CS162 Operating Systems and Systems Programming Lecture 2

## History of the World Parts 1—5 Operating Systems Structures

August 31, 2005 Prof. John Kubiatowicz http://inst.eecs.berkeley.edu/~cs162

#### Review: What does an Operating System do?

- Silerschatz and Gavin: "An OS is Similar to a government"
  - Begs the question: does a government do anything useful by itself?
- Coordinator and Traffic Cop:
  - Manages all resources
  - Settles conflicting requests for resources
  - Prevent errors and improper use of the computer
- Facilitator:
  - Provides facilities that everyone needs
  - Standard Libraries, Windowing systems
  - Make application programming easier, faster, less error-prone
- Some features reflect both tasks:
  - E.g. File system is needed by everyone (Facilitator)
  - But File system must be Protected (Traffic Cop)

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#### **Review: Virtual Machine Abstraction**

# Application

– Virtual Machine Interface

## **Operating System**

– Physical Machine Interface

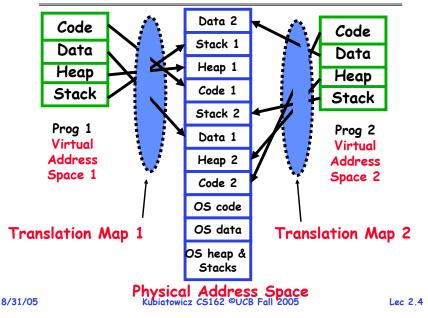
# Hardware

- Software Engineering Problem:
  - Turn hardware/software quirks  $\Rightarrow$  what programmers want/need
  - Optimize for convenience, utilization, security, reliability, etc...
- For Any OS area (e.g. file systems, virtual memory, networking, scheduling):
  - What's the hardware interface? (physical reality)
  - What's the application interface? (nicer abstraction)

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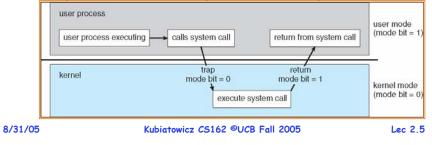
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#### Review: Example of Address Translation



#### **Review: Dual Mode Operation**

- Hardware provides at least two modes:
  - "Kernel" mode (or "supervisor" or "protected")
  - "User" mode: Normal programs executed
- Some instructions/ops prohibited in user mode:
  - Example: cannot modify page tables in user mode
    - » Attempt to modify  $\Rightarrow$  Exception generated
- Transitions from user mode to kernel mode:
  - System Calls, Interrupts, Other exceptions



### Moore's law change

	1981	2005	Factor
CPU MHz,	10	3800	380
Cycles/inst	3—10	0.25-0.5	6—40
DRAM capacity	128KB	4GB	32,768
Disk capacity	10MB	1TB	100,000
Net bandwidth	9600 b/s	1 Gb/s	110,000
# addr bits	16	32	2
#users/machine	10s	≤ <b>1</b>	≤ 0.1
Price	\$25,000	\$4,000	0.2

### Typical academic computer 1981 vs 2005

### Goals for Today

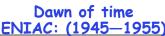
- History of Operating Systems
   Really a history of resource-driven choices
- Operating Systems Structures
- Operating Systems Organizations

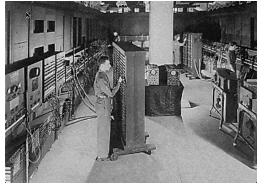
Note: Some slides and/or pictures in the following are adapted from slides ©2005 Silberschatz, Galvin, and Gagne

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• "The machine designed by Drs. Eckert and Mauchly was a monstrosity. When it was finished, the ENIAC filled an entire room, weighed thirty tons, and consumed two hundred kilowatts of power."

• http://ei.cs.vt.edu/~history/ENIAC.Richey.HTML

#### History Phase 1 (1948—1970) Hardware Expensive, Humans Cheap

- When computers cost millions of \$'s, optimize for more efficient use of the hardware!
   Lack of interaction between user and computer
- User at console: one user at a time
- Batch monitor: load program, run, print
- Optimize to better use hardware
  - When user thinking at console, computer idle $\Rightarrow$ BAD!
  - Feed computer batches and make users wait
  - Autograder for this course is similar
- No protection: what if batch program has bug?

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# Core Memories (1950s & 60s)

The first magnetic core memory, from the IBM 405 Alphabetical Accounting Machine.

