CS 188 Fall 2019

Exam Prep 1

Q1. Search problems

It is training day for Pacbabies, also known as Hungry Running Maze Games day. Each of k Pacbabies starts in its own assigned start location s_i in a large maze of size $M \ge N$ and must return to its own Pacdad who is waiting patiently but proudly at g_i ; along the way, the Pacbabies must, between them, eat all the dots in the maze.

At each step, all k Pacbabies move one unit to any open adjacent square. The only legal actions are Up, Down, Left, or Right. It is illegal for a Pacbaby to wait in a square, attempt to move into a wall, or attempt to occupy the same square as another Pacbaby. To set a record, the Pacbabies must find an optimal collective solution.

(a) Define a minimal state space representation for this problem.

(b) How large is the state space?

(c) What is the maximum branching factor for this problem?

\bigcirc	4^k	\bigcirc	$4^k 2^{MN}$
\bigcirc	8^k	\bigcirc	$4^{k}2^{4}$

(d) Let MH(p,q) be the Manhattan distance between positions p and q and F be the set of all positions of remaining food pellets and p_i be the current position of Pacbaby i. Which of the following are admissible heuristics?

$h_A: \frac{\sum_{i=1}^{k} MH(p_i, g_i)}{k}$	$\square h_D: \max_{1 \le i \le k} [\min_{f \in F} MH(p_i,$	f)]
$h_B: \max_{1 \le i \le k} MH(p_i, g_i)$	$ \qquad \qquad$	f)]
$h_C: \max_{1 \le i \le k} [\max_{f \in F} MH(p_i, f)]$		f)]

Now suppose that some of the squares are flooded with water. In the flooded squares, it takes two timesteps to travel through the square, rather than one. However, the Pacbabies don't know which squares are flooded and which aren't, until they enter them. After a Pacbaby enters a flooded square, its howls of despair instantly inform all the other Pacbabies of this fact.

(e) Define a minimal space of belief states for this problem.

- (f) How many possible environmental configurations are there in the initial belief state, before the Pacbabies receive any wetness percepts?
- (g) Given the current belief state, how many different belief states can be reached in a single step?

$$\bigcirc \begin{array}{c} 4^k \\ \bigcirc 8^k \end{array} \qquad \bigcirc \begin{array}{c} 4^k 2^{MN} \\ \bigcirc 4^k 2^4 \end{array}$$

Q2. Search

(a) Rubik's Search

Note: You do not need to know what a Rubik's cube is in order to solve this problem.

A Rubik's cube has about 4.3×10^{19} possible configurations, but any configuration can be solved in 20 moves or less. We pose the problem of solving a Rubik's cube as a search problem, where the states are the possible configurations, and there is an edge between two states if we can get from one state to another in a single move. Thus, we have 4.3×10^{19} states. Each edge has cost 1. Note that the state space graph does contain cycles. Since we can make 27 moves from each state, the branching factor is 27. Since any configuration can be solved in 20 moves or less, we have $h^*(n) \leq 20$.

For each of the following searches, estimate the approximate number of states expanded. Mark the option that is closest to the number of states expanded by the search. Assume that the shortest solution for our start state takes exactly 20 moves. Note that 27^{20} is much larger than 4.3×10^{19} .

