61A Lecture 25

Friday, October 26

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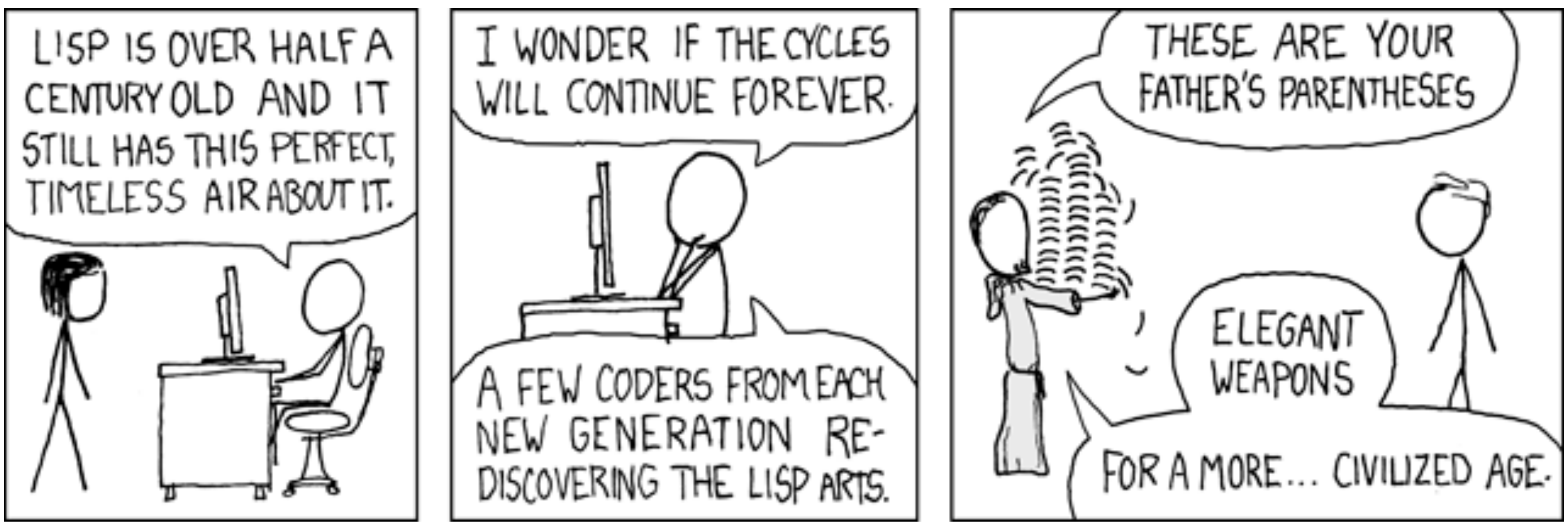

http://imgs.xkcd.com/comics/lisp_cycles.png

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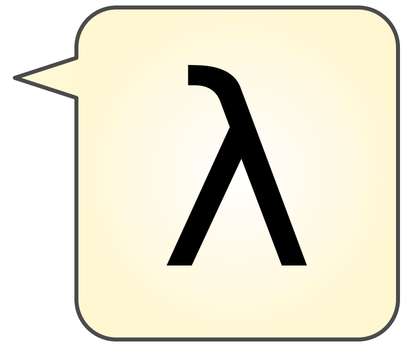

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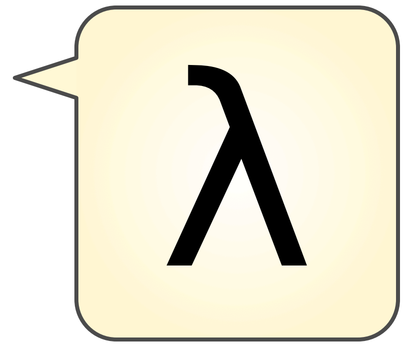

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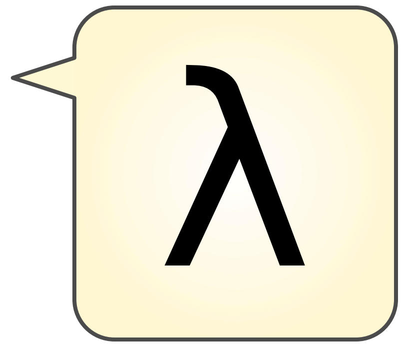

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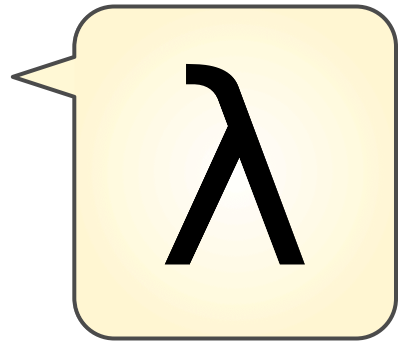

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Symbols are now values

Quotation can also be applied to combinations to form lists.

