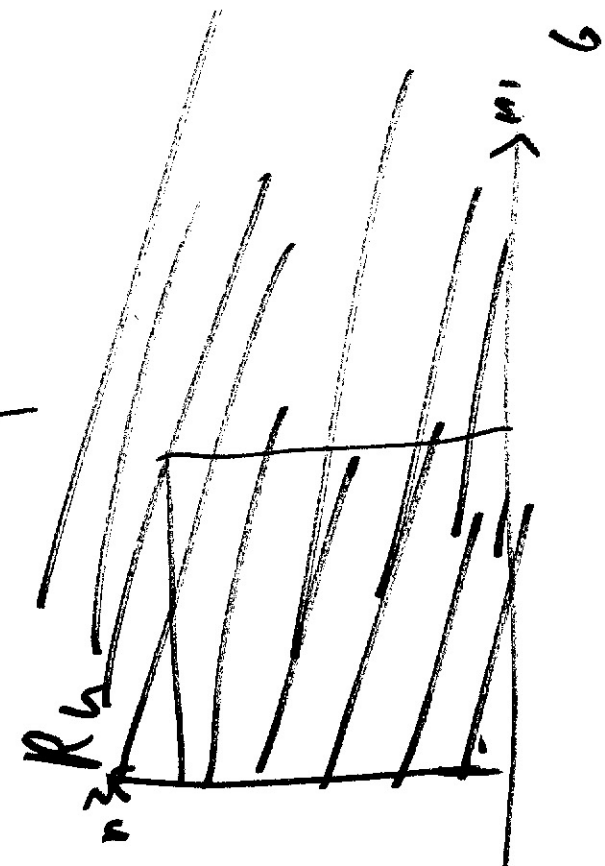
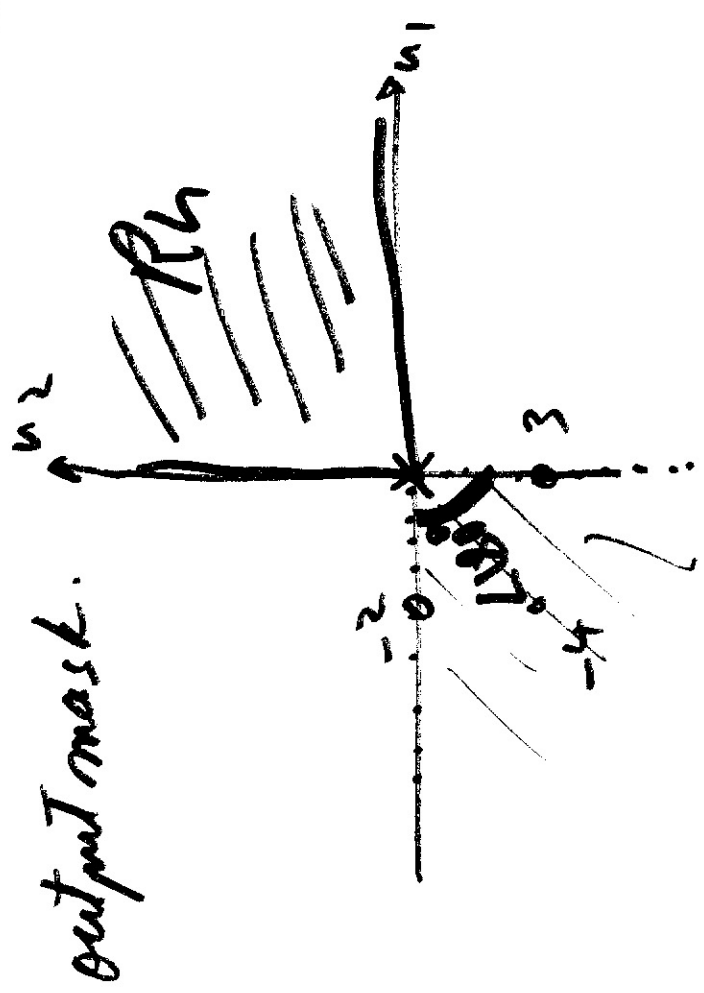
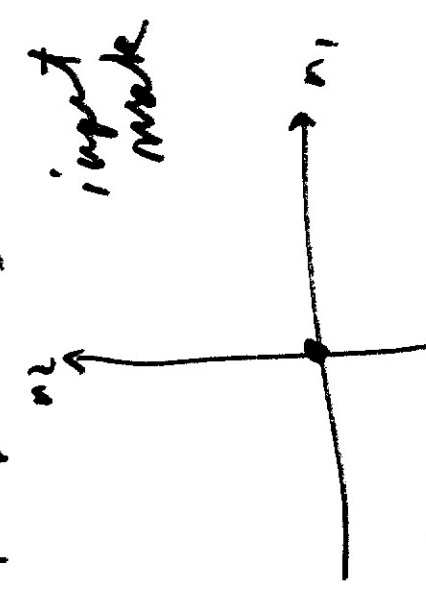
② Given input  $\longrightarrow R_x$

