

Random Number Generation

- Many crypto protocols require parties to generate random numbers
 - Key generation
 - Generating nonces
- How to generate random numbers?
 - Step 1: how to generate truly random bits?
 - Step 2: crypto methods to stretch a little bit of true randomness into a large stream of pseudorandom values that are indistinguishable from true random bits (PRNG)

Case Study

- Random number generation is easy to get wrong
- Can you spot the problems in this example?

```
unsigned char key[16];
```

```
srand(time(null));
for (=0; i<16; i++)
key[i] = rand(). & 0xFF;
where
static unsigned int next = 0;
void srand(unsigned int seed) {
    next = seed;
    j
    int rand(vr:id) {
```

```
int rand(void) {
    next = next * 1103515245 + 12345;
    return next % 32768;
```

}

Real-world Examples

- X Windows "magic cookie" was generated using rand()
- Netscape browsers generated SSL session keys using time & process ID as seed (1995)
- Kerberos
 - First discover to be similarly flawed
 - -4 yrs later, discovered flaw with memset()
- PGP used return value from read() to seed its PRNG, rather than the contents of buffer
- On-line poker site used insecure PRNG to shuffle cards

Lessons Learned

- Seeds must be unpredictable
- Algorithm for generating pseudorandom bits must be secure

Generating Pseudorandom Numbers

- True random number generator (TRNG)
 - Generates bits that are distributed uniformly at random, so that all outputs are equally likely, with no patterns, correlations, etc.
- Cryptographically secure pseudorandom number generator (CS-PRNG)
 - Taking a short true-random seed, and generates long sequence of bits that is computationally indistinguishable from true random bits

CS-PRNG

- CS-PRNG: cryptographically secure pseudorandom number generator
 - G: maps a seed to an output G(S)
 - » E.g., G: {0,1}¹²⁸ -> {0,1}¹⁰⁰⁰⁰⁰⁰
 - Let K denote a random variable distributed uniformly at random in domain of G
 - Let U denote a random variable distributed uniformly at random in range of G
 - G is secure if output G(K) is computationally indistinguishable from U

Sample construction

-Use the seed as a key k, and compute AES-CBC(k, 0ⁿ)

TRNG (I)

- TRNG should be random and unpredictable
- Good or bad choices?
 - IP addresses
 - Contents of network packets
 - Process IDs
 - High-speed clock
 - Soundcard
 - Keyboard input
 - Disk timings

TRNG (II)

How to convert non-uniform sources of randomness into TRNG?

- Use a cryptographic hash function, such as SHA1
- Suppose x is a value from an imperfect source, or a concatenation of values from multiple sources, and it is impossible for an attacker to predict the exact value x except with probability 1/2ⁿ
- Then hash(x) truncated to n bits should provide a n-bit value that is uniformly distributed, if hash() is secure

Administrative Matters

HW2 graded
Mean: 41.7
Standard deviation: 13.2
1st quartile: 39.8
2nd quartile (median): 44.5
3rd quartile: 50.0
Maximum: 57.0











Desired Properties for Ecash

- · Anonymous: bank should not know how Alice spends her money
- Prevent forging
- Prevent double spending



Ecash Using Blind Signature

- How to use blind signature to build ecash?
- A valid \$1 bill is a pair (x,y), where y = hash(x)^d mod N, hash() is one-way function
- How does the ecash protocol work?
- Why do we need hash()?
- How to prevent double spending?
- What to do for different denominations? -Nickles, dimes, dollars

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Other Methods for Ecash

- Use zero-knowledge proofs (out of scope)
 - More building blocks of ZKP
 - Support many properties
 - » Identifying double spenders

Conclusion

- Random number generator
 - CS-PRNG
 - » Definition» How to construct it?
 - TRNG
- Ecash
 - Example of the power of crypto
 - Blind signatures

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