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> (car '(a b c))
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Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?

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> (define a 1)
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No sign of “a” and “b” in the resulting value

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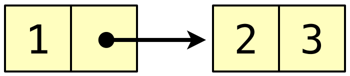

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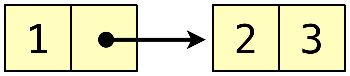

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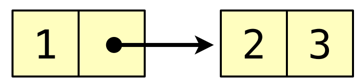

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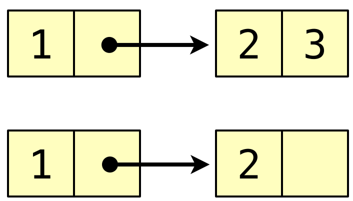

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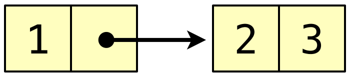

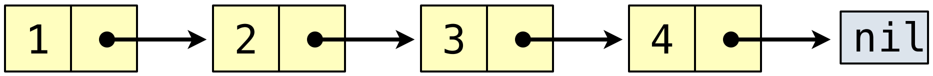

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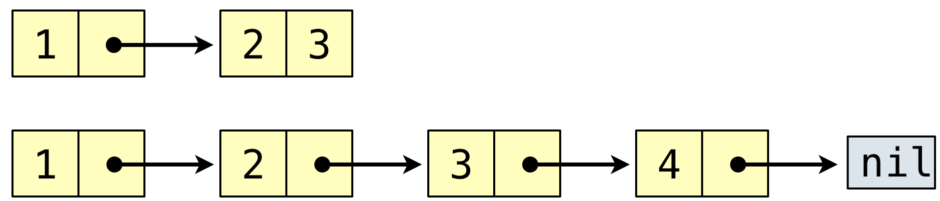

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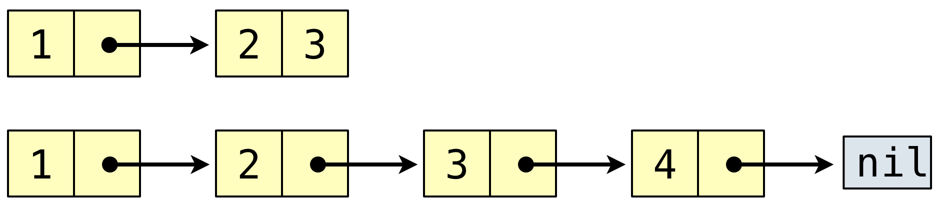

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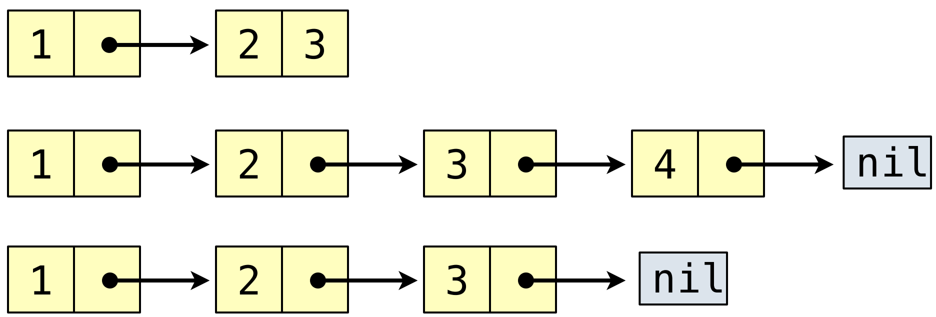

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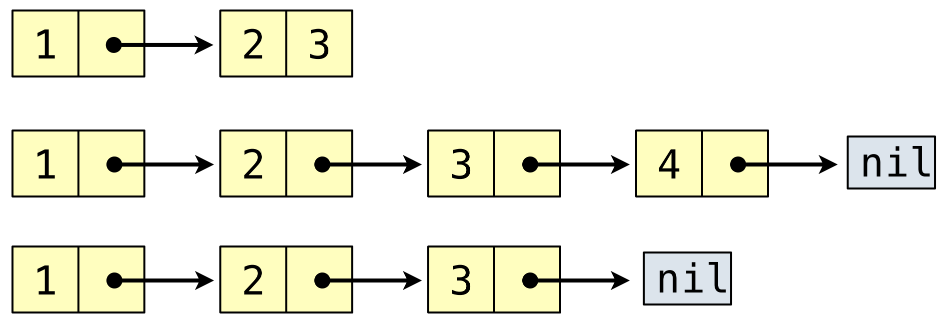