③  $R_x + R_h \longrightarrow R_y$

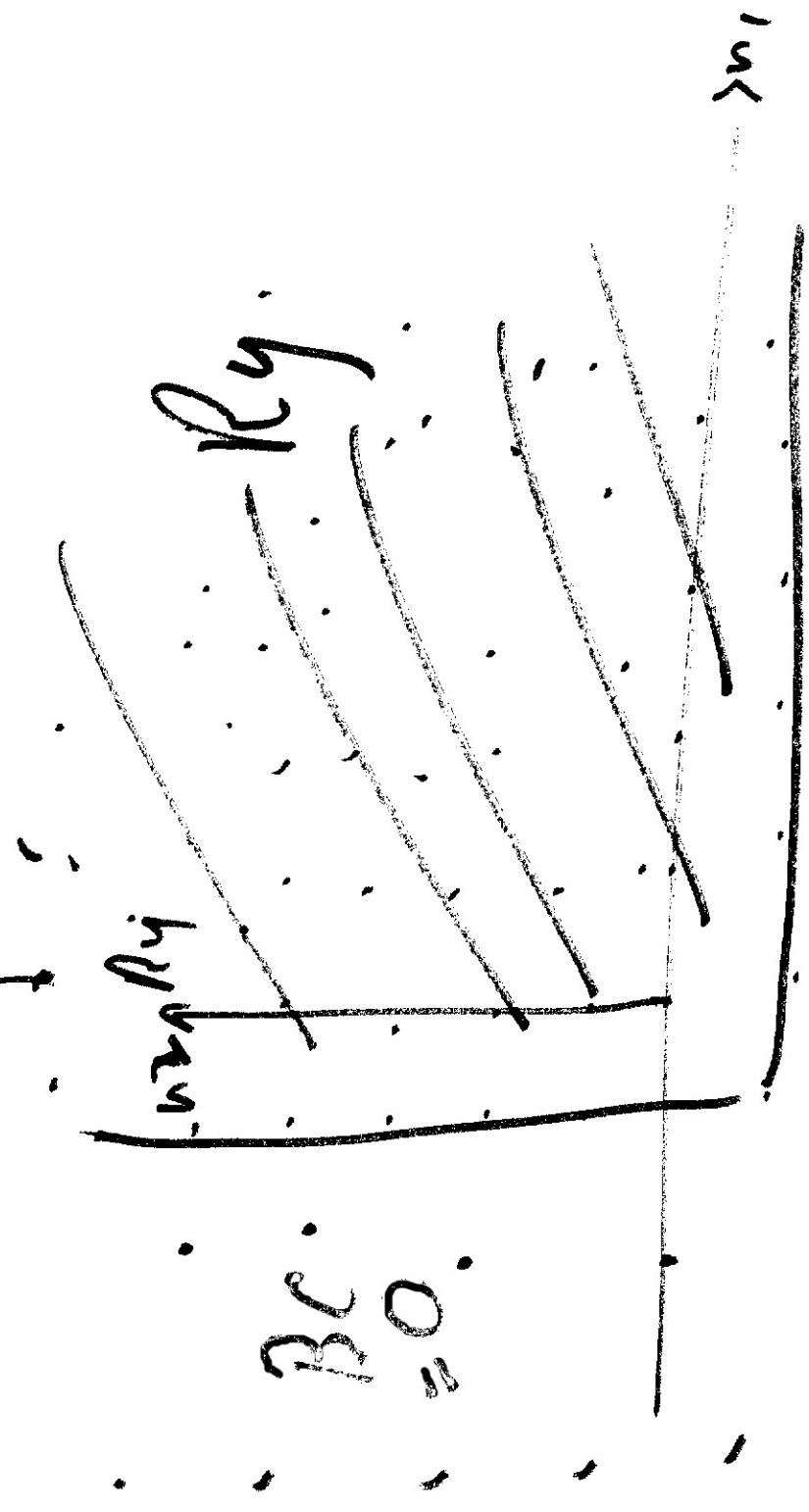
④ P.C = 0  $\quad \forall n_1, n_2 \neq R_y$ .

Ex ~~Comp.~~ Comp. Proc:

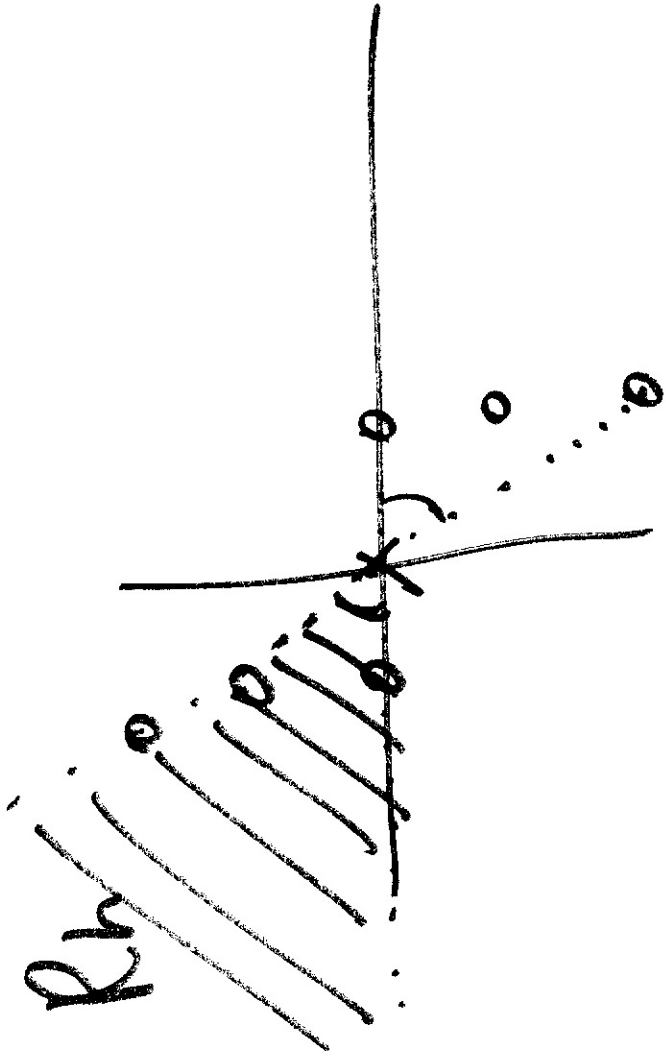
$$y(n_1, n_2) \leftarrow -2y(n_1-1, n_2) - 3y(n_1, n_2-1) - 4y(n_1-1, n_2-1) + x(n_1, n_2)$$



