# University of California, Berkeley 

College of Engineering
Computer Science Division - EECS
Fall 2011
Anthony D. Joseph and Ion Stoica

## Midterm Exam

October 13, 2011
CS162 Operating Systems

| Your Name: |  |
| :--- | :--- |
| SID AND 162 Login: |  |
| TA Name: |  |
| Discussion Section <br> Time: |  |

General Information:
This is a closed book and one 2-sided note examination. You have 80 minutes to answer as many questions as possible. The number in parentheses at the beginning of each question indicates the number of points for that question. You should read all of the questions before starting the exam, as some of the questions are substantially more time consuming.

Write all of your answers directly on this paper. Make your answers as concise as possible. If there is something in a question that you believe is open to interpretation, then please ask us about it!

Good Luck!!

| QUESTION | POINTS ASSIGNED | POINTS OBTAINED |
| :---: | :---: | :---: |
| 1 | 12 |  |
| 2 | 42 |  |
| 3 | 18 |  |
| 4 | 28 |  |
| TOTAL | 100 |  |

1. (12 points total) Short answer questions:
a. (4 points) List the four requirements for deadlock:

We gave one point for each requirement for deadlock.
b. (3 points) Operating System terms.
i) (1 point) List a synchronization mechanism that is an abstraction of a counting number (non-negative integer).
1 point for correctly described term; in very few cases we gave partial credit.
ii) (1 point) What is the portion of a process' code that manipulates shared variables?
1 point for correctly described term; in very few cases we gave partial credit.
iii) (1 point) What is the OS data structure that represents a running process? 1 point for correctly described term; in very few cases we gave partial credit.
c. ( 5 points) What are the differences between a Hoare monitor and a Mesa monitor?

Your answer should be brief (no more than 3 sentences).
For Hoare Monitors we gave 1 point for transfer lock, 1 point for transfer cpu, and 1 point to thread being woken.

For Mesa Monitors, we gave 1 point for putting the thread into ready queue, and 1 point for rechecking the condition and content for resources (lpt) We subtracted lpt for mentioning that the lock and cpu are not transferred for Mesa monitors.
2. (42 points total) CPU Scheduling.
a. (4 points) Provide a brief definition of what are preemptive and non-preemptive scheduling policies. Your answer should not exceed four sentences.

We gave 2 points for defining preemptive scheduling (i.e., "scheduler can take away CPU from the job before the CPU burst is over), and 2 points for defining the nonpreemptive scheduling (i.e., "jobs run until they explicitly yield the resource").

We took away 1 point for each obviously incorrect statement, e.g. "preemptive scheduling tries to minimize job completion time".
b. (3 points) Why is SJF or SRTF scheduling difficult to implement in a real OS?

Give an alternative approach with the same goals that can be implemented easily in a real OS

We gave points for saying "it's hard to predict the future", and 1 point for giving Multilevel Feedback Queue as an example. We didn't accept lottery scheduling.
c. (35 points) Here is a table of processes and their associated arrival and running times.

| Process ID | Arrival Time | Expected CPU <br> Running Time |
| :--- | :--- | :--- |
| Process 1 | 0 | 4 |
| Process 2 | 2 | 5 |
| Process 3 | 3 | 3 |
| Process 4 | 8 | 4 |

i) ( 15 points) Show the scheduling order for these processes under First-In-FirstOut (FIFO), Shortest-Job First (SJF), and Round-Robin (RR) with a quantum $=1$ time unit. Assume that the context switch overhead is 0 and new processes are added to the head of the queue except for FIFO.

| Time | FIFO | SJF | RR |
| :--- | :--- | :--- | :--- |
| $\mathbf{0}$ |  |  |  |
| $\mathbf{1}$ |  |  |  |
| 2 |  |  |  |
| 3 |  |  |  |
| 4 |  |  |  |
| 5 |  |  |  |
| 7 |  |  |  |
| $\mathbf{7}$ |  |  |  |
| $\mathbf{8}$ |  |  |  |
| 11 |  |  |  |
| 15 |  |  |  |
| 13 |  |  |  |
| 15 |  |  |  |

ii) (18 points) For each process in each schedule above, indicate the queue wait time and turnaround time (TRT).

The queue wait time is the total time a thread spends in the wait queue.
The turnaround time is defined as the time a process takes to complete after it arrives.

| Scheduler | Process 1 | Process 2 | Process 3 | Process 4 |
| :--- | :--- | :--- | :--- | :--- |
| FIFO queue wait |  |  |  |  |
| FIFO TRT |  |  |  |  |
| SJF queue wait |  |  |  |  |
| SJF TRT |  |  |  |  |
| RR queue wait |  |  |  |  |
| RR TRT |  |  |  |  |

iii) (2 points) Which of these three scheduling disciplines has the lowest average TRT?

