

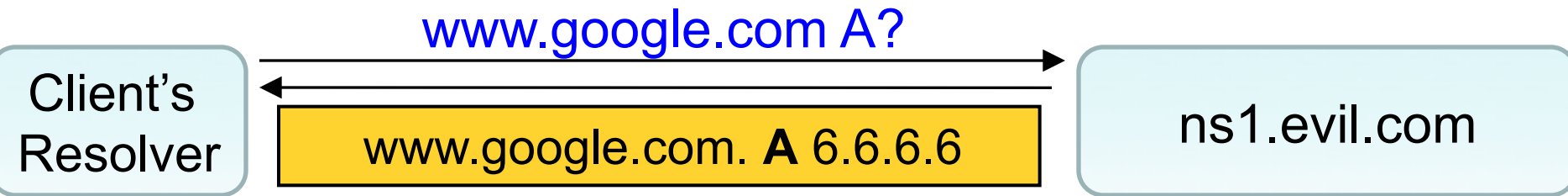
# TLS and DNSSEC wrap-up

***CS 161: Computer Security***

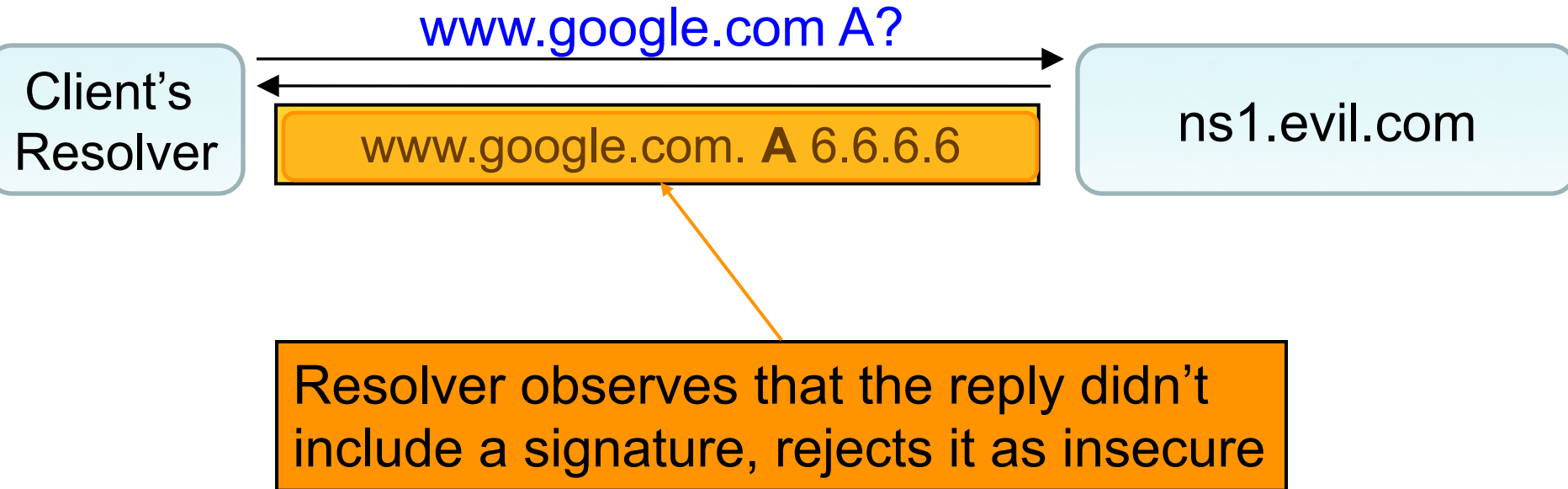
**Prof. David Wagner**

**April 14, 2013**

# DNSSEC – Mallory attacks!



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www.google.com A?

Client's  
Resolver

www.google.com. **A** 6.6.6.6  
www.google.com **RRSIG A**  
*signature-of-the-A-record-using-evil.com's-key*

ns1.evil.com

# DNSSEC – Mallory attacks!

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www.google.com A?

ns1.evil.com

www.google.com. **A** 6.6.6.6

www.google.com **RRSIG A**

*signature-of-the-A-record-using-evil.com's-key*

(1) If resolver didn't receive a signature from .com for evil.com's key, then it can't validate this signature & ignores reply since it's not properly signed ...