Input  $x(n_1, n_2)$



P. 98-100 J.S. Lina's book



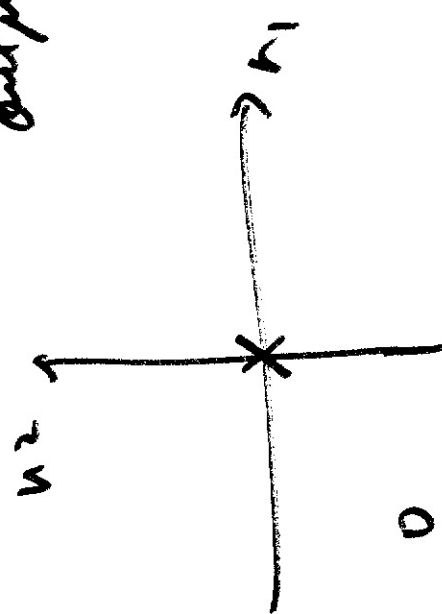


Ex  $y(n_1, n_2) = y(n_1-1, n_2-1) + x(n_1, n_2)$

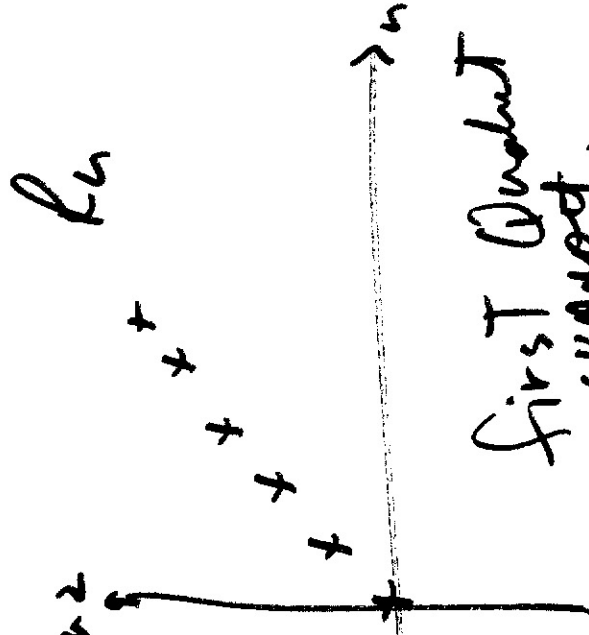
Find Comp. Proc  $\longrightarrow$   $h(n_1, n_2)$  is 3rd Quadrant.

Only 2 long. proc.