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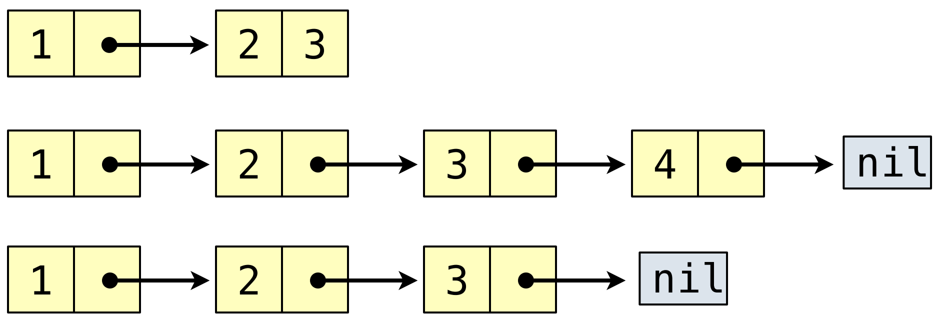

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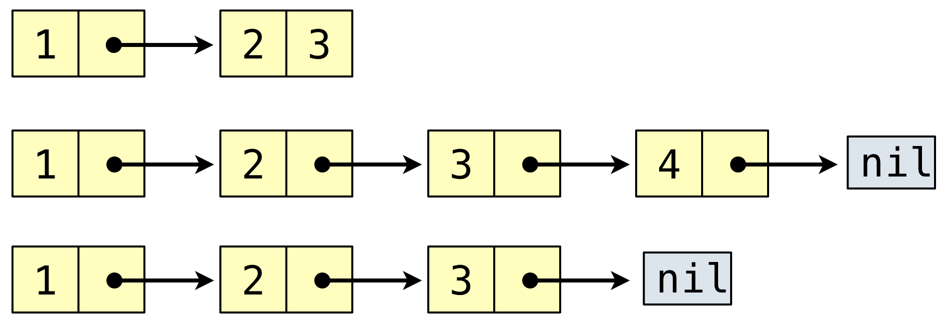

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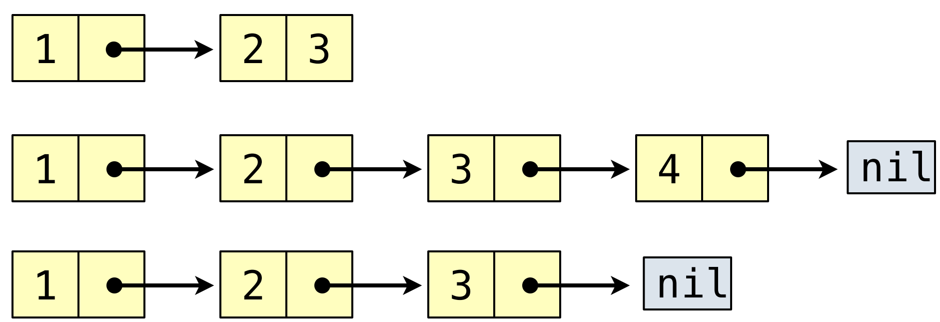

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Coercing a Sorted List to a Binary Search Tree

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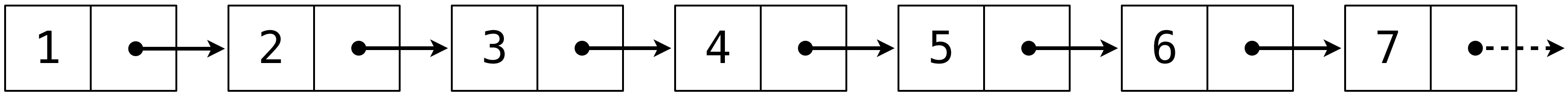

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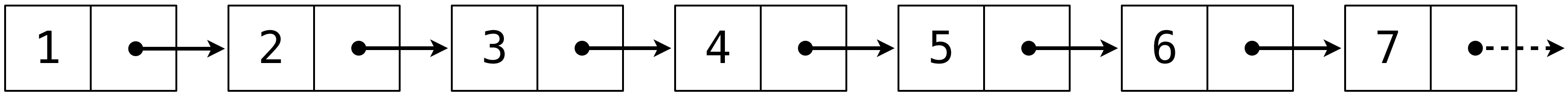

Divide length n into 3 parts: [(n-1)/2 , 1 , (n-1)/2]