# No Credit - Problem X (000000000000 points) 

## The 2010 Stella Awards

The annual 'Stella Awards' are named after 81-year-old Stella Liebeck who spilled hot coffee on herself and successfully sued the McDonald's in New Mexico, where she purchased coffee. You remember, she took the lid off the coffee and put it between her knees while she was driving. Who would ever think one could get burned doing that, right?

## *SEVENTH PLACE*

Kathleen Robertson of Austin, Texas was awarded $\$ 80,000$ by a jury of her peers after breaking her ankle tripping over a toddler who was running inside a furniture store. The store owners were understandably surprised by the verdict, considering the running toddler was her own son.

## *SIXTH PLACE *

Carl Truman, 19, of Los Angeles, California won $\$ 74,000$ plus medical expenses when his neighbor ran over his hand with a Honda Accord. Truman apparently didn't notice there was someone at the wheel of the car when he was trying to steal his neighbor's hubcaps.

## * FIFTH PLACE *

Terrence Dickson, of Bristol, Pennsylvania, who was leaving a house he had just burglarized by way of the garage. Unfortunately for Dickson, the automatic garage door opener malfunctioned and he could not get the garage door to open. Worse, he couldn't re-enter the house because the door connecting the garage to the house locked when Dickson pulled it shut. Forced to sit for eight, count 'em, EIGHT days and survive on a case of Pepsi and a large bag of dry dog food, he sued the homeowner's insurance company claiming undue mental Anguish. Amazingly, the jury said the insurance company must pay Dickson $\$ 500,000$ for his anguish.

## *FOURTH PLACE*

Jerry Williams, of Little Rock, Arkansas, garnered 4th Place in the Stella's when he was awarded \$14,500 plus medical expenses after being bitten on the butt by his next door neighbor's beagle - even though the beagle was on a chain in its owner's fenced yard. Williams did not get as much as he asked for because the jury believed the beagle might have been provoked at the time of the butt bite because Williams had climbed over the fence into the yard and repeatedly shot the dog with a pellet gun.

## *THIRD PLACE *

Amber Carson of Lancaster, Pennsylvania because a jury ordered a Philadelphia restaurant to pay her $\$ 113,500$ after she slipped on a spilled soft drink and broke her tailbone. The reason the soft drink was on the floor: Ms.Carson had thrown it at her boyfriend 30 seconds earlier during an argument.

## *SECOND PLACE*

Kara Walton, of Claymont, Delaware sued the owner of a night club in a nearby city because she fell from the bathroom window to the floor, knocking out her two front teeth. Even though Ms. Walton was trying to sneak through the ladies room window to avoid paying the $\$ 3.50$ cover charge, the jury said the night club had to pay her $\$ 12,000 \ldots$. oh, yeah, plus dental expenses.

## *FIRST PLACE *

This year's runaway First Place Stella Award winner was: Mrs. Merv Grazinski, of Oklahoma City, Oklahoma, who purchased new 32 -foot Winnebago motor home. On her first trip home, from an OU football game, having driven on to the freeway, she set the cruise control at 70 mph and calmly left the driver's seat to go to the back of the Winnebago to make herself a sandwich. Not surprisingly, the motor home left the freeway, crashed and overturned. Also not surprisingly, Mrs. Grazinski sued Winnebago for not putting in the owner's manual that she couldn't actually leave the driver's seat while the cruise control was set.. The Oklahoma jury awarded her, $\$ 1,750,000$ PLUS a new motor home. Winnebago actually
changed their manuals as a result of this suit, just in case Mrs. Grazinski has any relatives who might also buy a motor home.
3. (18 points) Networking.
a. (5 points) In a network, both the link layer and the transport layer may provide error detection facilities. Give 2 reasons that a transport-level error detection mechanism is needed to ensure end-to-end data integrity.
We gave 2.5 points per reason Two reasons that were basically the same idea counted as one reason. Additional acceptable answers (+2.5 each):

- Data could be corrupted, lost, etc. at higher layers than link layer
- Can't check for reordering of packets with some explanation of why this is true (packets can take different links, etc.)
- Enforcing at the link layer would impose unnecessary burden on applications that don't need reliability
Only endpoint can guarantee that the endpoint received a packet Partially acceptable answers (+1 each, more or less depending on the case):
- Explanation of E2E principle without tying it back to the actual issue
- Talking about reordering of packets, but inadequate explanation of what could go wrong with that
- Mentioning that endpoints would "have to check anyway" but no explanation of why
Unacceptable answers (+0 each):
- Just stating that packets can be dropped / corrupted / reordered / etc. in general
- Explanations of what reliability means
- Very general statements about network architecture
- Link layer "doesn't have enough information" with no / inadequate additional explanation
- Saying that reliability is the transport layer's responsibility, or that it's just the way the network is made
- Comparing transport layer vs. application layer (instead of link layer)
- Not mentioning anything about the link layer at all-Saying that the transport layer is in charge of routing
- Saying that the transport layer can ask for packets to be resent, with no additional explanation
b. (8 points) Statistical Multiplexing.
i) (5 points) Briefly explain the concept of statistical multiplexing. Your answer should not exceed three sentences.

