

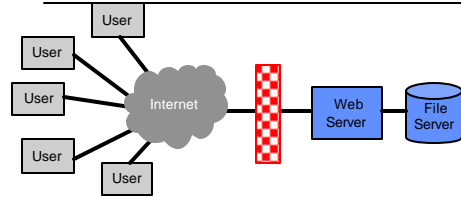
CS 194-1 (CS 161)  
Computer Security

Lecture 12

Web Security and Intrusion Detection

October 11, 2006  
Prof. Anthony D. Joseph  
<http://cs161.org/>

Review: Firewalls



- Default firewall rule: deny all
- Other firewall rules:
  - allow tcp \*.\*/\*out -> <web server IP>:80/in
  - allow tcp \*.\*/\*out -> <web server IP>:443/in
  - drop \* \*.\* \* -> \*.\* \*

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Goals for Today

- Web Servers
  - Static and Dynamic Content
- Adding a DMZ to a Firewall
- Secure Topologies
- Intrusion Detection History
- Network-based Host Compromise
- Host-based Network Intrusion Detection
  - Signature-based, Anomaly-based
- Distributed Network Intrusion Detection
  - Honeypots, Tarpits
- An attack against an IDS

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Polls

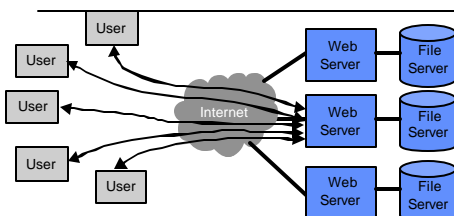
- How many people have set up a personal web server?
- How many people have set up a business web server?

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Web Servers



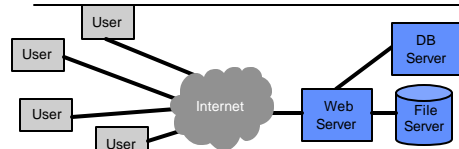
- Web server serves up static, read-only content from file server
- Scales up by replicating web servers
  - Can use DNS round-robin or load balancer

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Web Servers



- Add a database server for dynamic content
  - DB used to store per-user info or site content
  - Also, used for authentication, read/write actions, e-commerce, ...
- Software connector to DB server
  - Object/Java DataBase Connectivity

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## Web Servers

- **Static content model:**
  - Web server uses file server for static content, templates, ...
- **Dynamic content model:**
  - Web server uses database server to retrieve/store dynamic content
- **Can have mixtures**
  - Ex: Storing dynamic content in FS
  - Ex: Storing static content in DB
- **What are the security issues?**

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## Some Web Server Threats and Attacks

- **Replace static content ("defacement")**
  - Exploit vulnerability to access Web or File servers
- **(Distributed) Denial of Service attack**
  - Request large image or emulate complex transaction
- **Unauthorized database access**
  - Exploit vulnerability (e.g., SQL injection) to read/write database
- **Attack server OS or other services**
  - Exploit vulnerability to disable server

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## Replace Static Content ("Defacement")

- **Cracker exploits a vulnerability to gain access to Web or File Servers**
- **Examples:**
  - Flaws in CGI programs
  - Flaws in URL processing
  - Buffer overflows
- **Replaces web pages with their own**
- **May also access protected content**

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## (Distributed) Denial of Service Attack

- **Cracker performs resource exhaustion attack**
  - Overwhelm network, CPU, disk bandwidth, ...
- **Examples**
  - Request large image or file
  - POST large image or file (requires many zombies)
  - Emulate complex transaction
- **Typically use large number of zombies (1,000's to 100,000's)**

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## Unauthorized Database Access

- **Cracker exploits vulnerability in Web server to DB server connection to read/write database**
- **Example:**
  - Use URL or POST attack to inject SQL code
  - Gain access to Web server, then connect to DB
- **Attacks can compromise DB integrity**

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## Attack Server OS or Other Services

- **Cracker exploits vulnerability to gain access to server**
  - Many OS and service vulnerabilities...
- **Can be a stepping stone to attacking web service or accessing database**

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## Stopping Some Attacks

- Replace static content (“defacement”)
  - Harden server (latest patch levels, minimum services)
  - Limit data on file server
- (Distributed) Denial of Service attack
  - Add load balancer, DNS round-robin, replicated clusters, ...
- Unauthorized database access
  - Harden server (latest patch levels, min. svcs)
  - Sanity check all arguments
- Attack server OS or other services
  - Harden servers (latest patch levels, min. svcs)

