UC Berkeley EECS
Sr Lecturer SOE
Dan Garcia
"KOOMEY'S LAW" - EFFICIENCY 2X EVERY 18 MO
Prof Jonathan Koomey looked at 6 decades of data and found that energy efficiency of computers doubles roughly every 18 months. This is even more relevant as battery-powered devices become more popular. Restated, it says that for a fixed computing load, the amount of battery you need drops by half
 every 18 months. This was true before transistors!

## Concurrency \& Parallelism, 10 mi up...

## Intra-computer

- Today's lecture
- Multiple computing "helpers" are cores within one machine
- Aka "multi-core"
- Although GPU parallism is also "intra-computer"


## Inter-computer

- Week 12's lectures
- Multiple computing "helpers" are different machines
- Aka "distributed computing"
- Grid \& cluster computing



## Anatomy: 5 components of any Computer



John von Neumann invented this architecture


## Computer

a) Control
b) Datapath

| Memory | Devices <br> Input <br>  <br> Output |
| :---: | :---: |

c) Memory
d) Input
e) Output

What causes the most headaches for SW and HW designers with multi-core computing?

## But what is INSIDE a Processor?



Bare Processor Die


Chip in Package

- Primarily Crystalline Silicon
- 1 mm - 25 mm on a side
- 2009 "feature size" (aka process) $\sim 45 \mathrm{~nm}=45 \times 10^{-9} \mathrm{~m}$ (then 32, 22, and 16 [by yr 2013])
- 100-1000M transistors
- 3-10 conductive layers
- "CMOS" (complementary metal oxide semiconductor) - most common
- Package provides:
- spreading of chip-level signal paths to board-level
- heat dissipation.
- Ceramic or plastic with gold wires.


## Moore's Law

Predicts: $\mathbf{2 X}$ Transistors / chip every 2 years



Gordon Moore Intel Cofounder
B.S. Cal 1950!

## Moore's Law and related curves



## Moore's Law and related curves



## Power Density Prediction circa 2000



## Going Mulii-core Helps Energy Efiiciency

- Power of typical integrated circuit ~ C V ${ }^{2} f$
- C = Capacitance, how well it "stores" a charge
- V = Voltage
- $f=$ frequency. l.e., how fast clock is (e.g., 3 GHz )

| In the same process technology... |  |  |
| :--- | :--- | :--- |
| Cache |  |  |
| Core |  | Cache |



Activity Monitor (on the lab Macs) shows how active
your cores are

## Energy \& Power Considerations



$$
\begin{aligned}
& \text { Power }=\frac{\text { Energy }}{\text { Second }}=\frac{\text { Energy }}{O p} \times \frac{\text { Ops }}{\text { Second }} \\
& \text { Power } \\
& \text { Chip Packaging } \\
& \text { Chip Cooling } \\
& \text { System Noise } \\
& \text { Case Temperature } \\
& \text { Data-Center Air } \\
& \text { Conditioning } \\
& \text { Energy } \\
& \text { Battery Life } \\
& \text { Electricity Bill } \\
& \text { Mobile Device } \\
& \text { Weight }
\end{aligned}
$$

## view. eecs.berkeley.edu <br> Parallelism again? What's different this time?

"This shift toward increasing parallelism is not a triumphant stride forward based on breakthroughs in novel soffware and archifectures for parallelism; instead, this plunge into parallelism is actually a retreat from even greater challenges that thwart efficient silicon implementation of traditional uniprocessor architectures."

- Berkeley View, December 2006
- HW/SW Industry bet its future that breakthroughs will appear before it's too late


## Background: Threads

" A Thread stands for "thread of execution", is a single stream of instructions

- A program / process can split, or fork itself into separate threads, which can (in theory) execute simultaneously.
- An easy way to describe/think about parallelism
- A single CPU can execute many threads by Time Division Multipexing
$\square$
Time

$\square$ Thread $_{0}$
$\square$ Thread, $_{1}$
$\square$ Thread $_{2}$
- Multithreading is running multiple threads through the same hardware


## Speedup Issues : Amdahl's Law

- Applications can almost never be completely parallelized; some serial code remains

- s is serial fraction of program, P is \# of cores (was processors)
- Amdahl's law:

Speedup(P) = Time(1) / Time(P)

$$
\begin{aligned}
& \leq 1 /(s+[(1-s) / P)] \text {, and as } P \rightarrow \infty \\
& \leq 1 / s
\end{aligned}
$$

- Even if the parallel portion of your application speeds up perfectly, your performance may be limited by the sequential portion


## Speedup lssues : Overhead

- Even assuming no sequential portion, there's...
- Time to think how to divide the problem up
- Time to hand out small "work units" to workers
- All workers may not work equally fast
- Some workers may fail
- There may be contention for shared resources
- Workers could overwriting each others' answers
- You may have to wait until the last worker returns to proceed (the slowest / weakest link problem)
- There's time to put the data back together in a way that looks as if it were done by one


## Life in a mulii-core wordd...

- This "sea change" to multicore parallelism means that the computing community has to rethink:
a) Languages
b) Architectures
c) Algorithms
d) Data Structures
e) All of the above



## But parallel programming is hard!

- What if two people were calling withdraw at the same time?
- Egg., balance=100 and two withdraw 75 each
- Can anyone see what the problem could be?
- This is a race condition
- In most languages, this is a problem.
- In Scratch, the system doesn't let two of these run at once.

YY Ne SA

## Another concurrency problem ... deadlock!

- Two people need to draw a graph but there is only one pencil and one ruler.
- One grabs the pencil
- One grabs the ruler
- Neither release what they hold, waiting for the other to release
- Livelock also possible
- Movement, no progress
- Dan and Luke demo


## Summary

- "Sea change" of computing because of inability to cool CPUs means we're now in multi-core world
- This brave new world offers lots of potential for innovation by computing professionals, but challenges persist