# DNSSEC – Mallory attacks!

Client's  
Resolver

www.google.com A?

ns1.evil.com

www.google.com. **A** 6.6.6.6

www.google.com **RRSIG A**

*signature-of-the-A-record-using-evil.com's-key*

(2) If resolver *did* receive a signature from .com for evil.com's key, then it knows the key is for evil.com and not google.com ... and ignores it

# DNSSEC – Mallory attacks!

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Client's  
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www.google.com **RRSIG A**  
*signature-of-the-A-record-using-*  
*google.com's-key*

ns1.evil.com

# DNSSEC – Mallory attacks!

www.google.com A?

Client's  
Resolver

ns1.evil.com

www.google.com. **A** 6.6.6.6

www.google.com **RRSIG A**

*signature-of-the-A-record-using-  
google.com's-key*

If signature **actually** comes from google.com's key, resolver will believe it ...

... but no such signature should exist unless either:

(1) google.com *intended* to sign the RR, or

(2) google.com's private key was compromised



# Issues With DNSSEC ?

- Issue #1: Replies are Big
  - E.g., “dig +dnssec berkeley.edu” can return 2100+ B
  - DoS **amplification**
  - Increased **latency** on low-capacity links
  - Headaches w/ older libraries that assume replies < 512B
- Issue #2: *Partial deployment*
  - Suppose **.com** not signing, though **google.com** is
  - Major practical concern. What do we do?
  - Can wire additional key into resolver (doesn't **scale**)
  - Or: outsource to trusted third party (“lookaside”)
    - Wire their key into resolver, they sign numerous early adopters

# Issues With DNSSEC, cont.

- Issue #1: *Partial deployment*
  - Suppose `.com` not signing, though `google.com` is. Or, suppose `.com` and `google.com` are signing, but `cnn.com` isn't. Major practical concern. What do we do?
  - What do you do with unsigned/unvalidated results?
  - If you trust them, **weakens incentive** to upgrade (man-in-the-middle attacker can defeat security even for `google.com`, by sending forged but unsigned response)
  - If you don't trust them, a whole lot of things **break**

# Issues With DNSSEC, cont.

- Issue #2: Negative results (“no such name”)
  - What statement does the nameserver sign?
  - If “gab1uph.google.com” doesn’t exist, then have to do dynamic key-signing (expensive) for any bogus request
  - Instead, sign (off-line) statements about order of names
    - E.g., sign “`gabby.google.com` is followed by `gabrunk.google.com`”
    - Thus, can see that `gab1uph.google.com` can’t exist
  - But: now attacker can **enumerate** all names that exist :-)

# Issues With DNSSEC, cont.

- Issue #3: *Whom do you really trust?*
  - For your laptop (say), who does all the “grunt work” of fetching keys & validating DNSSEC signatures?
- Your laptop’s local resolver?
  - ... which you acquire via DHCP in your local coffeeshop
  - I.e., exactly the most-feared potentially **untrustworthy** part of the DNS resolution process!
- Alternatives?
  - ⇒ Your laptop needs to do all the validation work itself :-)

# TLS/SSL Trust Issues

- “*Commercial certificate authorities protect you from anyone from whom they are unwilling to take money.*”  
– Matt Blaze, circa 2001
- ... and there are lots of CAs, and we must trust them all.
- Of course, it's not just their greed that matters ...

# Fraudulent Google certificate points to Internet attack

Is Iran behind a fraudulent Google.com digital certificate? The situation is similar to one that happened in March in which spoofed certificates were traced back to Iran.