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## Problems

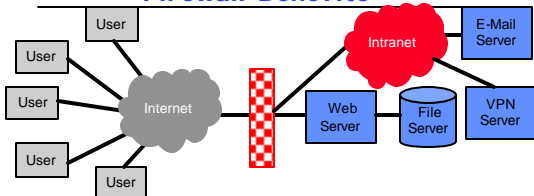
- Hard to keep servers up-to-date with patches
  - Zero-day exploits
  - Delays in releasing, retrieving, testing, installing patches
- DDoS attacks still impose load on servers
- Add layered defense
  - Place firewall between Internet and Web server

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## Firewall Benefits



- Helps harden servers by blocking all but web traffic
- DDoS attacks: add stateful rules or block zombie IP subnets
  - Doesn't work for all content attacks
- Intranet and e-mail server access?

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## Firewall Issues

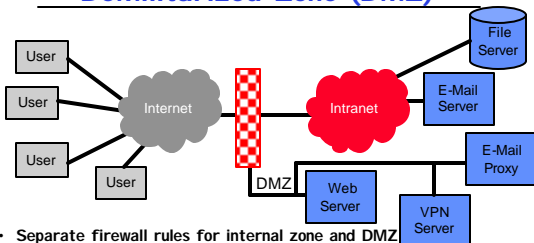
- We can add more rules
  - For access to Intranet, E-mail server, and other “public” servers
- But, what happens if one server or Intranet machine is compromised?
- This is the classic firewall problem:
  - All our machines are now vulnerable!
- Real issue:
  - We need to both protect public servers and Intranet
- Solution: Place public servers in a DMZ

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## DeMilitarized Zone (DMZ)



- Separate firewall rules for internal zone and DMZ
  - Internet-DMZ rules only allow web, e-mail traffic
  - DMZ-Intranet rules only allow file, e-mail, remote login from DMZ
  - No Internet-Intranet access
- Where to place e-mail server?
  - Add proxy to isolate e-mail access/storage from e-mail forwarding

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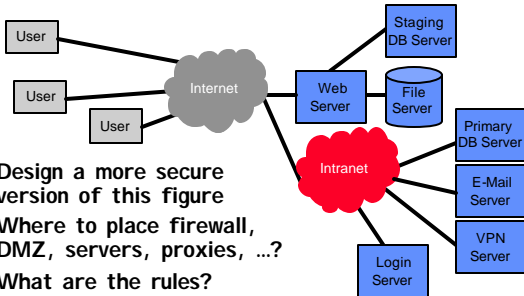
## Administrivia

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## Designing a More "Secure" Topology



- Design a more secure version of this figure
- Where to place firewall, DMZ, servers, proxies, ...?
- What are the rules?

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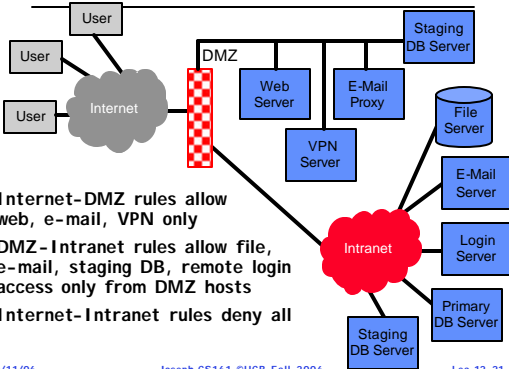
## Design

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## One Solution



- Internet-DMZ rules allow web, e-mail, VPN only
- DMZ-Intranet rules allow file, e-mail, staging DB, remote login access only from DMZ hosts
- Internet-Intranet rules deny all

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## Web Security Summary

- Public servers are vulnerable to attack
  - OS and services
- Eliminate unnecessary services
- Apply all patches
- Use a DMZ to provide layered defense
  - Place server/proxy in DMZ
  - Place database/file/"real" servers in Intranet
  - Deny all default for Internet-Intranet traffic

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## Intrusion Detection History

- Detecting attempts to penetrate our systems
  - Used for post-mortem activities
  - Related problem of extrusion (info leaking out)
- In pre-network days (centralized mainframes)...
  - Primary concern is abuse and insider information access/theft
  - Reliance on logging and audit trails
- But, highly labor intensive to analyze logs
  - What is abnormal activity?
  - Ex: IRS employees snooping records
  - Ex: Moonlighting police officers

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## Network-based Host Compromises

- How do remote intruders gain access?
- They attempt network-based attacks that exploit OS & app bugs
  - Ex: Denial of service, spyware install, zombie, ...

