

**EECS 151/251A**

**Discussion 8**

04/13/2018

# Announcements

- That extra discussion with Taehwan will be in two weeks
  - Location/time TBA, slides will be available if you can't make it.
- Homework 10 out by Sunday

# Agenda

- Memories: Adders
- Your questions

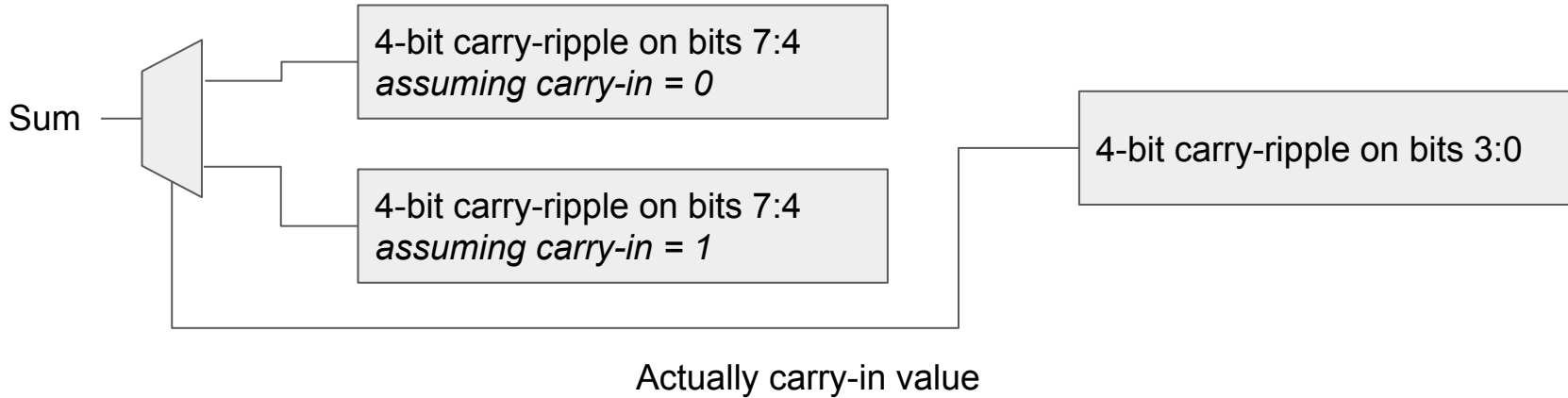
# Carry-ripple adder

Problem?

Slow!

# Carry-select adder

e.g. adding two 8-bit numbers



What do you need to consider when picking: size of group, number of groups?

What's best?  $\sqrt{N}$  stages of  $\sqrt{N}$  bits

## Better (from your notes)

- Compare to ripple adder delay:

$$T_{\text{total}} = 2 \sqrt{N} T_{\text{FA}} - T_{\text{FA}}, \text{ assuming } T_{\text{FA}} = T_{\text{MUX}}$$

$$\text{For ripple adder } T_{\text{total}} = N T_{\text{FA}}$$

“cross-over” at  $N=3$ , Carry select faster for any value of  $N>3$ .

- Is  $\sqrt{N}$  really the optimum?

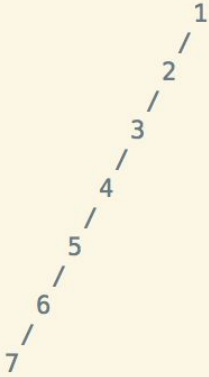
- From right to left increase size of each block to better match delays
- Ex: 64-bit adder, use block sizes [12 11 10 9 8 7 7], the exact answer depends on the relative delay of mux and FA

# Carry-select adder

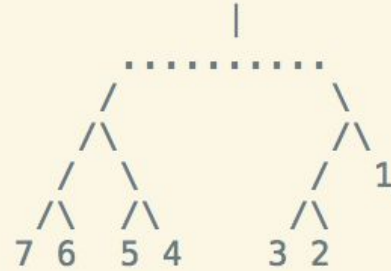
Problem?

*Groups still need to propagate addition in series.*

# Reduction trees



$$((((((7 + 6) + 5) + 4) + 3) + 2) + 1)$$



$$((7 + 6) + (5 + 4)) + ((3 + 2) + 1)$$

*6 time steps to compute*

*$\text{ceil}(\log_2(6)) = 3$  steps to compute*

Because? Associativity of addition.