(1)  $y(n_1, n_2)$   $\longleftarrow$  output mask



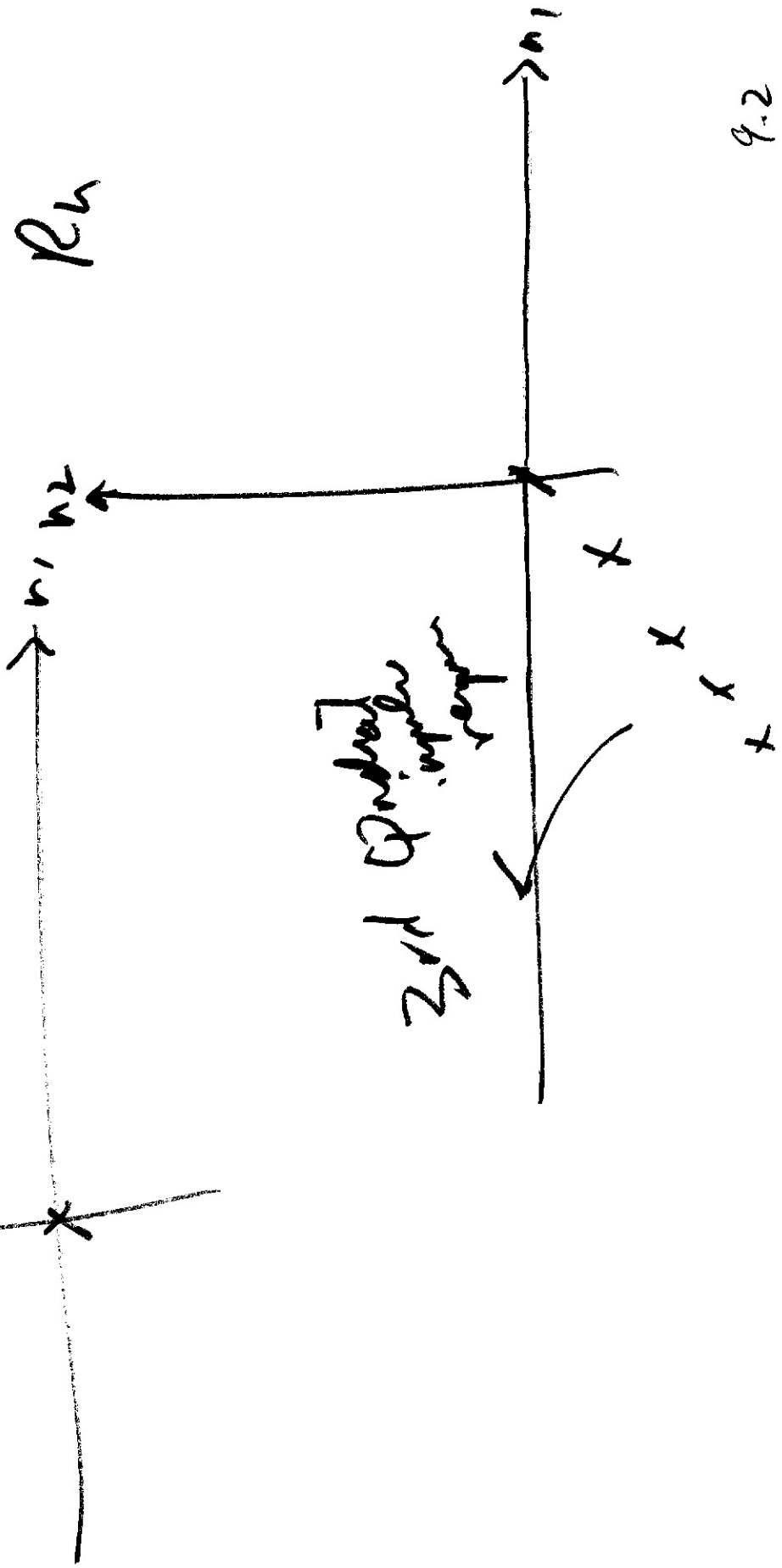
$y(n_1-1, n_2-1) + x(n_1, n_2)$

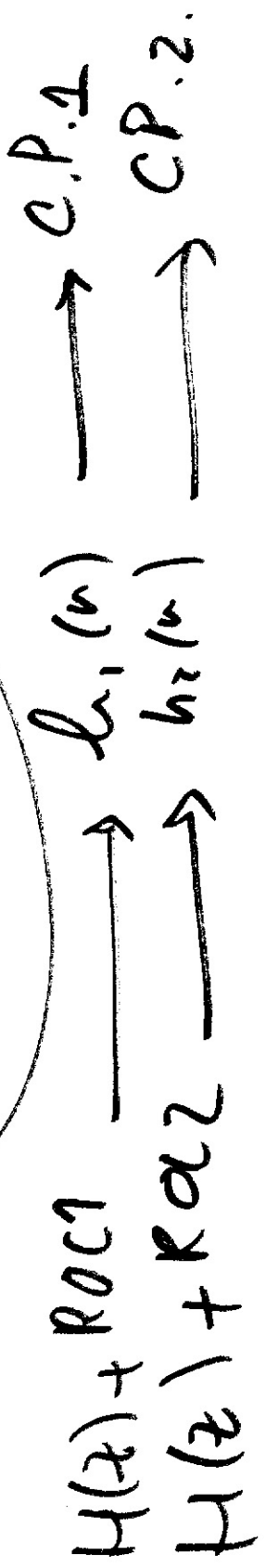
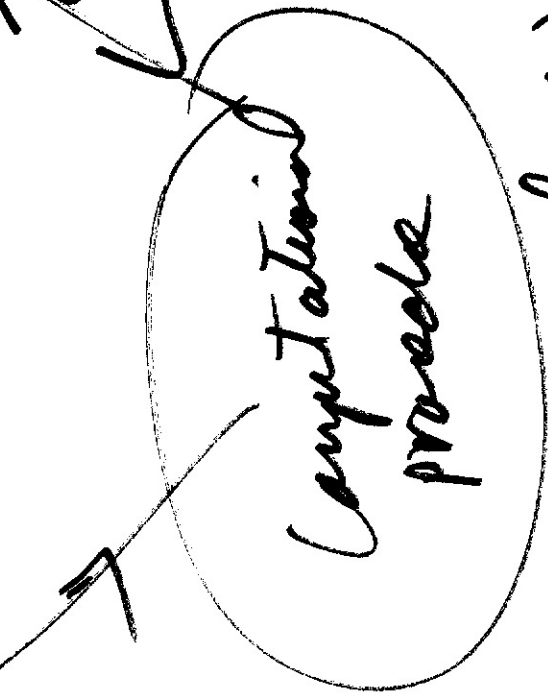
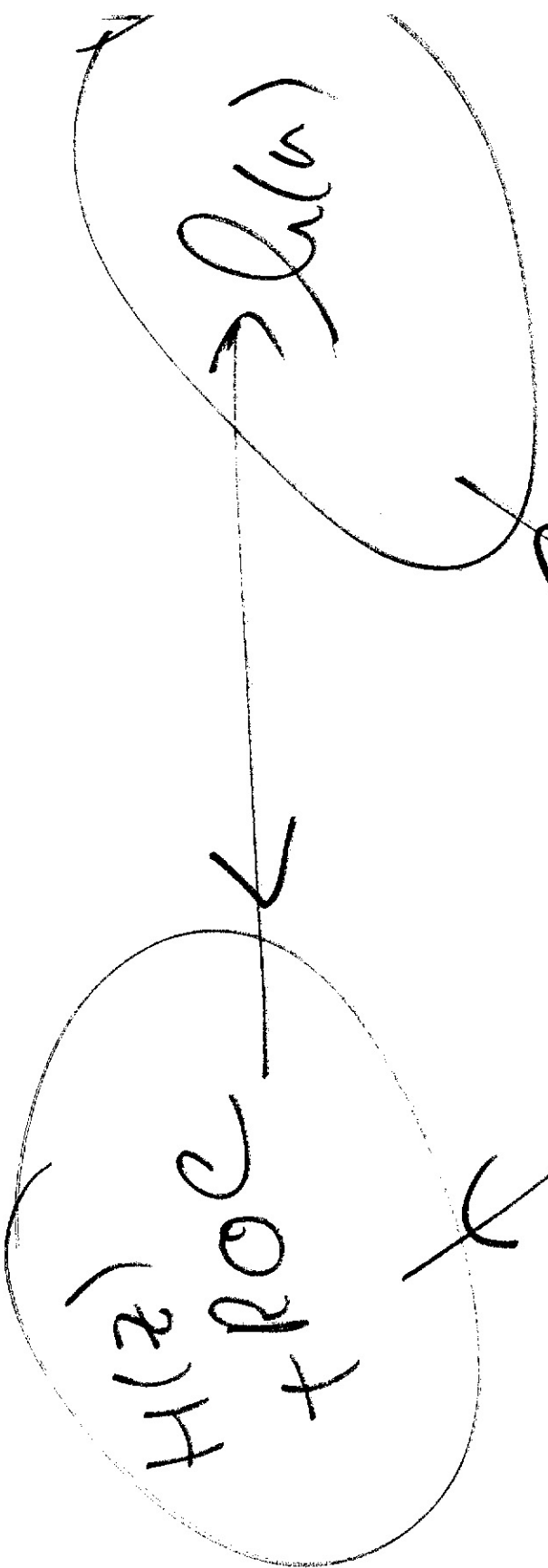


first Quadrant support.

$$y(n_1, n_2) \xrightarrow{\text{budget constraint}} y(n_1+1, n_2+1) - x(n_1+1, n_2+1)$$