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## Host-based Net Intrusion Detection

- At each host, monitor all incoming and outgoing network traffic - for each packet:
  - Analyze 4-tuple and protocol
  - Examine contents
  - ...
- Challenge: Separate "signal" from "noise"
  - **Signal** is an attack (intrusion)
  - **Noise** is normal "background" traffic
  - Assumption: can separate signal and noise...

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## Some Challenges

- What is normal traffic?
  - Server, desktop, PDA, PDA/phone, ...
  - My normal traffic ? your normal traffic
  - Lots of data for servers
- Why do we need sufficient signal and noise separation?
  - To avoid too many false alarms!
- What happens if signals are missed?
  - Possible intrusion!

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## Some Common False Positives

- Proximity probes
  - Website load balancers will probe your machine for proximity
  - Connect to website hosted by mirror-image.com, and >10 load balancers in 6 countries probe your machine
- Stale IP caches
  - Using dynamic IP addresses, you may get the "old" address of someone who was running a P2P app
  - Peers continue to try to "re-connect"
- Web posts with dynamic IP addresses
  - Spiders crawl machine currently using IP address

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## Lots and Lots of Data!!

- Network trace from Win2K desktop

```

ZoneAlarm Logging Client v3.7.202
Windows 2000 -5.0.2195-Service Pack 4-SP
type,date,time,source,destination,transport
FWIN,2004/01/15,13:17:38 -8:00 GMT,216.183.33.67,42445,128.32.168.229-6129,TCP (flags:S)
FWOUT,2004/01/15,13:18:00 -8:00 GMT,128.32.168.229-5000,68.26.217.204-5000,UDP
FWIN,2004/01/15,13:42:38 -8:00 GMT,61.178.60.11:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,13:42:48 -8:00 GMT,62.177.227.10:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,13:48:12 -8:00 GMT,128.32.41.80,1040,128.32.168.229-38293,UDP
FWIN,2004/01/15,13:58:30 -8:00 GMT,24.224.253.230-2446,128.32.168.229-6129,TCP (flags:S)
FWIN,2004/01/15,14:04:40 -8:00 GMT,80.116.4.42:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWOUT,2004/01/15,14:04:44 -8:00 GMT,128.32.168.229-5000,68.26.217.204-5000,UDP
FWIN,2004/01/15,14:07:36 -8:00 GMT,210.217.129.194-3598,128.32.168.229-1433,TCP (flags:S)
FWIN,2004/01/15,14:15:00 -8:00 GMT,128.32.30.70:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,14:23:20 -8:00 GMT,80.56.148.243:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,14:41:48 -8:00 GMT,194.23.44.215:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,14:43:08 -8:00 GMT,61.64.246.192:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWOUT,2004/01/15,14:43:16 -8:00 GMT,128.32.168.229-5000,68.26.217.204-5000,UDP
FWIN,2004/01/15,15:02:00 -8:00 GMT,128.32.168.21:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,15:06:28 -8:00 GMT,81.185.244.166:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,15:43:44 -8:00 GMT,217.255.55.163:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWOUT,2004/01/15,15:44:16 -8:00 GMT,128.32.168.229-5000,68.26.217.204-5000,UDP
FWIN,2004/01/15,15:50:06 -8:00 GMT,65.78.10.110-3071,128.32.168.229-3410,TCP (flags:S)
FWIN,2004/01/15,15:59:42 -8:00 GMT,202.42.49.198:0,128.32.168.229:0,ICMP (type:8/subtype:0)
FWIN,2004/01/15,16:07:40 -8:00 GMT,68.22.89.249-4088,128.32.168.229-1433,TCP (flags:S)
FWIN,2004/01/15,16:08:36 -8:00 GMT,193.95.219.45:0,128.32.168.229:0,ICMP (type:3/subtype:0)
FWIN,2004/01/15,16:23:50 -8:00 GMT,67.37.40.15-4299,128.32.168.229-3410,TCP (flags:S)
FWOUT,2004/01/15,16:24:16 -8:00 GMT,128.32.168.229-5000,68.26.217.204-5000,UDP
    
```