- Core Memory stored data as magnetization in iron rings
  - Iron "cores" woven into a 2-dimensional mesh of wires
  - Origin of the term "Dump Core"
  - Rumor that IBM consulted Life Saver company
- See: http://www.columbia.edu/acis/history/core.html

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#### History Phase $1\frac{1}{2}$ (late 60s/early 70s)

- · Data channels, Interrupts: overlap I/O and compute
  - DMA Direct Memory Access for I/O devices
  - I/O can be completed asynchronously
- Multiprogramming: several programs run simultaneously
  - Small jobs not delayed by large jobs
  - More overlap between I/O and CPU
  - Need memory protection between programs and/or OS
- Complexity gets out of hand:

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- Multics: announced in 1963, ran in 1969
  - » www.multicians.org lists 1777 people who "contributed to Multics". Probably 30-40 core developers.
  - » Turing award lecture from Fernando Corbató (key researcher): "On building systems that will fail"
- OS 360: released with 1000 known bugs (APARs)
   » "Anomalous Program Activity Report"
- OS finally becomes and important science:
  - How to deal with complexity???

-UNIX based on Multics but vastly simplified

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### A Multics System (Circa 1976)

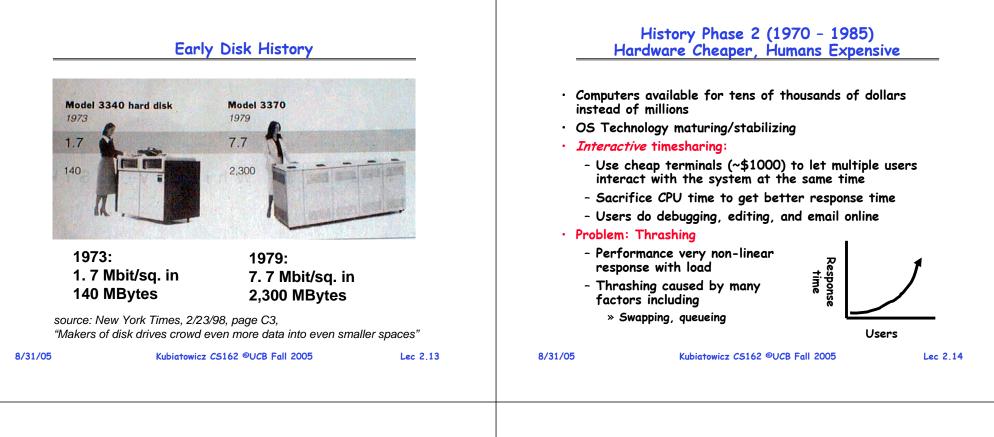


• The 6180 at MIT IPC, skin doors open, circa 1976:

- "We usually ran the machine with doors open so the operators could see the AQ register display, which gave you an idea of the machine load, and for convenient access to the EXECUTE button, which the operator would push to enter BOS if the machine crashed."

• http://www.multicians.org/multics-stories.html 8/31/05 Kubiatowicz CS162 @UCB Fall 2005

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#### Administriva: Time for Project Signup

