

CS162 Operating Systems and Systems Programming Lecture 7

Mutual Exclusion, Semaphores, Monitors, and Condition Variables

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Review: A Concurrent Program Example

- Two threads, A and B, compete with each other
 - One tries to increment a shared counter
 - The other tries to decrement the counter

<u>Thread A</u>	<u>Thread B</u>
i = 0;	i = 0;
while (i < 10)	while (i > -10)
i = i + 1;	i = i - 1;
printf("A wins!");	printf("B wins!");

- Assume that memory loads and stores are atomic, but incrementing and decrementing are *not* atomic
- Who wins? Could be either
- Is it guaranteed that someone wins? Why or why not?
- What if both threads have their own CPU running at same speed? Is it guaranteed that it goes on forever?

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Review: Hand Simulating Multiprocessor Example

- Inner loop looks like this:

<u>Thread A</u>	<u>Thread B</u>
r1=0 load r1, M[i]	r1=0 load r1, M[i]
r1=1 add r1, r1, 1	r1=-1 sub r1, r1, 1
M[i]=1 store r1, M[i]	M[i]=-1 store r1, M[i]

- **Hand Simulation:**
 - And we're off. A gets off to an early start
 - B says "hmp, better go fast" and tries really hard
 - A goes ahead and writes "1"
 - B goes and writes "-1"
 - A says "HUH??? I could have sworn I put a 1 there"
- Could this happen on a uniprocessor?
 - Yes! Unlikely, but if you depending on it not happening, it will and your system will break.

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Review: Too Much Milk Solution #3

- Here is a possible two-note solution:

<u>Thread A</u>	<u>Thread B</u>
leave note A;	leave note B;
while (note B) {	if (noNote A) {
do nothing;	if (noMilk) {
}	buy milk;
if (noMilk) {	}
buy milk;	remove note B;
}	}
remove note A;	

- Does this work? Yes. Both can guarantee that:
 - It is safe to buy, or
 - Other will buy, ok to quit
- At A:
 - if no note B, safe for A to buy,
 - otherwise wait to find out what will happen
- At B:
 - if no note A, safe for B to buy
 - Otherwise, A is either buying or waiting for B to quit

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Goals for Today

- Hardware Support for Synchronization
- Higher-level Synchronization Abstractions
 - Semaphores, monitors, and condition variables
- Programming paradigms for concurrent programs



Note: Some slides and/or pictures in the following are adapted from slides ©2005 Silberschatz, Galvin, and Gagne

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High-Level Picture

- The abstraction of threads is good:
 - Maintains sequential execution model
 - Allows simple parallelism to overlap I/O and computation
- Unfortunately, still too complicated to access state shared between threads
 - Consider "too much milk" example
 - Implementing a concurrent program with only loads and stores would be tricky and error-prone
- Today, we'll implement higher-level operations on top of atomic operations provided by hardware
 - Develop a "synchronization toolbox"
 - Explore some common programming paradigms



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Where are we going with synchronization?

Programs	Shared Programs				
Higher-level API	Locks	Semaphores	Monitors	Send/Receive	
Hardware	Load/Store	Disable Ints	Test&Set	Comp&Swap	

- We are going to implement various higher-level synchronization primitives using atomic operations
 - Everything is pretty painful if only atomic primitives are load and store
 - Need to provide primitives useful at user-level

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How to implement Locks?

- **Lock:** prevents someone from doing something
 - Lock before entering critical section and before accessing shared data
 - Unlock when leaving, after accessing shared data
 - Wait if locked
 - » Important idea: all synchronization involves waiting
- Atomic Load/Store: get solution like Milk #3
 - Looked at this last lecture
 - Pretty complex and error prone
- Hardware Lock instruction
 - Is this a good idea?
 - Complexity?
 - » Done in the Intel 432.
 - » Each feature makes hardware more complex and slow
 - What about putting task to sleep?
 - » How do you handle the interface between the hardware and scheduler?