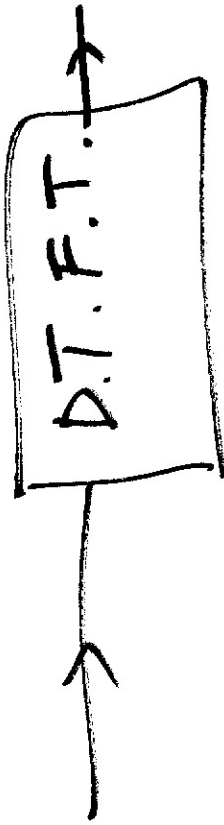
$n_2$  budget constraint





# 2.D. D.F.T

$x(n)$   
discrete



$$X(\omega) = \sum_n x(n) e^{-j\omega n}$$

continuous variable  
real.

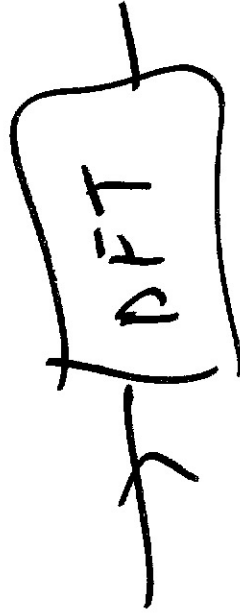
$x(n)$



$$X(z) = \sum_n x(n) z^{-n}$$

cont. variable  
complex.