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## Trace Analysis

- ZoneAlarm Logging Client v3.7.202
  - Windows 2000-5.0.2195-Service Pack 4-SP
  - type,date,time,source,destination,transport
  - FWIN,2004/01/15,13:17:38 -8:00 GMT,216.183.33.67,42445,128.32.168.229-6129,TCP (flags:S)
  - FWOUT,2004/01/15,13:18:00 -8:00 GMT,128.32.168.229-5000,68.26.217.204-5000,UDP
  - FWIN,2004/01/15,13:42:38 -8:00 GMT,61.178.60.11:0,128.32.168.229:0,ICMP (type:8/subtype:0)
  - FWIN,2004/01/15,13:42:48 -8:00 GMT,62.177.227.10:0,128.32.168.229:0,ICMP (type:8/subtype:0)
  - FWIN,2004/01/15,13:42:48 -8:00 GMT,128.32.168.229:0,128.32.168.229:0,ICMP (type:8/subtype:0)
- b2b-33-67.ip.granderiver.com
- "ping" probe
- Used by the Dameware remote admin sw (old versions have a bug allowing unauthorized login). Dameware also installed by some viruses

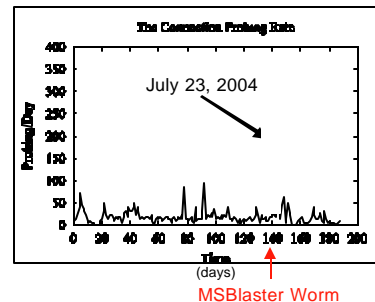
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## Analyzing Host-based Trace Data

- TCP connection probes on port 445
- Day 0 is 2003/03/04



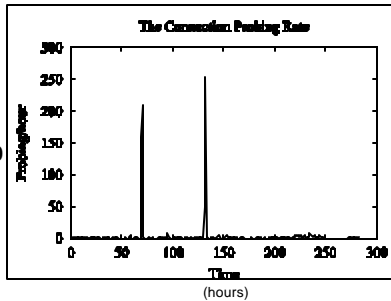
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## MSBlaster in Detail

- TCP 445 probes/hr
- Hour 0 is 15:20 on 2003/07/20



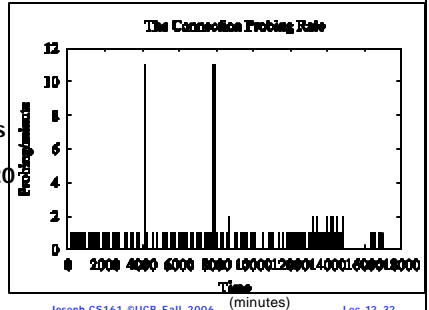
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## MSBlaster in More Detail

- TCP 445 probes / 10 min
- Minute 0 is 15:20 on 2003/07/20



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## Example Common Attack

- Port scanning a host
  - Trying to connect/send data to different ports/protocols: sequential scan of host
  - Nmap tool (<http://www.insecure.org/nmap/>)
    - » Determines OS/hostname/device type detection via service fingerprinting (ex: SGI IRIX has svc on TCP port 1)
    - » Determines what svc is really listening on a port and can even determine app name and version
    - » Operates in optional obfuscation mode
- How to detect attack?

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## Intrusion Detection Using Signals

- This is a misuse detection problem
  - Similar problem to virus detection
  - "Match what you know"
- High-level solution:
  - Collect info about attack methods and types
    - » 4-tuple/protocol
    - » Packet contents
  - Create and look for signatures
    - » Slammer packet, port scan, ...

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## Intrusion Detection Using Noise

- This is an anomaly detection problem
  - Need to learn normal behavior
  - "Match what's different"
- High-level solution:
  - Try to identify what is normal traffic
    - » Common 4-tuple/protocol
  - Heuristic: Look for major deviations (outliers)
    - » Ex: unusual target port, source addr, or port sequence (scan)
  - Apply AI: Statistical Learning Techniques

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## Signature Detection

- Language to specify intrusion patterns
  - 4-tuple/protocol and potential intrusion values
    - » Ex: External host → file server (port 110, 135, ...)
    - » Ex: Internal workstation → external P2P host
  - Packet contents
    - » Could be single or multiple packets (stream reconstruction)
  - Sequence of 4-tuple/protocol and packets
    - » Also, model of protocol/app finite state machine
- Lots of state in pattern matching engine
- Example rule:
 

```
-alert tcp any any -> myip 21 (content:"site exec"; content:"%"; msg:"site exec buffer overflow attempt");
```

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## Signature Detection

- Snort tool (<http://www.snort.org/>)
  - 2 million downloads, 100,000+ active users,
- Advantages
  - Very low false positive (alarm) rate
- Disadvantages
  - Only able to detect already known attacks
  - Simple changes to attack can defeat detection
    - » Ex: Scan every even port, then every odd port...