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Naïve use of Interrupt Enable/Disable

- How can we build multi-instruction atomic operations?
 - Recall: dispatcher gets control in two ways.
 - » Internal: Thread does something to relinquish the CPU
 - » External: Interrupts cause dispatcher to take CPU
 - On a uniprocessor, can avoid context-switching by:
 - » Avoiding internal events (although virtual memory tricky)
 - » Preventing external events by disabling interrupts
- Consequently, naïve Implementation of locks:

```
LockAcquire { disable Ints; }
LockRelease { enable Ints; }
```

- Problems with this approach:

- **Can't let user do this!** Consider following:

```
LockAcquire();
While(TRUE) {;
```

- Real-Time system—no guarantees on timing!
 - » Critical Sections might be arbitrarily long
- What happens with I/O or other important events?
 - » "Reactor about to meltdown. Help?"




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Better Implementation of Locks by Disabling Interrupts

- Key idea: maintain a lock variable and impose mutual exclusion only during operations on that variable

```
int value = FREE; 
```

```
Acquire() {
    disable interrupts;
    if (value == BUSY) {
        put thread on wait queue;
        Go to sleep();
        // Enable interrupts?
    } else {
        value = BUSY;
    }
    enable interrupts;
}

Release() {
    disable interrupts;
    if (anyone on wait queue) {
        take thread off wait queue;
        Place on ready queue;
    } else {
        value = FREE;
    }
    enable interrupts;
}
```

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New Lock Implementation: Discussion

- Why do we need to disable interrupts at all?
 - Avoid interruption between checking and setting lock value
 - Otherwise two threads could think that they both have lock

```
Acquire() {
    disable interrupts;
    if (value == BUSY) {
        put thread on wait queue;
        Go to sleep();
        // Enable interrupts?
    } else {
        value = BUSY;
    }
    enable interrupts;
}
```

} Critical Section

- Note: unlike previous solution, the critical section (inside Acquire()) is very short
 - User of lock can take as long as they like in their own critical section: doesn't impact global machine behavior
 - Critical interrupts taken in time!

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Interrupt re-enable in going to sleep

- What about re-enabling ints when going to sleep?

```
Acquire() {
    disable interrupts;
    if (value == BUSY) {
        put thread on wait queue;
        Go to sleep();
    } else {
        value = BUSY;
    }
    enable interrupts;
}
```

Enable Position
Enable Position
Enable Position

→ → →

- Before Putting thread on the wait queue?
 - Release can check the queue and not wake up thread
- After putting the thread on the wait queue
 - Release puts the thread on the ready queue, but the thread still thinks it needs to go to sleep
 - Misses wakeup and still holds lock (deadlock!)
- Want to put it after sleep(). But - how?

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Administrivia

- First Design Document due Monday 9/26
 - Subsequently need to schedule design review with TA (through web form)
 - Note that Much of the design document grade comes from first version (some from final version)
- CVS group accounts should be setup
 - Check out the CVS Quick Start Guide for instructions on how to get your CVS repository working
 - If you change your key - need to let us know!
- Anyone interested in being a note-taker?
 - Have a student who needs help with note taking
 - Can receive payment for this help

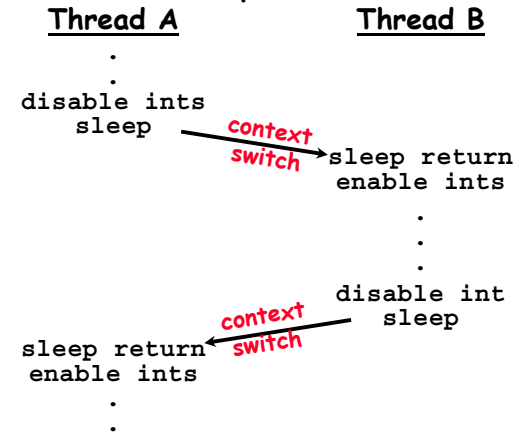
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How to Re-enable After Sleep()?