$x(n)$



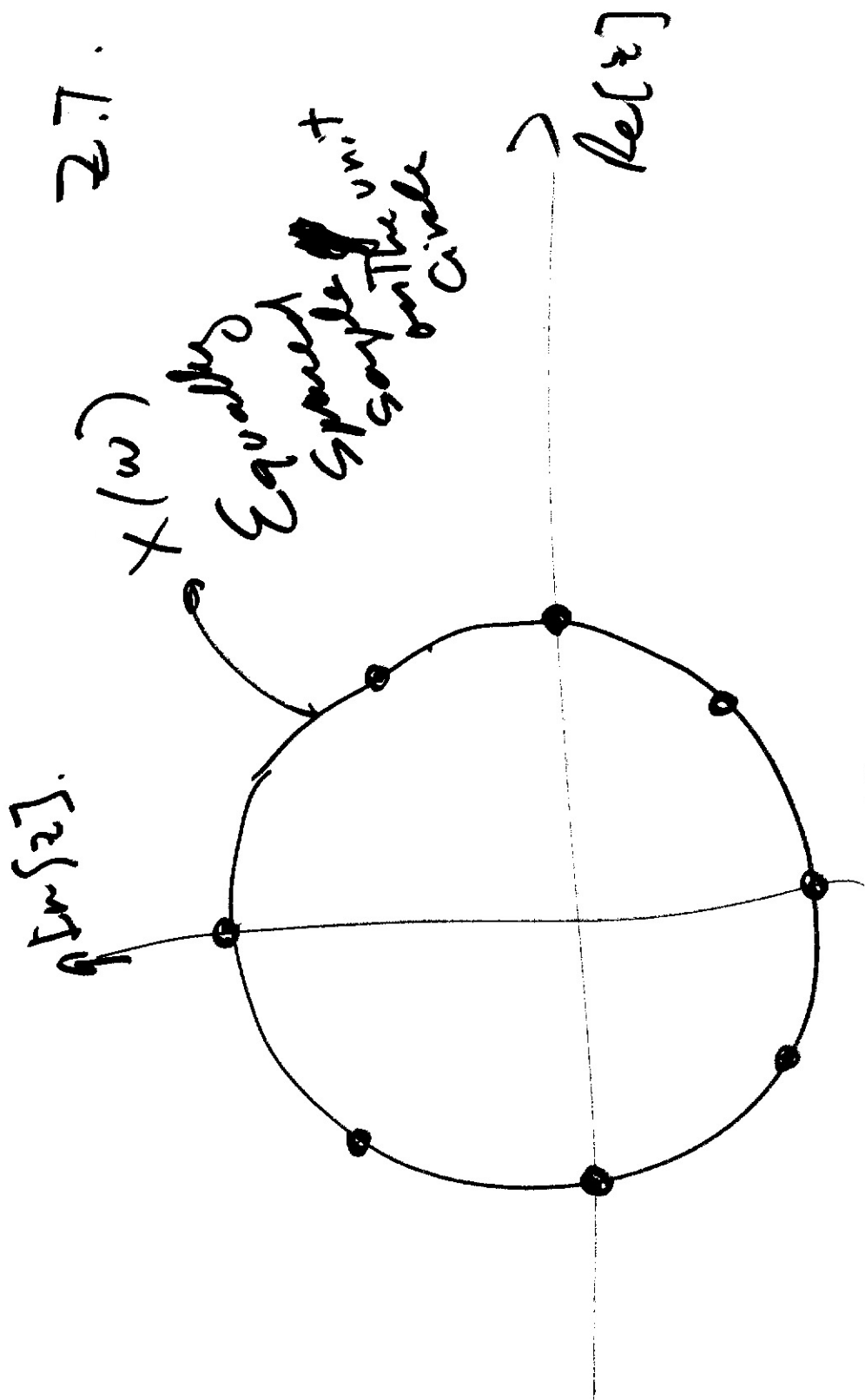
$$X(k) = \sum_n x(n) e^{-j2\pi nk/N}$$

integer  
discrete



= Fast Fourier Transform (1)

2.7.



$$X(k) = [X(w)]_{w = \frac{2\pi k}{N}}$$

$n_1, n_2$

$$\text{DFT} \{ x(n_1, n_2) \} = X(k_1, k_2)$$

2. D. DFT

$$\sum_{n_1=0}^{N_1-1} \sum_{n_2=0}^{N_2-1} x(n_1, n_2) e^{-j2\pi(n_1 k_1 / N_1 + n_2 k_2 / N_2)}$$

$0 \leq k_1 < N_1$   
 $0 \leq k_2 < N_2$

0

otherwise

$$\sum_{n_1=0}^{N_1-1} \sum_{n_2=0}^{N_2-1} X(k_1, k_2) e^{j2\pi(n_1 k_1 / N_1 + n_2 k_2 / N_2)}$$

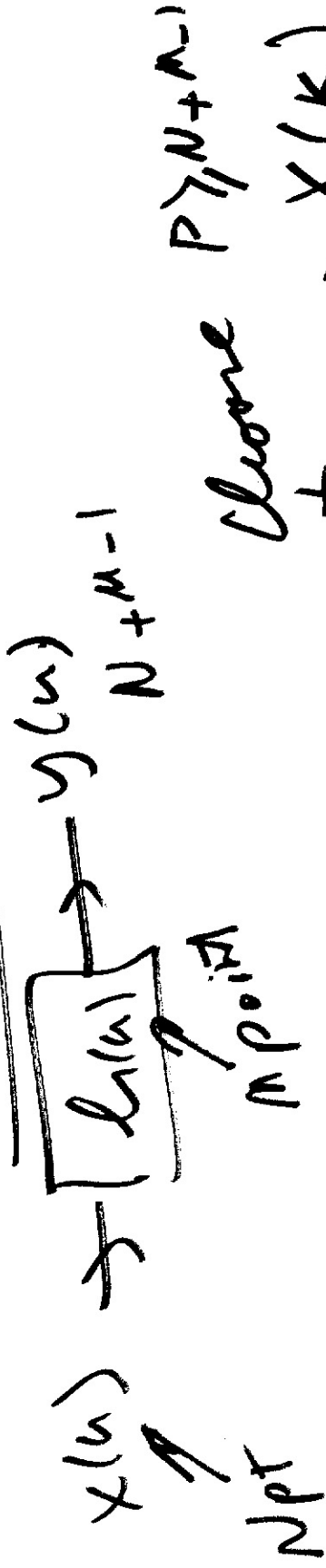
$0 \leq n_1 < N_1$   
 $0 \leq n_2 < N_2$