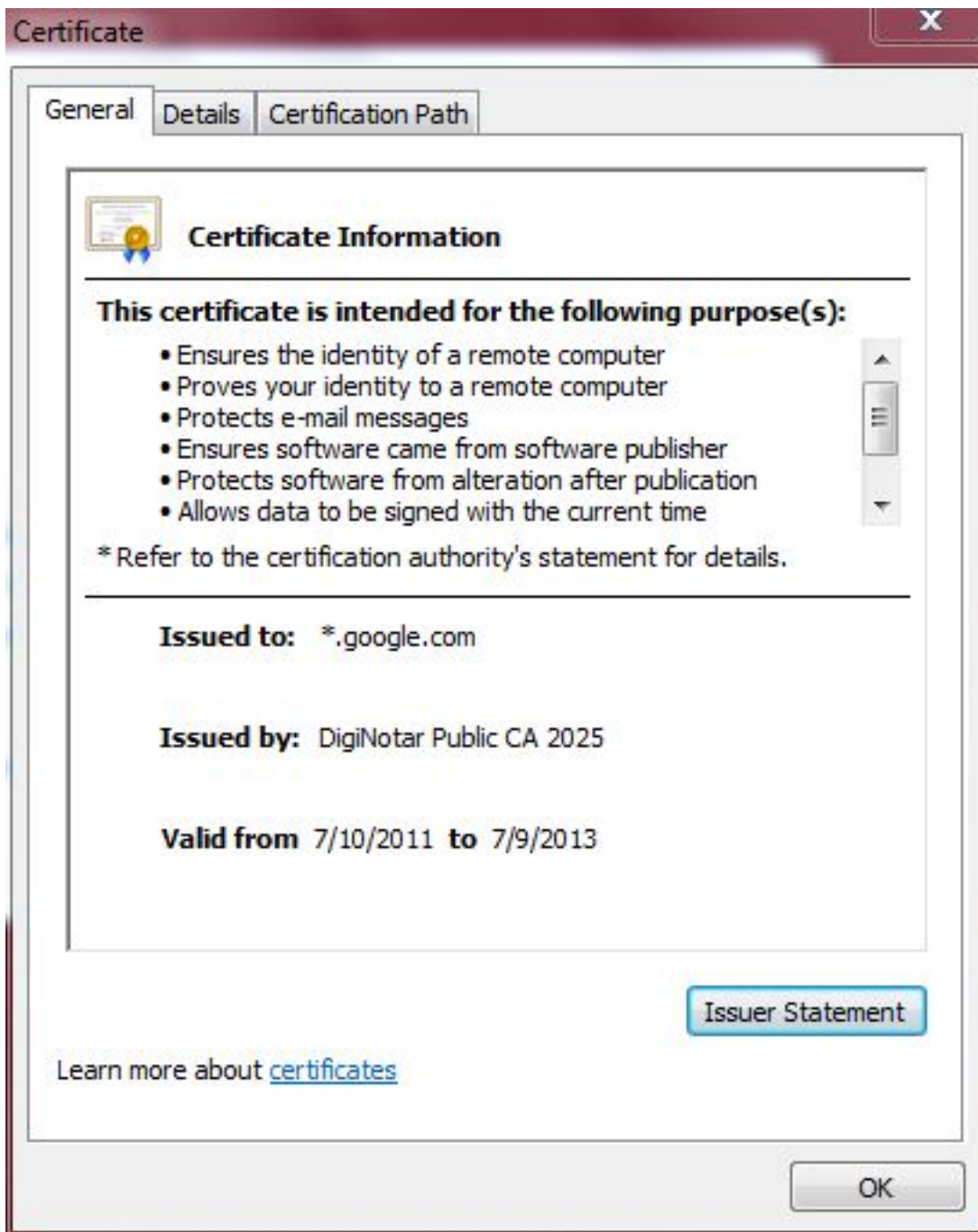


by [Elinor Mills](#) | August 29, 2011 1:22 PM PDT



A Dutch company appears to have issued a digital certificate for Google.com to someone other than Google, who may be using it to try to re-direct traffic of users based in Iran.

Yesterday, someone reported on a Google support site that when attempting to log in to Gmail the browser issued a warning for the digital certificate used as proof that the site is legitimate, according to [this thread](#) on a Google support forum site.



This appears to be a **fully valid** cert using normal browser validation rules.

Only detected by Chrome due to its recent introduction of cert “**pinning**” – requiring that certs for certain domains **must** be signed by specific CAs rather than any generally trusted CA

October 31, 2012, 10:49AM

# Final Report on DigiNotar Hack Shows Total Compromise of CA Servers

The attacker who penetrated the Dutch CA DigiNotar last year had complete control of all eight of the company's certificate-issuing servers during the operation and he may also have issued some rogue certificates that have not yet been identified. The final report from a

## Evidence Suggests DigiNotar, Who Issued Fraudulent Google Certificate, Was Hacked *Years Ago*

from the *diginot* dept

The big news in the security world, obviously, is the fact that a **fraudulent Google certificate made its way out into the wild**, apparently targeting internet users in Iran. The Dutch company DigiNotar has put out a statement saying that **it discovered a breach** back on July 19th during a security audit, and that fraudulent certificates were generated for "several dozen" websites. The only one known to have gotten out into the wild is the Google one.