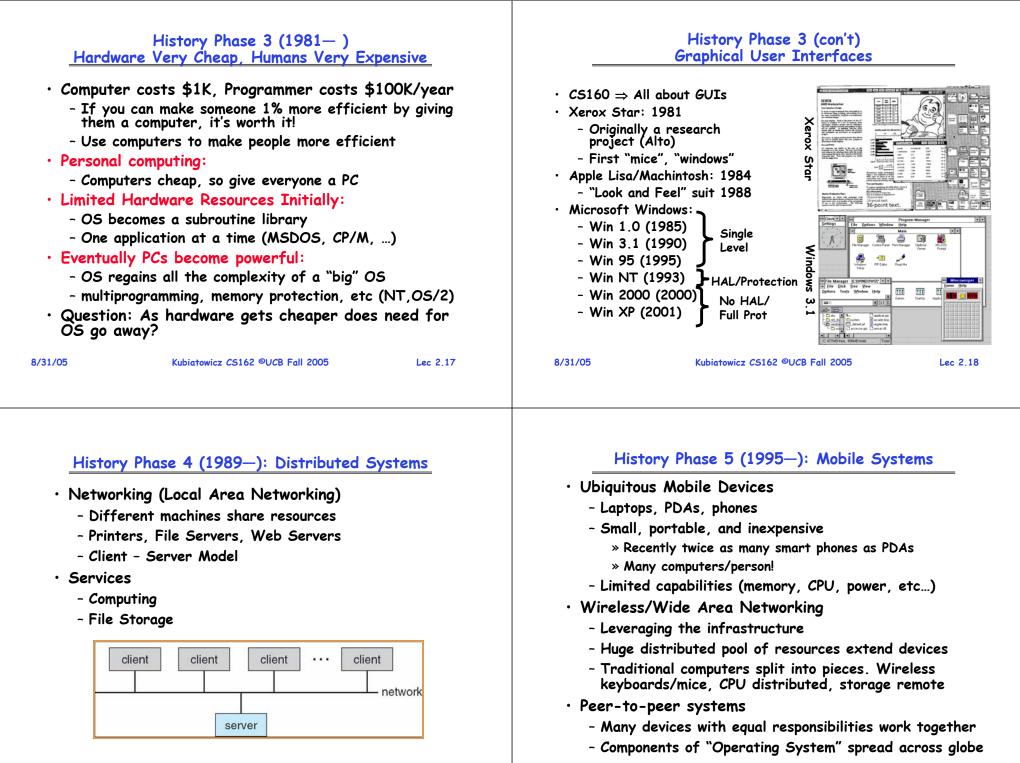
- Project Signup:
  - The group signup page is now working
  - Only submit once per group!
  - Everyone in group must have logged into their cs162-xx accounts before you register the group
  - Make sure that you select at least 2 potential sections
  - Due date: Wednesday 9/7 by 11:59pm
- Next week, go to your pre-assigned section

Section	Time	Location	ТА
101	Tu 1:00-2:00P	310 Hearst Mining	Dominic
102	W 10:00-11:00A	2 Evans	Rajesh
103	W 11:00-12:00P	85 Evans	Rajesh
104	W 1:00-2:00P	85 Evans	Chris
105	W 2:00-3:00P	85 Evans	Chris

#### Administrivia (2)

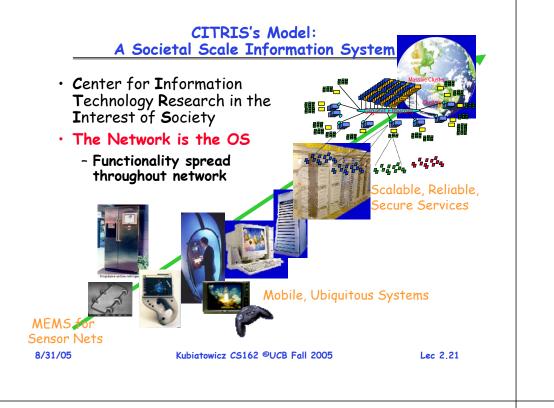
- Cs162-xx accounts:
  - Make sure you got an account form
  - If you haven't logged in yet, you need to do so
- Clarification of late policy:
  - I need to enforce the midnight deadlines
- Nachos readers:
  - Available Friday from Northside Copy Central
  - Includes printouts of all of the code
- Book Status:
  - Campus bookstore is out of the books new shipment scheduled next week
  - Ned's has 19 of the Silbershatz books and one used copy of the Free BSD book
- Web cast archives available off lectures page
  - Just click on the title of a lecture for webcast
  - Only works for lectures that I have already given!

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#### Moore's Law Reprise: My new X41 tablet

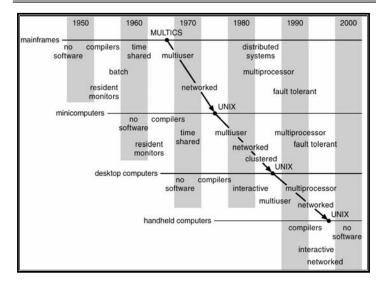
	1981	2005	2005 Laptop
CPU MHz,	10	3800	1500
Cycles/inst	3—10	0.25-0.5	0.25-0.5
DRAM capacity	128KB	4GB	1.5GB
Disk capacity	10MB	1TB	60GB
Net bandwidth	9600 b/s	1 Gb/s	1 Gb/s (wired) 54 Mb/s (wireless) 2 Mb/s (wide-area)
# addr bits	16	32	32
#users/machine	10s	≤ <b>1</b>	$\leq \frac{1}{4}$
Price	\$25,000	\$4,000	\$2000