(i)	DFS Tree Search Best Case:	\bigcirc 20	$\bigcirc 4.3 \times 10^{19}$	\bigcirc 27 ²⁰	\bigcirc	∞ (never finishes)				
	Worst Case:	\bigcirc 20	\bigcirc 4.3 × 10 ¹⁹	$\bigcirc 27^{20}$	0	∞ (never finishes)				
(ii)	DFS graph search Best Case:	\bigcirc 20	\bigcirc 4.3 × 10 ¹⁹	$\bigcirc 27^{20}$	\bigcirc	∞ (never finishes)				
	Worst Case:	\bigcirc 20	$\bigcirc 4.3 \times 10^{19}$	\bigcirc 27 ²⁰	0	∞ (never finishes)				
(iii)	BFS tree search Bost Case:	\bigcirc 20	\bigcirc 4.3 \times 10 ¹⁹	\bigcirc 27 ²⁰	\bigcirc	ac (nover finishes)				
	Worst Case:	\bigcirc 20 \bigcirc 20	$\bigcirc 4.3 \times 10 \\ \bigcirc 4.3 \times 10^{19}$	\bigcirc 27 \bigcirc 27 ²⁰	0	∞ (never finishes) ∞ (never finishes)				
(:)	PES graph gaarch									
(1v)	Br 5 graph search Best Case:	\bigcirc 20	\bigcirc 4.3 \times 10 ¹⁹	\bigcirc 27 ²⁰	\bigcirc	∞ (never finishes)				
	Worst Case:	\bigcirc 20	\bigcirc 4.3 × 10 ¹⁹	\bigcirc 27 ²⁰	\bigcirc	∞ (never finishes)				
(v)	A^* tree search with a	a perfect heurist	ic, $h^*(n)$, Best Case							
	\bigcirc 20	$\bigcirc 4.3 \times 10^{19}$	\bigcirc	27^{20}	\bigcirc	∞ (never finishes)				
(vi)	A^* tree search with a	a bad heuristic.	$h(n) = 20 - h^*(n).$	Worst Case						
()	\bigcirc 20	$\bigcirc 4.3 \times 10^{19}$	\bigcirc	27^{20}	\bigcirc	∞ (never finishes)				
(vii) A^* graph source with a perfect houristic $h^*(n)$. Rest Case										
(*11)	\bigcirc 20	$\bigcirc 4.3 \times 10^{19}$	$\bigcirc \qquad \qquad$	27^{20}	\bigcirc	∞ (never finishes)				
() / *	4 h - h - i h	:= h() = 00 = h*()) Went Core						
(vili	\bigcirc 20	\bigcirc 4.3 \times 10 ¹⁹	$\int n(n) = 20 - h^2(n)$	27^{20}	\bigcirc	∞ (never finishes)				

(b) Limited A^* Graph Search

Consider a variant of A^* graph search called Limited A^* graph search. It is exactly like the normal algorithm, but instead of keeping all of the fringe, at the end of each iteration of the outer loop, the fringe is reduced to just a certain amount of the best paths. I.e. after all children have been inserted, the fringe is cut down to the a certain length. The pseudo-code for normal A^* graph search is reproduced below, the only modification being an argument W for the limit.

1: function A* GRAPH SEARCH(problem, W)

```
fringe \leftarrow an empty priority queue
2:
       fringe \leftarrow \text{INSERT}(\text{MAKE-NODE}(\text{INITIAL-STATE}[problem]), fringe)
3:
       closed \leftarrow an empty set
4:
       ADD INITIAL-STATE [problem] to closed
5:
6:
       loop
           if fringe is empty then
7:
               return failure
8:
           end if
9:
           node \leftarrow \text{REMOVE-FRONT}(fringe)
10:
11:
           if GOAL-TEST(problem, STATE[node]) then
               return node
12:
           end if
13:
           if STATE[node] not in closed then
14:
               ADD STATE[node] to closed
15:
               for successor in GETSUCCESSORS(problem, STATE[node]) do
16:
17:
                   fringe \leftarrow \text{INSERT}(\text{MAKE-SUCCESSOR-NODE}(successor, node), fringe)
               end for
18:
19:
           end if
           fringe = fringe[0:W]
20:
21:
       end loop
22: end function
```

(i) For a positive W, limited A^* graph search is complete.

○ True ○ False

- (ii) For a positive W, limited A^* graph search is optimal.
 - True False
- (iii) Provide the smallest value of W such that this algorithm is equivalent to normal A^* graph search (i.e. the addition of line 20 makes no difference to the execution of the algorithm).

Q3. Pacman's Life

Suppose a maze has height M and width N and there are F food pellets at the beginning. Pacman can move North, South, East or West in the maze.

(a) In this subquestion, the position of Pacman is known, and he wants to pick up all F food pellets in the maze. However, Pacman can move North at most two times overall.

What is the size of a minimal state space for this problem? Give your answer as a product of terms that reference problem quantities such as (but not limited to) M, N, F, etc. Below each term, state the information it encodes. For example, you might write $4 \times MN$ and write number of directions underneath the first term and Pacman's position under the second.

(b) In this subquestion, Pacman is lost in the maze, and does not know his location. However, Pacman still wants to visit every single square (he does not care about collecting the food pellets any more). Pacman's task is to find a sequence of actions which guarantees that he will visit every single square.

What is the size of a minimal state space for this problem? As in part(a), give your answer as a product of terms along with the information encoded by each term. You will receive partial credit for a complete but non-minimal state space.