We gave 2 points for the idea of multiple flows/flows with different peaks or burst times, and 3 points for the idea that as the number of flows increases, the peak to the average ratio decreases.
ii) (3 points) Briefly explain why statistical multiplexing is helpful for network designers. Your answer should not exceed two sentences.
We gave one point for the idea that multiplexing many flows "reduces burstiness"/smooths out peaks, and 2 points for the idea that this means that network designers need to provision less bandwidth, i.e., slightly over the average, rather than the peak.
c. (5 points) Flow control. Assume a sender sending 1000 byte packets to a receiver with a window size of 10 packets. Assume the round-trip time (RTT) is 50 ms . What is the maximum throughput (i.e., bits/sec) that the sender can sustain? (Ignore the packet header in your answer.)
points for correct expression and answer; answers with correct expression and correct unit conversion but with minor arithmetic errors also received full credit.for failing to account for receiver window, 3 points for incorrect RTT, and 1 point each for unit conversion errors (msec to sec, bytes to bits etc).
4. (28 points total) Concurrency control.
a. (12 points total) Dining Philosophers.

The goal of this exercise is to implement a solution to the Dining Philosophers problem - 5 philosophers sitting at a round table. Philosophers repeat (forever) the following three things in order: (1) think, (2) eat, and (3) sleep. Each philosopher has a plate of spaghetti in front of him or her and a chopstick between each plate.

Important rule: A philosopher must have two (2) chopsticks to be able to eat! A philosopher must acquire the chopstick to the left of their plate and the chopstick to the right of their plate; they cannot reach across the table to use other chopsticks. When they are done eating, they put down (and free) the two chopsticks for someone else to use. Multiple philosophers should be able to acquire chopsticks at the same time.
Below is code that attempts to solve this problem. Each philosopher will call this method with the parameter ( 0 to 4 ) indicating which philosopher they are. Unfortunately, this code can get stuck (deadlock).
You need to fix this code so it cannot deadlock and so that more than one philosopher can eat at the same time.

The constructor Philosopher () is called once to create the 5 chopsticks and 5 philosopher threads are created, each one executing the Dine () method.

Think(), Eat (), and Sleep () are methods provided to you (already written). You do not know how long each of these takes to complete.

## Your solution must avoid deadlock.

```
    Philosopher () {
        for (int i=0; i<=4; i++)
            chopstick[i] = new Semaphore(1);
        for (int j=0; j<=4; j++)
            thread_fork(Dine(j));
    }
    void Dine(int p) {
        while (1) {
            Think();
            chopstick[p].P();
            chopstick[(p + 1) % 5].P();
            Eat();
            chopstick[(p + 1) % 5].V();
            chopstick[p].v();
            Sleep();
    }
}
```

i) (4 points) Specify the correctness constraints. Be succinct and explicit in your answer.
ii) (8 points) Correctly re-implement the Dine () method.

We subtracted six points if your solution could still deadlock and 1 point if your solution was missing think/eat/sleep. We took four points off if your solution broke or violated the correctness constraints
b. (10 points) Synchronization. The following Java code implements three operations to handle a linked list: (a) insertion at the head of the list, (b) deletion from the head of the list, and (c) counting and printing the number of the items in the list.

```
void insertHead(int value) \{
        item = new Item(value);
        if (head == null) \{
                head = item;
        \} else \{
                item->next = head;
                head = item;
        \}
    \}
    item deleteHead(int value) \{
        item result = head;
        if (head != null) \{
            head = head->next;
        \}
        return result;
    \}
```

```
int getLen() {
    int len = 0;
    item tmp = head;
    while (tmp) {
        len = len + l;
        tmp = tmp->next;
}
    printf(len);
    return len;
}
```

Initially, the list contain the following three elements:


Assume there are three threads, each of them performing the following operations. (Note that the code is not thread safe. Assume that only reads and writes from/to memory are atomic.)

| Thread1 | Thread2 | Thread3 |
| :---: | :---: | :---: |
| insertHead(2) | deleteHead() | getLen() |

i) (6 points) What are the possible contents of the lists after the first two threads finish?

We subtracted 1 point for every missing outcome, and one point for every wrong one.
ii) (4 points) What are the possible outputs printed by Thread 3?

We subtracted 1 point for every missing outcome, and one point for every wrong one.
c. (6 points) We discussed the Therac- 25 in lecture and in an assigned reading.
i) (3 points) Explain the Therac-25's problems and the underlying causes.

We gave full points for specifying at least one of the variety of causes listed in the paper.
ii) (3 points) What would have been a good software development model for the Therac-25's manufacturer to use?

We gave full points for every answer that mentioned "waterfall". Otherwise, we gave 1 point if the answer mentioned testing.