- In Nachos, since ints are disabled when you call sleep:
 - Responsibility of the next thread to re-enable ints
 - When the sleeping thread wakes up, returns to acquire and re-enables interrupts



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Interrupt disable and enable across context switches

- An important point about structuring code:
 - In Nachos code you will see lots of comments about assumptions made concerning when interrupts disabled
 - This is an example of where modifications to and assumptions about program state can't be localized within a small body of code
 - In these cases it is possible for your program to eventually "acquire" bugs as people modify code
- Other cases where this will be a concern?
 - What about exceptions that occur after lock is acquired? Who releases the lock?

```
mylock.acquire();  
a = b / 0;  
mylock.release();
```

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Atomic Read-Modify-Write instructions

- Problems with previous solution:
 - Can't give lock implementation to users
 - Doesn't work well on multiprocessor
 - » Disabling interrupts on all processors requires messages and would be very time consuming
- Alternative: atomic instruction sequences
 - These instructions read a value from memory and write a new value atomically
 - Hardware is responsible for implementing this correctly on both uniprocessors (not too hard) and multiprocessors (requires help from cache coherence protocol)
 - Unlike disabling interrupts, can be used on both uniprocessors and multiprocessors

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Examples of Read-Modify-Write

```
• test&set (&address) { /* most architectures */
    result = M[address];
    M[address] = 1;
}
• swap (&address, register) { /* x86 */
    temp = M[address];
    M[address] = register;
    register = temp;
}
• compare&swap (&address, reg1, reg2) { /* 68000 */
    if (reg1 == M[address]) {
        M[address] = reg2;
        return success;
    } else {
        return failure;
    }
}
• load-linked&store conditional(&address) {
    /* R4000, alpha */
    loop:
        ll r1, M[address];
        movi r2, 1; /* Can do arbitrary comp */
        sc r2, M[address];
        beqz r2, loop;
}
```

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Implementing Locks with test&set

- Another flawed, but simple solution:

```
int value = 0; // Free
Acquire() {
    while (test&set(value)); // while busy
}
Release() {
    value = 0;
}
```

- Simple explanation:

- If lock is free, test&set reads 0 and sets value=1, so lock is now busy. It returns 0 so while exits.
- If lock is busy, test&set reads 1 and sets value=1 (no change). It returns 1, so while loop continues
- When we set value = 0, someone else can get lock

- **Busy-Waiting:** thread consumes cycles while waiting

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Problem: Busy-Waiting for Lock

- Positives for this solution
 - Machine can receive interrupts
 - User code can use this lock
 - Works on a multiprocessor
- Negatives
 - This is very inefficient because the busy-waiting thread will consume cycles waiting
 - Waiting thread may take cycles away from thread holding lock (no one wins!)
 - **Priority Inversion:** If busy-waiting thread has higher priority than thread holding lock \Rightarrow no progress!
- Priority Inversion problem with original Martian rover
- For semaphores and monitors, waiting thread may wait for an arbitrary length of time!
 - Thus even if busy-waiting was OK for locks, definitely not ok for other primitives
 - Homework/exam solutions should not have busy-waiting!



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Better Locks using test&set

- Can we build test&set locks without busy-waiting?
 - Can't entirely, but can minimize!
 - Idea: only busy-wait to atomically check lock value

```
int guard = 0;
int value = FREE;
```



```
Acquire() {
    // Short busy-wait time
    while (test&set(guard));
    if (value == BUSY) {
        put thread on wait queue;
        go to sleep() & guard = 0;
    } else {
        value = BUSY;
        guard = 0;
    }
}
Release() {
    // Short busy-wait time
    while (test&set(guard));
    if anyone on wait queue {
        take thread off wait queue
        Place on ready queue;
    } else {
        value = FREE;
        guard = 0;
    }
}
```

- Note: sleep has to be sure to reset the guard variable
 - Why can't we do it just before or just after the sleep?