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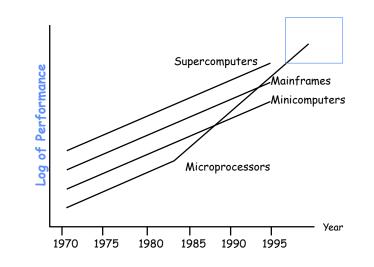
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#### Migration of Operating-System Concepts and Features



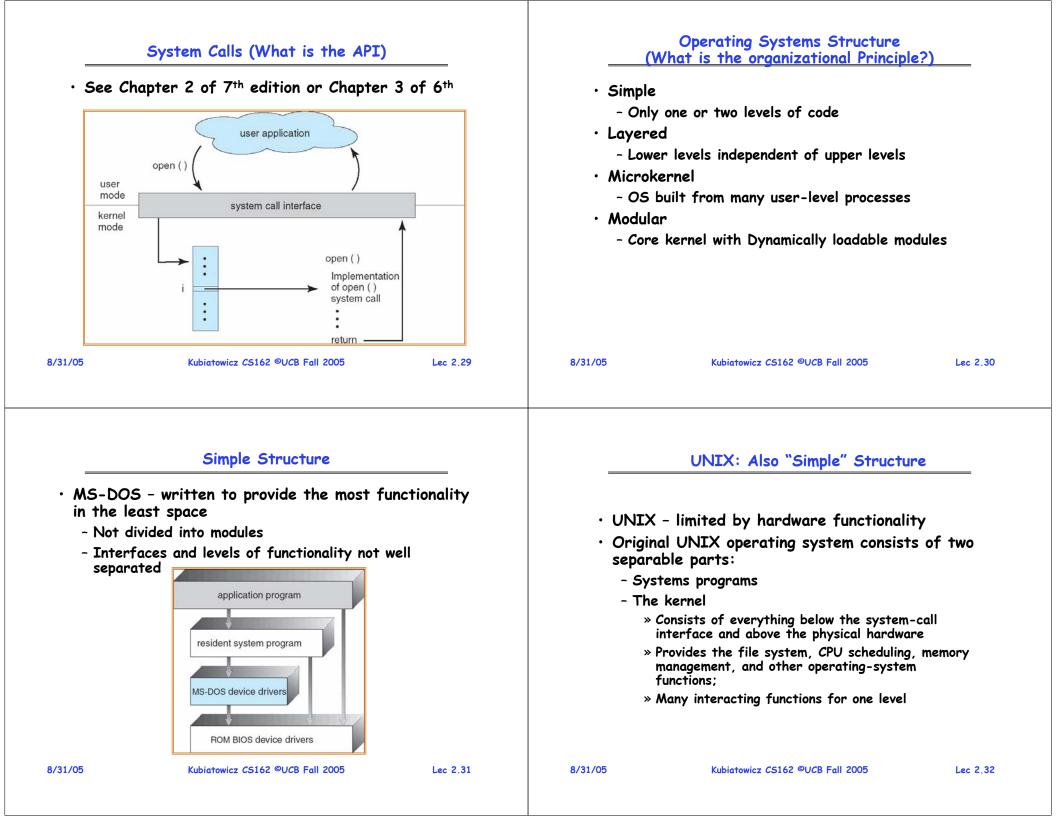
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#### Compare: Performance Trends (from CS152)



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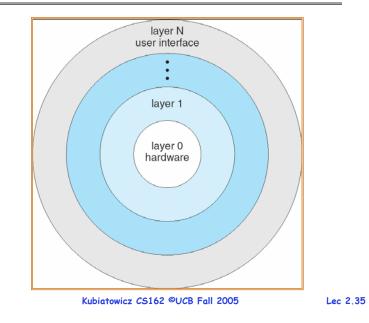
	History of OS: Summary				
<ul> <li>Change is continuous and OSs should adapt <ul> <li>Not: look how stupid batch processing was</li> <li>But: Made sense at the time</li> </ul> </li> <li>Situation today is much like the late 60s <ul> <li>Small OS: 100K lines</li> <li>Large OS: 10M lines (5M for the browser!)</li> <li>100-1000 people-years</li> </ul> </li> <li>Complexity still reigns <ul> <li>NT under development from early 90's to late 90's</li> <li>» Still doesn't work very well</li> <li>Jury still out on Windows 2000/XP</li> <li>Windows "Longhorn" delayed many times</li> <li>» Latest release date of 2006+</li> <li>» Promised by removing some of the intended technology</li> </ul> </li> </ul>			Now for a quick tour of OS Structures		
/31/05	Kubiatowicz CS162 ©UCB Fall 2005	Lec 2.25	8/31/05	Kubiatowicz CS162 ©UCB Fall 2005	Lec 2.26
(	Operating Systems Components What are the pieces of the OS	)	• Services ·	Operating System Services <u>(What things does the OS do?)</u> that (more-or-less) map onto com	
• Main-Mo	king		<ul> <li>» How</li> <li>- I/O ope</li> <li>» Stan</li> <li>- File sys</li> <li>» How</li> <li>» Loor</li> <li>- Commun</li> <li>» Netv</li> <li>• Cross-cut</li> <li>- Error d</li> </ul>	dardized interfaces to extremely divers tem manipulation do you read/write/preserve files? ning concern: How do you even find files nications vorking protocols/Interface with CyberS ting capabilities etection & recovery a allocation ting	e devices ???
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#### UNIX System Structure