otherwise

0

$$x(n_1, n_2) = \text{IDFT} \{ X(k_1, k_2) \}$$

# Use of DFT in Convolution



- ① Compute P point DFT of input  $\rightarrow X(k)$
- ② " " " " " "  $\rightarrow H(k)$
- ③ " " " " " "  $\rightarrow Y(k) = X(k)H(k) = \text{DFT}\{y(n)\}$
- ④ " " " " " "  $\rightarrow y(n)$

All of this applies in 2D

## How To Compute 2D DFT

- 3 ways
- ① Direct method
  - ② Row/Col. decomposition
  - ③ Vector-Radix  $\leftarrow$  True FFT

## Factors affect speed

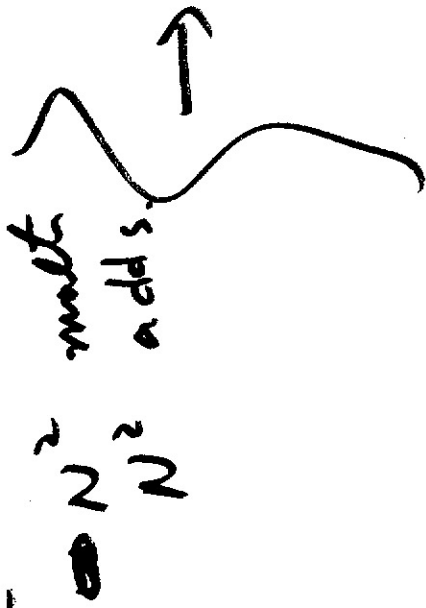
- # of arithmetic ops
- complexity of program
- memory for execution  $\rightarrow$  footprint
- memory usage on processor.
- language / platform.



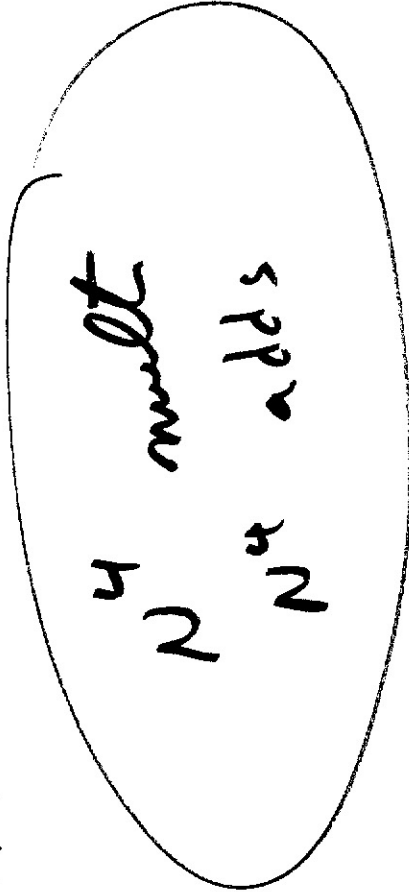
Direct Method

$N_1 = N_2 = N$

① For each  $(k_1, k_2)$   $\rightarrow$



② how many  $(k_1, k_2) ? N^2$ .