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Higher-level Primitives than Locks

- Goal of last couple of lectures:
 - What is the right abstraction for synchronizing threads that share memory?
 - Want as high a level primitive as possible
- Good primitives and practices important!
 - Since execution is not entirely sequential, really hard to find bugs, since they happen rarely
 - UNIX is pretty stable now, but up until about mid-80s (10 years after started), systems running UNIX would crash every week or so - concurrency bugs
- Synchronization is a way of coordinating multiple concurrent activities that are using share state
 - This lecture and the next presents a couple of ways of structuring the sharing

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Semaphores



- Semaphores are a kind of generalized lock
 - First defined by Dijkstra in late 60s
 - Main synchronization primitive used in original UNIX
- Definition: a Semaphore has a non-negative integer value and supports the following two operations:
 - **P()**: an atomic operation that waits for semaphore to become positive, then decrements it by 1
 - » Think of this as the wait() operation
 - **V()**: an atomic operation that increments the semaphore by 1, waking up a waiting P, if any
 - » Think of this as the signal() operation
 - Note that **P()** stands for "*proberen*" (to test) and V stands for "*verhogen*" (to increment) in Dutch

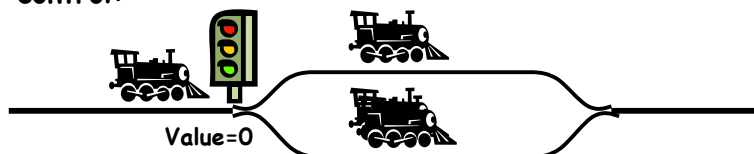
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Semaphores Like Integers Except

- Semaphores are like integers, except
 - No negative values
 - Only operations allowed are P and V - can't read or write value, except to set it initially
 - Operations must be atomic
 - » Two P's together can't decrement value below zero
 - » Similarly, thread going to sleep in P won't miss wakeup from V - even if they both happen at same time
- Semaphore from railway analogy
 - Here is a semaphore initialized to 2 for resource control:



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Two Uses of Semaphores

- Mutual Exclusion (initial value = 1)
 - Also called "Binary Semaphore".
 - Can be used for mutual exclusion:

```
semaphore.P();
// Critical section goes here
semaphore.V();
```
- Scheduling Constraints (initial value = 0)
 - Locks are fine for mutual exclusion, but what if you want a thread to wait for something?
 - Example: suppose you had to implement ThreadJoin which must wait for thread to terminate:

```
Initial value of semaphore = 0
ThreadJoin {
    semaphore.P();
}
ThreadFinish {
    semaphore.V();
}
```

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Producer-consumer with a bounded buffer



- Problem Definition
 - Producer puts things into a shared buffer
 - Consumer takes them out
 - Need synchronization to coordinate producer/consumer
- Don't want producer and consumer to have to work in lockstep, so put a fixed-size buffer between them
 - Need to synchronize access to this buffer
 - Producer needs to wait if buffer is full
 - Consumer needs to wait if buffer is empty
- Example 1: GCC compiler
 - `cpp | cc1 | cc2 | as | ld`
- Example 2: Coke machine
 - Producer can put limited number of cokes in machine
 - Consumer can't take cokes out if machine is empty



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Correctness constraints for solution

- Correctness Constraints:
 - Consumer must wait for producer to fill buffers, if none full (scheduling constraint)
 - Producer must wait for consumer to empty buffers, if all full (scheduling constraint)
 - Only one thread can manipulate buffer queue at a time (mutual exclusion)
- Remember why we need mutual exclusion
 - Because computers are stupid
 - Imagine if in real life: the delivery person is filling the machine and somebody comes up and tries to stick their money into the machine
- General rule of thumb: **Use a separate semaphore for each constraint**
 - Semaphore `fullBuffers`; // consumer's constraint
 - Semaphore `emptyBuffers`; // producer's constraint
 - Semaphore `mutex`; // mutual exclusion