User Mode		Applications         (the users)           Standard Libs         shells and commands compilers and interpreters system libraries			
Kernel Mode	Kernel	signals terminal handling character I/O system terminal drivers	n-call interface to the ke file system swapping block I/O system disk and tape drivers	CPU scheduling page replacement demand paging virtual memory	
Hardware		terminal controllers terminals	memory controllers physical memory		
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## Layered Operating System

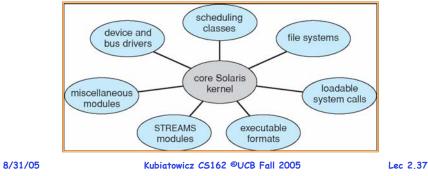


## Layered Structure

• Operating	system is divided many layers (le	evels)
•	ilt on top of lower layers	
	layer (layer 0) is hardware	
- Highest	layer (layer N) is the user interfac	e
<ul> <li>Each laye only lower</li> </ul>	r uses functions (operations) and -level layers	services of
	age: modularity ⇒ Easier debugging/ ays possible: Does process scheduler irtual memory layer?	
» Need	t to reschedule processor while waiting f need to page in information about tasks	or paging
•	: Machine-dependent vs independ	
•	nigration between platforms	•
	evolution of hardware platform	
- Good id	ea for you as well!	
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	Microkernel Structure	
· Moves as	much from the kernel into " <i>user</i> "	shaca
	ore OS running at kernel level	spuce
	vices built from many independent us	ser-level
<ul> <li>Communic</li> <li>Benefits:</li> </ul>	ation between modules with messo	ige passing
	to extend a microkernel	
_	to port OS to new architectures	
	eliable (less code is running in kernel	mode)
	solation (parts of kernel protected fi	-
- More se	ecure	
• Detriment	'S:	
- Perform	aance overhead severe for naïve impl	ementation
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#### Modules-based Structure

- Most modern operating systems implement modules
  - Uses object-oriented approach
  - Each core component is separate
  - Each talks to the others over known interfaces
  - Each is loadable as needed within the kernel
- $\cdot$  Overall, similar to layers but with more flexible



#### Operating System Design Goals (What is this OS trying to achieve?)

- \$2000 price point?
- Fault tolerance/Fast failover/High Availability?
- High Performance?
- Real Time Capable?

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#### Implementation Issues (How is the OS implemented?)

- Policy vs. Mechanism
  - Policy: What do you want to do?
  - Mechanism: How are you going to do it?
  - Should be separated, since policies change
- Algorithms used
  - Linear, Tree-based, Log Structured, etc...
- Event models used
  - threads vs event loops
- Backward compatability issues
  - Very important for Windows 2000
- System generation/configuration
  - How to make generic OS fit on specific hardware

#### Conclusion

- Rapid Change in Hardware Leads to changing OS
  - Batch ⇒ Multiprogramming ⇒ Timeshare ⇒ Graphical UI ⇒ Ubiquitous Devices ⇒ Cyberspace/Metaverse/??
- $\cdot$  OS features migrated from mainframes  $\Rightarrow$  PCs
- Standard Components and Services
  - Process Control
  - Main Memory
  - I/O
  - File System
  - UI
- Policy vs Mechanism
  - Crucial division: not always properly separated!
- · Complexity is always out of control
  - However, "Resistance is NOT Useless!"

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