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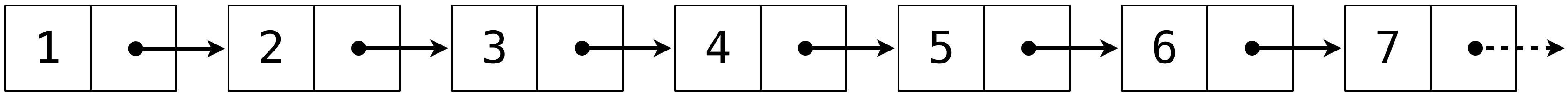

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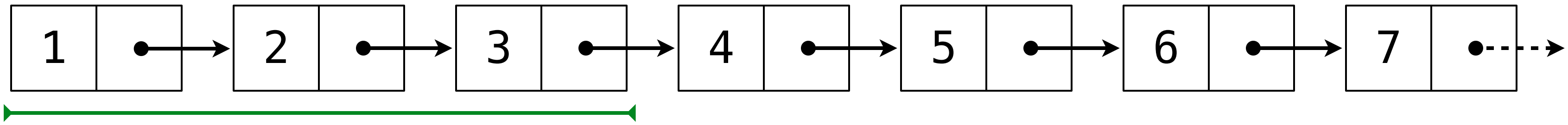

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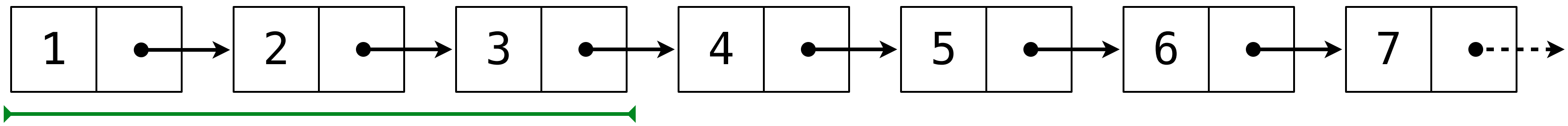

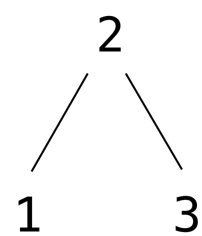

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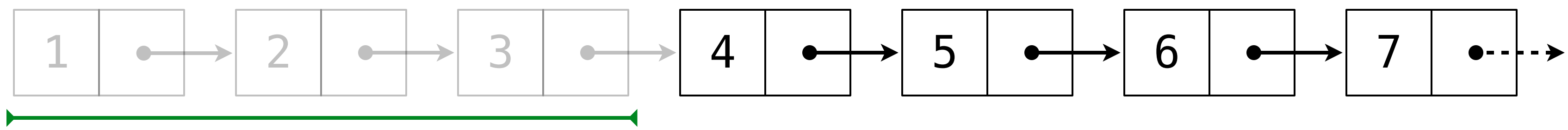

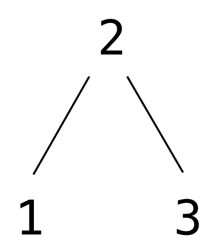

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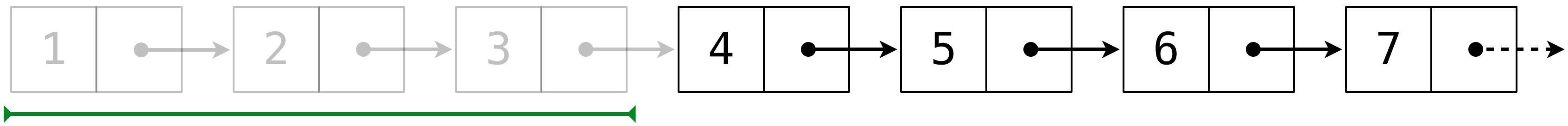