# Carry-lookahead adders

- Try to turn addition into reduction operation
- Need associative computation for carry values

$a$	$b$	$c_i$	$c_{i+1}$	$s$	
0	0	0	0	0	carry "kill"
0	0	1	0	1	$k_i = a_i' b_i'$
<hr/>					
0	1	0	0	1	
0	1	1	1	0	carry "propagate"
1	0	0	0	1	$p_i = a_i \oplus b_i$
1	0	1	1	0	
<hr/>					
1	1	0	1	0	carry "generate"
1	1	1	1	1	$g_i = a_i b_i$

$$\begin{aligned} c_{i+1} &= g_i + p_i c_i \\ s_i &= p_i \oplus c_i \end{aligned}$$

(From your notes)

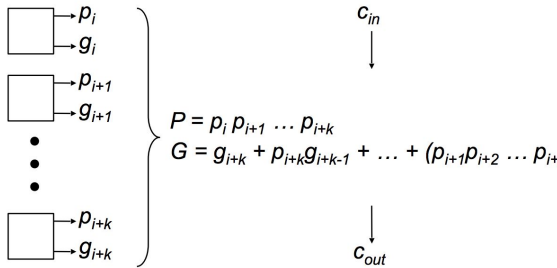
But no advantage if we do this for every bit as before...

# Carry-lookahead adders

## Carry Look-ahead Adders

### Carry Look-ahead Adders

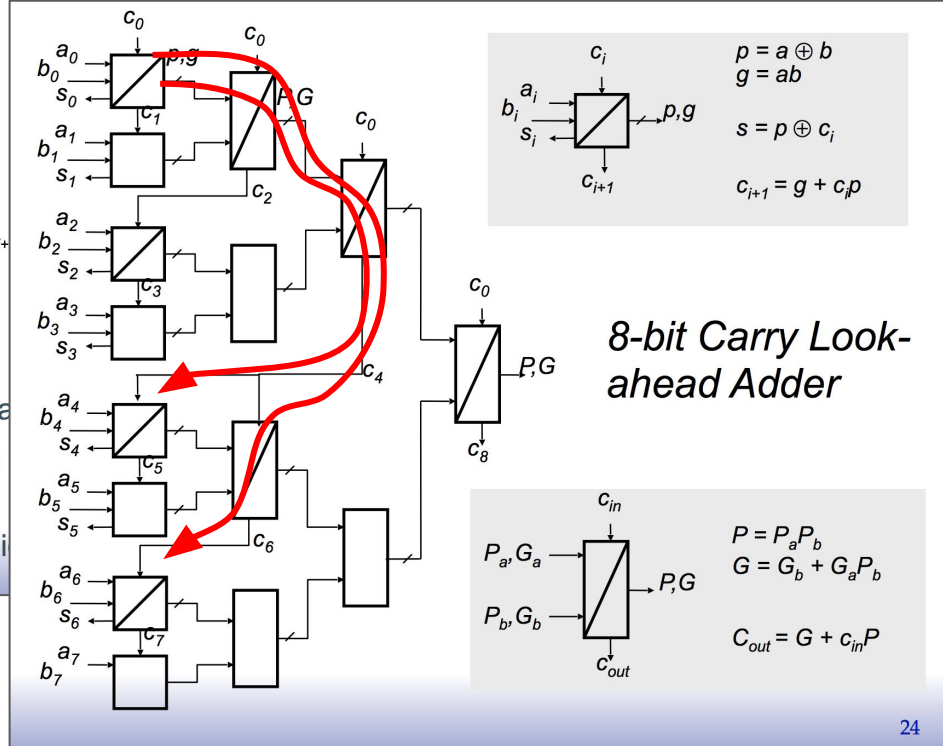
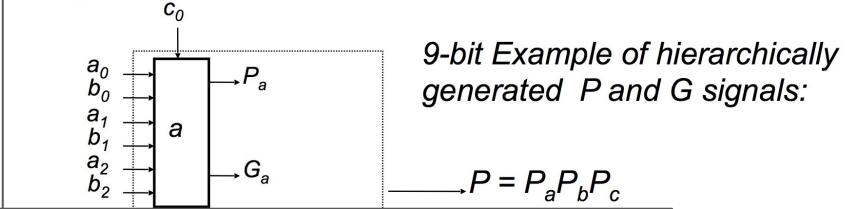
- “Group” propagate and generate signals:



