# CS 162 Operating Systems and Systems Programming Professor: Anthony D. Joseph Spring 2003

# **Lecture 25: Course Review**

#### 25.0 Course Goals

- 1. Provide you with the knowledge you need to make informed decisions.
  - Is it better to buy a computer with more memory or a faster processor (or faster memory versus faster processor)?
  - Why is my company's Web server slow? Is it the network, the server, the application?
- 2. Experience with different design tradeoffs, choices, and decisions.
  - What is the cost of using a software modem instead of a software modem? Everything that's done in hardware can be done in software, but when does it make sense?
  - How do I enable my company's users to share information with collaborators at other companies? With good performance. Without compromising security.
- 3. Design abstractions: separating policy from mechanism
  - What abstractions should the operating system provide?
  - How should I implement privacy controls?

#### 25.1 OS as Illusionist

We used the Operating System as starting point for understanding/analyzing the issues.

Operating systems have two functions:

- 1. Coordinator and traffic cop
- 2. Standard services

Physical Reality	Abstraction
Single CPU	Infinite # of CPUs
	(multiprogramming)
Interrupts	Cooperating sequential threads

Limited memory	Unlimited virtual memory
No protection	Each address space has its own
	machine
Unreliable, fixed size messages	Reliable, arbitrary messages and
	network services

## 25.2 Concepts

We've abstracted out three key concepts. They apply to more than just operating systems

- Locality/Caching basis for TLB's, paging, file systems, distributed systems, etc.
  - · Spatial versus temporal locality
  - Thrashing
  - · Multi-level hierarchies
  - · Same issue in HW and SW
- 2. Scheduling adaptive management of resources
  - · Constrained resources require careful management
  - · Multi-level adaptive feedback
  - · Countermeasures for misbehaving users and applications
- 3. Layering Abstraction on top of abstraction
  - Use divide and conquer to simplify a hard problem.
  - Makes it easier to design, debug, extend
  - · Performance penalty

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## 25.3 Major topics

- 1. Threads: state, creation, dispatching
  - Why Abstraction for concurrency: overlap I/O and computation, share HW resources (and information) across multiple users and programs. Modularity makes system easier to extend.
  - How Context switching, and thread dispatching (mechanism).and scheduling. Decompose task into smaller units/functions.
  - But performance overhead for context switching.
- Synchronization: races, inconsistency, semaphores, monitors, and condition variables.
  - The cost of concurrency! Without sharing concurrency is useless, but remember the "Too Much Milk Lecture"
  - Non-reproducibility Hard to debug!
  - Use atomic operations as a start, but complicated to use and OS interactions (load/store, interrupt disable, test&set).
  - Create higher level abstractions to ease the burden:
    - o Critical sections and mutual exclusion policy.
    - o Locks and semaphores mechanism.
    - Monitors: separate mutex (locks) and scheduling constraints (condition variables) – mechanisms.
  - Language-level interactions with primitives. Be careful!
  - · Biggest caveats: Deadlock and starvation
    - Starvation: Indefinite waiting for a resource by a thread (can end, but doesn't have to).
    - Deadlock: Circular chain of waiting (doesn't end without external intervention). Requires: limited resource, no resource preemption, multiple independent requests, circular chain of requests. Break the chain – detect/fix or prevent
- 3. Scheduling: shortest (remaining) time to completion first, round robin, FIFO
  - · Policy: minimize response time, maximize throughput, fair.
  - Lots of choices: algorithm, time slice, dynamic adaptation (multi-level feedback), etc. – most choices don't really matter unless resources are constrained
- 4. Memory management & address spaces:
  - Isolate processes/programs from all others and OS- protection.

- Dual mode operation: kernel versus user mode operations themselves must be protected: How do you enter/leave kernel mode?
- o This can be done without hardware support:
  - Strong typing
  - Software fault isolation
- o But inter-process communication breaks this (bugs can leak).
- Illusion of infinite memory:
  - o Build a hierarchy out of fast, small -> large, small technologies
- Transparent (can't tell if physical memory is shared)
  - o Address translation
  - Base & bounds, paging, segmentation, multi-level translation, TLB's for caching/performance (replacement policy and write-back/write-through are considerations – thrashing).
  - o Complexity versus functionality tradeoffs
- 5. Virtual memory: demand paging, thrashing
  - · Exploit spatial and temporal locality
  - Caching misses: compulsory, capacity, conflict, policy
  - Lots of page replacement policies: Again, most important when resources are limited! Approximations work well.
  - · Application working set size is important
- 6. File systems:
  - I/O system performance: overhead, latency, bandwidth
    - o Disk seeks, rotational delay, sector sizes
    - o Scheduling is important: FIFO, elevator (SCAN)
  - File headers and directories: abstraction of bytes, named files, protection, durability
    - o Management policies based upon file usage patterns
    - Caching for performance
    - o Protection and access control are important
    - o Transactions: Implement atomic, persistent operations (durability) for unreliable components.
      - Two-phase locking for coordinating multiple threads
- 7. Distributed computing
  - · Cheaper, more reliable, incremental scalability

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- In reality, not more reliable
- Coordination is more difficult than in single machine case.
- 8. Networks: protocol layers, windowing, RPC
  - Build protocols layer-by-layer
  - · Lots of different network technologies
  - Goals: arbitrary message size, ordered, reliable, process-to-process, routed anywhere, secure
  - Goals are hard (lots can go wrong)
  - Remote Procedure Call is key abstraction for 2-way communication:
    - o Cross-domain communication
    - o Location-transparency
    - o Microkernel is ultimate in RPC usage
- 9. Network file systems: cache coherence
  - Transparent access to files on a remote disk: NFS, AFS
  - · Caching, consistency, and false sharing issues
  - Multiprocessors: shared-bus, switched, Network of Workstations (Similar problems to filesystems)
- 10. Security: access control, encryption, Trojan horses
  - Why you should never trust a computer!
  - · Intentional and accidental misuse
  - · Three parts:
    - o Authentication who user is
      - Passwords, encryption (private and public key encryption)
    - o Authorization who is allowed to do what
      - Access control lists
    - Enforcement make sure people do what they're supposed to do
      - Kernel does this for OS

## 25.4 Problem Areas

1. Performance

Abstractions like threads, RPC aren't free

Remember threads in OS/2

Caching doesn't work when there's little locality

- Failures how do we build systems that continue to work even when parts of the system break?
   Still a problem today!
- Security basic tradeoff between making computer systems easy to use vs. hard to misuse

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