# TLS/SSL Trust Issues

- *“Commercial certificate authorities protect you from anyone from whom they are unwilling to take money.”*
  - Matt Blaze, circa 2001
- ... and there are lots of CAs, and we must trust them all.
- Of course, it's not just their greed that matters ...
- ... and it's not just their diligence & security that matters ...
  - *“A decade ago, I observed that commercial certificate authorities protect you from anyone from whom they are unwilling to take money. That turns out to be wrong; they don't even do that much.”* - Matt Blaze, circa 2010

# Law Enforcement Appliance Subverts SSL

By [Ryan Singel](#)  March 24, 2010 | 1:55 pm | Categories: [Surveillance](#), [Threats](#)



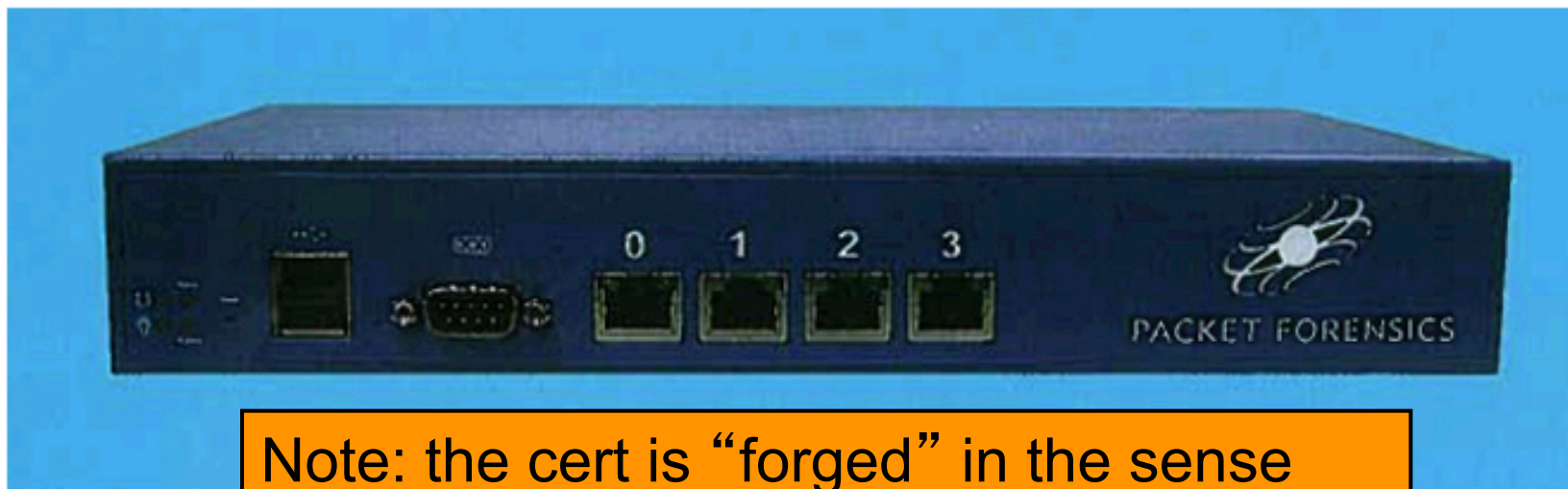
That little lock on your browser window indicating you are communicating securely with your bank or e-mail account may not always mean what you think it means.

Normally when a user visits a secure website, such as Bank of America, Gmail, PayPal or eBay, the browser examines the website's certificate to verify its authenticity.

At a recent wiretapping convention, however, security researcher Chris Soghoian discovered that a small company was marketing internet spying boxes to the feds. The boxes were designed to intercept those communications — without breaking the encryption — by using forged security certificates, instead of the real ones that websites use to verify secure connections. To use the appliance, the government would need to acquire a forged certificate from any one of more than 100 trusted Certificate Authorities.