- P true if the group as a whole propagates a carry
- G true if the group as a whole generates a carry

$$c_{out} = G + P c_{in}$$

- Group P and G can be generated hierarchically



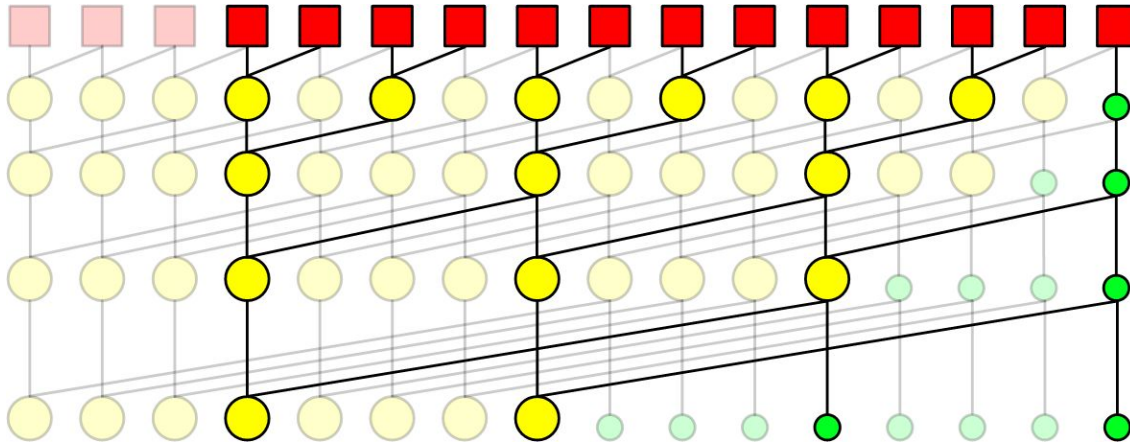
$$+ P_b P_c G_a$$

# Carry-lookahead

- Use combinational logic to calculate propagate (P) and generate (G) values for groups of bits
- Provide generated values to other groups well before their sum would ripple through
- Can combine with other techniques depending on specific design

# Parallel-prefix adders

- Generate carries for each bit directly; don't do any grouping there
- E.g. Kogge-Stone adder (radix 2, sparsity 4)



- (each line is a bit, each level is a successive computation)
- Each of these carries can feed a carry-select or carry-ripple (or other) adder (shown by transparency)

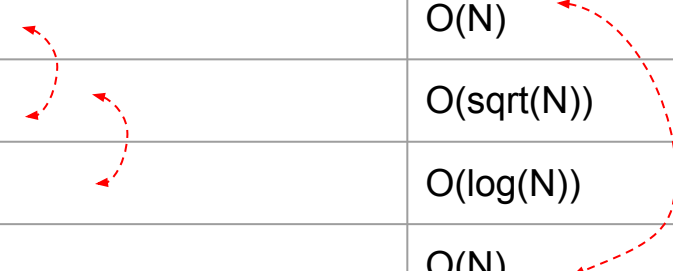
# Bit serial adders

- Draw one
- Pros?
  - Simple, require very few components, can feed other serial computations
- Cons?
  - Take  $\sim N$  cycles

# Adder summary

Where  $N$  is the operand width in bits:

Adder	Cost	Delay
Ripple	$O(N)$	$O(N)$
Carry-select	$O(N)$	$O(\sqrt{N})$
Carry-lookahead	$O(N)$	$O(\log(N))$
Bit-serial	$O(1)$	$O(N)$

The table contains red dashed arrows that highlight comparisons between the Cost and Delay columns. A curved arrow points from the  $O(N)$  cost of the Ripple adder to its  $O(N)$  delay. Another curved arrow points from the  $O(N)$  cost of the Carry-select adder to its  $O(\sqrt{N})$  delay. A third curved arrow points from the  $O(N)$  cost of the Carry-lookahead adder to its  $O(\log(N))$  delay. Finally, a long curved arrow points from the  $O(N)$  delay of the Bit-serial adder back to the  $O(N)$  delay of the Ripple adder.

Usual  $O$ -notation caveat: constants are hidden, but here may be important. **See notes!** *Lecture 20, slide 33.*

# References

- Lecture notes
- Wikipedia