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Full Solution to Bounded Buffer

```
Semaphore fullBuffer = 0; // Initially, no coke
Semaphore emptyBuffers = numBuffers;
                               // Initially, num empty slots
Semaphore mutex = 1; // No one using machine

Producer(item) {
    emptyBuffers.P(); // Wait until space
    mutex.P(); // Wait until buffer free
    Enqueue(item);
    mutex.V();
    fullBuffers.V(); // Tell consumers there is
                    // more coke
}

Consumer() {
    fullBuffers.P(); // Check if there's a coke
    mutex.P(); // Wait until machine free
    item = Dequeue();
    mutex.V();
    emptyBuffers.V(); // tell producer need more
    return item;
}
```

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Discussion about Solution

- Why asymmetry?
 - Producer does: `emptyBuffer.P()`, `fullBuffer.V()`
 - Consumer does: `fullBuffer.P()`, `emptyBuffer.V()`
- Is order of P's important?
 - Yes! Can cause deadlock
- Is order of V's important?
 - No, except that it might affect scheduling efficiency
- What if we have 2 producers or 2 consumers?
 - Do we need to change anything?

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Motivation for Monitors and Condition Variables

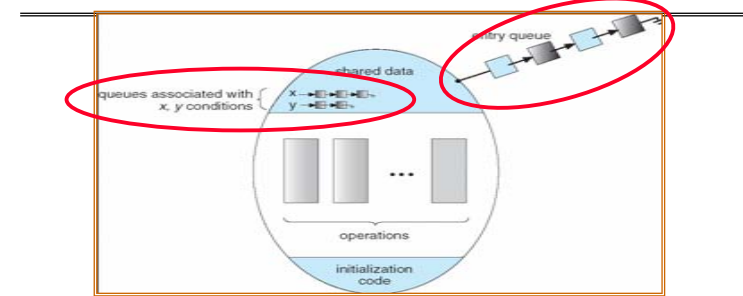
- Semaphores are a huge step up; just think of trying to do the bounded buffer with only loads and stores
 - Problem is that semaphores are dual purpose:
 - » They are used for both mutex and scheduling constraints
 - » Example: the fact that flipping of P's in bounded buffer gives deadlock is not immediately obvious. How do you prove correctness to someone?
- Cleaner idea: Use *locks* for mutual exclusion and *condition variables* for scheduling constraints
- Definition: **Monitor**: a lock and zero or more condition variables for managing concurrent access to shared data
 - Some languages like Java provide this natively
 - Most others use actual locks and condition variables

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Monitor with Condition Variables



- **Lock**: the lock provides mutual exclusion to shared data
 - Always acquire before accessing shared data structure
 - Always release after finishing with shared data
 - Lock initially free
- **Condition Variable**: a queue of threads waiting for something *inside* a critical section
 - Key idea: make it possible to go to sleep inside critical section by atomically releasing lock at time we go to sleep
 - Contrast to semaphores: **Can't wait inside critical section**

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Simple Monitor Example

- Here is an (infinite) synchronized queue

```
Lock lock;
Condition dataready;
Queue queue;

AddToQueue(item) {
    lock.Acquire();           // Get Lock
    queue.enqueue(item);     // Add item
    dataready.signal();     // Signal any waiters
    lock.Release();         // Release Lock
}

RemoveFromQueue() {
    lock.Acquire();           // Get Lock
    while (queue.isEmpty()) {
        dataready.wait(&lock); // If nothing, sleep
    }
    item = queue.dequeue();  // Get next item
    lock.Release();         // Release Lock
    return(item);
}
```

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Summary

- Important concept: Atomic Operations
 - An operation that runs to completion or not at all
 - These are the primitives on which to construct various synchronization primitives
- Talked about hardware atomicity primitives:
 - Disabling of Interrupts, test&set, swap, comp&swap, load-locked/store conditional
- Showed several constructions of Locks
 - Must be very careful not to waste/tie up machine resources
 - » Shouldn't disable interrupts for long
 - » Shouldn't spin wait for long
 - Key idea: Separate lock variable, use hardware mechanisms to protect modifications of that variable
- Talked about Semaphores, Monitors, and Condition Variables
 - Higher level constructs that are harder to "screw up"

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