Row 1

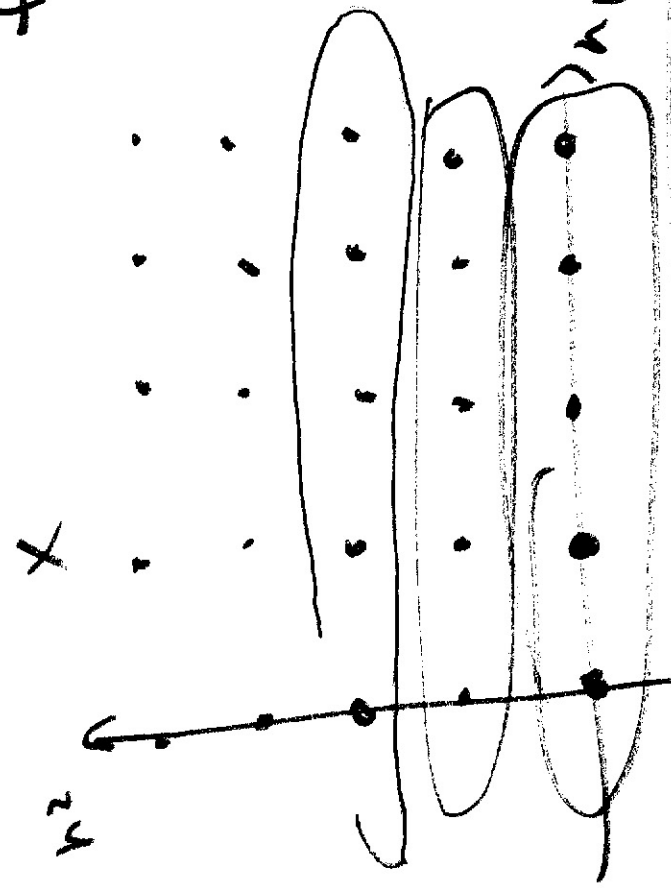
# Row Column Decomposition

$$X(k_1, k_2) = \sum_{n_2=0}^{N-1}$$

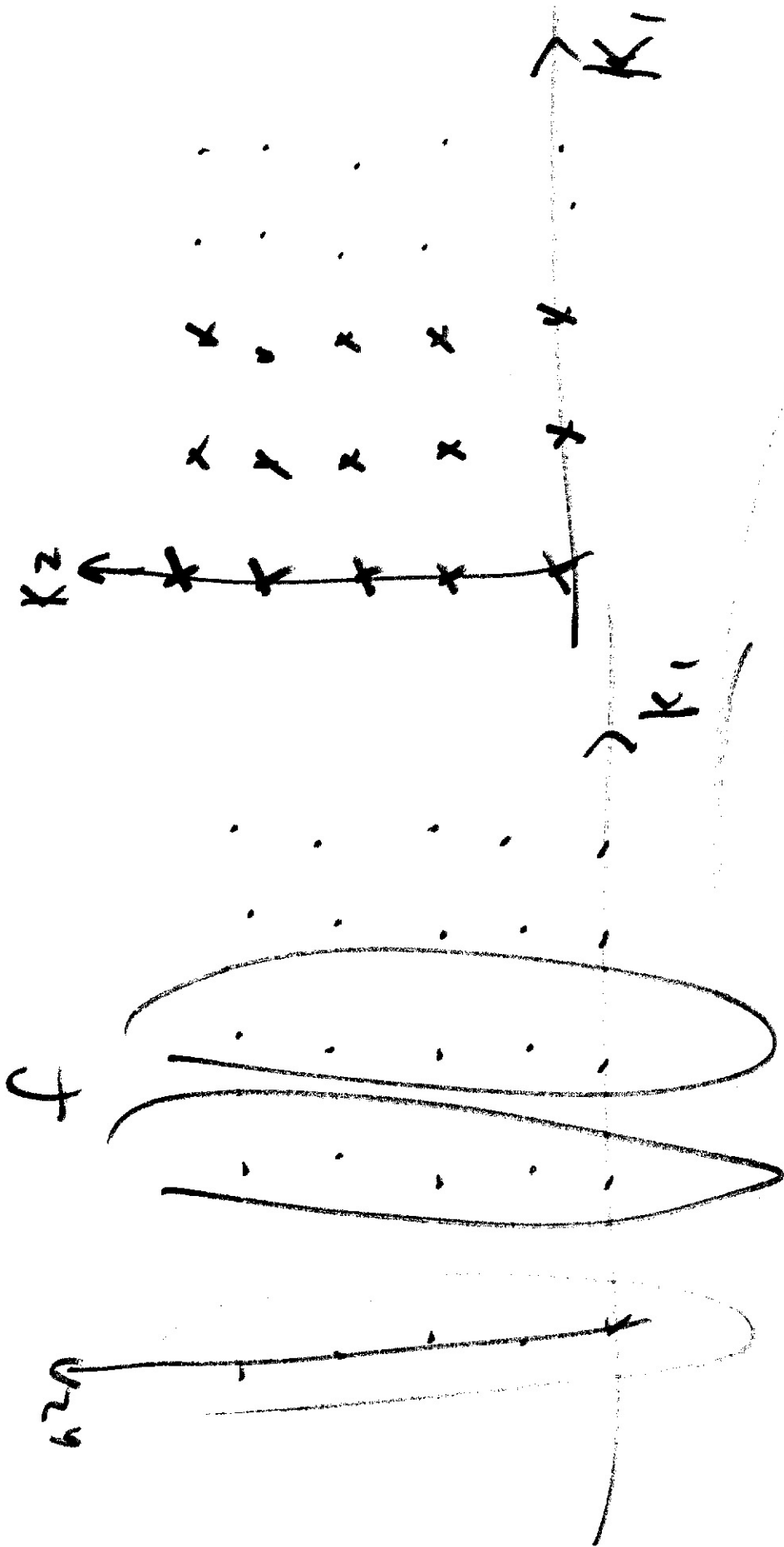
$$\left( \sum_{n_1=0}^{N-1} x(n_1, n_2) e^{-j2\pi n_1 k_1} \right) e^{-j2\pi n_2 k_2}$$

$$e^{-j2\pi n_2 k_2}$$

$f(k_1, n_2)$



$N$  - 1-D DFT of row  
each 1-D DFT is  $N$  point



$N - 1$  ID DFT.  $N$  PT  
 each ID DFT was NPT

1-D DFT  $\longrightarrow$  FFT

$\cdot N \log_2 N$  multi

$\frac{N}{2} \log_2 N$  adds

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R/C decomp  $N^2 \log_2 N$  operations

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lot lower direct method