# Law Enforcement Appliance Subverts SSL

By Ryan Singel  March 24, 2010 | 1:55 pm | Categories: [Surveillance](#), [Threats](#)



Note: the cert is “forged” in the sense that it doesn’t really belong to Gmail, PayPal, or whomever. But it does not *appear* forged because it includes a legitimate signature from a trusted CA.

That little lock of  
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## Security Warning: Do you trust the Russian government?

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Firefox has detected that your connection to this website is probably not secure. If you are attempting to access or transmit sensitive data, you should **stop** this task, and try again using a **different Internet connection**.

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Firefox has detected a potential security problem while trying to access [www.bankofamerica.com](http://www.bankofamerica.com), a website visited at least 131 times in the past by persons using this computer.

In these previous browsing sessions, [www.bankofamerica.com](http://www.bankofamerica.com) provided a security certificate verified by a company in the **United States**.

However, this website is now presenting a different security certificate verified by a company based in **Russia**.

If you do not trust the government of Russia with your private data, or think it unlikely that Bank of America would obtain a security certificate from a company based there, this could be a sign that someone is attempting to intercept your secure communications.

[Click here](#) to learn more about security certificates and this potentially risky situation.

If you trust the government of Russia and companies located there to protect your privacy and security, [click here](#) to accept this new certificate and continue with your visit to the site.

Get me out of here!

# Summary of TLS & DNSSEC Technologies

- **TLS:** provides **channel security** (for communication over TCP)
  - Confidentiality, integrity, authentication
  - Client & server agree on crypto, session keys
  - Underlying security dependent on:
    - Trust in **Certificate Authorities** / decisions to sign keys
    - (as well as implementors)
- **DNSSEC:** provides **object security** (for DNS results)
  - Just **integrity** & **authentication**, not confidentiality
  - No client/server setup “dialog”
  - Tailored to be **caching-friendly**
  - Underlying security dependent on trust in Root Name Server’s key, and all other signing keys





**POLICE**





# Tamper-Evident Logging

- We work for the police Electronic Records office.
- To ensure that evidence can't be questioned in court, we want to make sure that evidence can't be tampered with, after it is logged with the office.
- In other words: a police officer can log an electronic file at any time; after it is logged, no back-dating or after-the-fact changes to evidence should be possible.
- How should we do it? What crypto or data structures could we use?



# Design Problem for You

- Idea: Each day, collect all the files  $(f_1, f_2, \dots, f_n)$  that are logged that day. Then, publish something in the next day's newspaper, to commit to these files.
- Question: What should we publish?  
Needs to be short, and ensure after-the-fact changes or backdating are detectable.
- When a file  $f_i$  is exhibited into evidence in a trial, how can judge verify it hasn't been modified post-facto?

# Your Solution

- Store in database:  $f_1, \text{Sign}(f_1), f_2, \text{Sign}(f_2), \dots, f_n, \text{Sign}(f_n)$
- Publish: public key
- To verify  $f_i$ : reveal  $f_1, \text{Sign}(f_i)$
  
- Critique: Sysadmin can get a copy of the private key, modify database, update the signature, and thus modify old entries or create new backdated entries.

# Your Solution

- Publish:  $H(f_1, f_2, \dots, f_n)$
- To verify  $f_i$ : reveal  $f_1, f_2, \dots, f_n$

# Solution

- Each day, collect all the files  $(f_1, f_2, \dots, f_n)$  that are logged that day. Then, publish  $H(f_1, f_2, \dots, f_n)$  in the next day's newspaper, to commit to these files.
- When a file  $f_i$  is exhibited into evidence in a trial, reveal  $f_1, f_2, \dots, f_n$  to judge. Judge can hash them, check that their hash was in the right day's newspaper, and check that  $f_i$  is in the list.

# Better Solution

- Each day, collect all the files  $(f_1, f_2, \dots, f_n)$  that are logged that day. Let  $f_0$  be the previous day's hash. Publish  $H(f_0, f_1, f_2, \dots, f_n)$  in the next day's newspaper, to commit to these files.
- Note that exhibiting file  $f_i$  into evidence still requires revealing entire list of other files  $(f_1, f_2, \dots, f_n)$  that were logged that day. Can you think of any way to avoid that?

# Take-away

- Using hash chaining, we can provide tamper-evident audit logs that let us detect after-the-fact modifications and backdated entries.