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## Anomaly Detection

- Analyze normal operation (behavior), look for anomalies
  - Uses AI techniques: Statistical Learning Techniques
  - Compute statistical properties of "features"
    - » 4-tuple, protocol, packet contents, packets/sec, range of port numbers, ...
  - Report errors if statistics are outside of "normal" range

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## Anomaly Detection

- Advantages
  - Can recognize "evolved" and new attacks
- Disadvantages
  - High false positive rate (alarms)
  - May have delayed alarm
  - Some attacks can hide in "normal" traffic
  - SLT requires training on known good data
  - Hard to capture protocol state behavior (FSM)
  - Problems when what's "normal" changes
    - » Ex: flash crowds

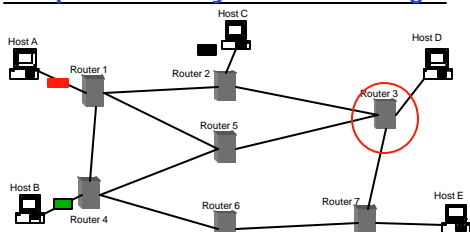
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**BREAK**

## Super Stealthy Port Scanning



- Use many zombies (each scans a few ports/hour of target)
  - Each zombie is assigned many machines to scan
- Fast to scan both one machine, and many
- Very hard to detect at targets!

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## Distributed Intrusion Detection

- Place appliance in the network at choke point or, *share results across machines*
- Apply signature or anomaly detection across larger data set
- Advantages:
  - Easier to detect stealth probes of large number of machines
- Disadvantages:
  - Large amount of data to communicate

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## Honeypots

- Closely monitored network decoys
- May distract adversaries from more valuable machines on a network
- May provide early warning about new attack and exploitation trends
  - Enables in-depth examination of adversaries during and after exploitation

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## Honeypots

- Can simulate one or more network services on one or more machines
  - Can have virtual cluster of machines
- Causes an attacker to think you're running vulnerable services that can be used to break into the machine
  - Can log access attempts to those ports, including the attacker's source IP and keystrokes
  - Can watch attacker in real-time and trace back/forward
- Provides advanced warning of an attack
  - Could use to automate generation of new firewall rules

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## Tarpits

- A very, very sticky honeypot...
- Set up network decoy
  - For each port we want to "tarpit," we allow connections to come in, but don't let them out
- Idea:
  - Slow down scanning tools/worms to kill their performance/propagation because they rely on quick turnarounds
  - Might also give us time to protect real hosts

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## Example Tarpit Implementation

- Accept any incoming TCP connection
- When data transfer begins to occur, set TCP window size to zero, so no data can be transferred within the session
- Hold the connection open, and ignore any requests by remote side to close session
- Attacker must wait for the connection to timeout in order to disconnect

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## Tarpits

- Advantages
  - Can customize for specific worms
    - » Ex: analyze incoming packets to port 80 and only tarpit web connections from worms
      - look for "cmd.exe" (CodeRed) or "default.ida" (Nimda)
- Disadvantages
  - Might trap valid host
  - Can cause some operating systems to crash

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## Intrusion Prevention Systems

- We can detect intrusions, so why not automatically cut off network connections to compromised hosts?
- Intrusion Prevention Systems do this
- But, what if we're wrong...
  - Possible Denial of Service - trick IPS into thinking host is compromised
  - Turn off access our airline reservation server when a fare deal causes very high/different traffic patterns

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## Witty Worm

- March 04: Attacked the IDS
- Targeted a buffer overflow vulnerability in several of a vendor's IDS products
- Deletes a randomly chosen sectors of hard drives over time killing system
- Payload contained phrase:
  - "(^.^) insert witty message here (^.^)"

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## Witty's Many Firsts

- First widely propagated Internet worm with a destructive payload
- First worm with order of magnitude larger hit list than any previous worm
- Shortest known interval between vulnerability disclosure and worm release - 1 day
- First to spread through nodes doing something proactive to secure their computers / networks
- Spread through a population almost an order of magnitude smaller than that of previous worms

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## Intrusion Detection Systems Summary

- Ongoing arms race between attackers and detection technologies
- Real challenge is false positive rate
  - Renders most IDS useless - alerts ignored
- Adaptive, anomaly detection is promising, but still lacking
- IPS products are still immature and problematic
- IDS products are now targets

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## Summary

- Avoiding attacks against public servers:
  - Eliminate unnecessary services
  - Apply all patches
  - Use a DMZ to provide layered defense
- Intrusion detection is hard!
  - Crying wolf syndrome
  - Immature products
  - We need new adaptive techniques
- Ongoing arms race between attackers and defenders

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