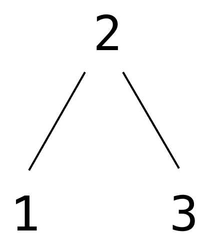

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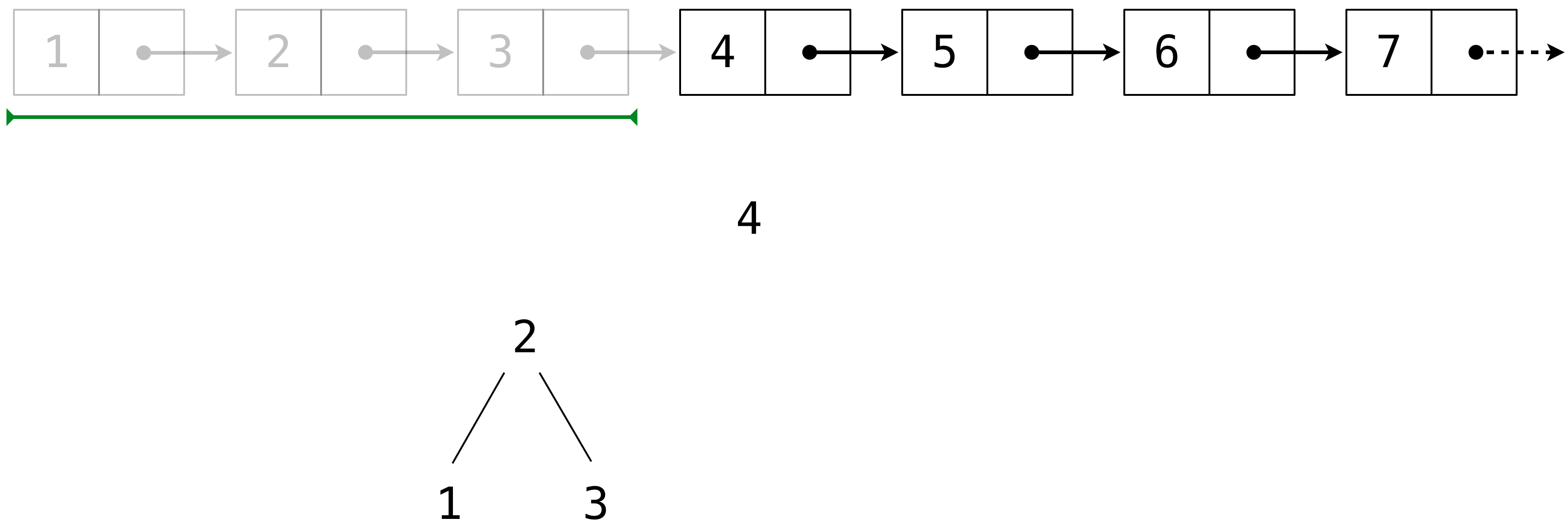


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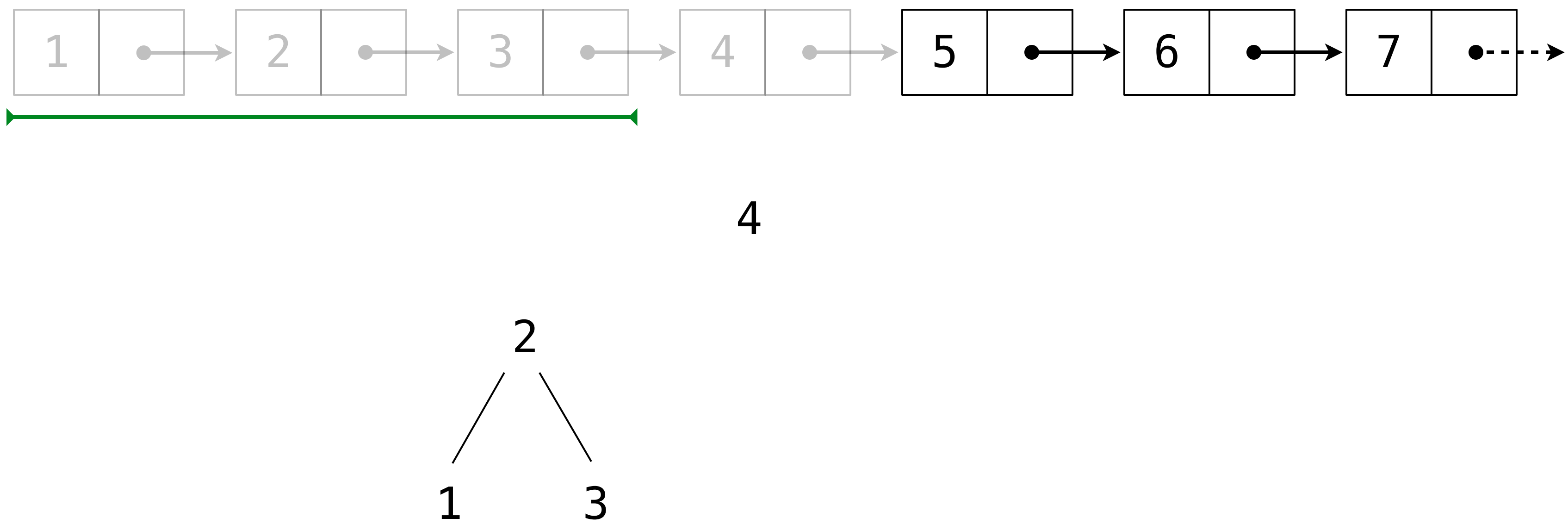

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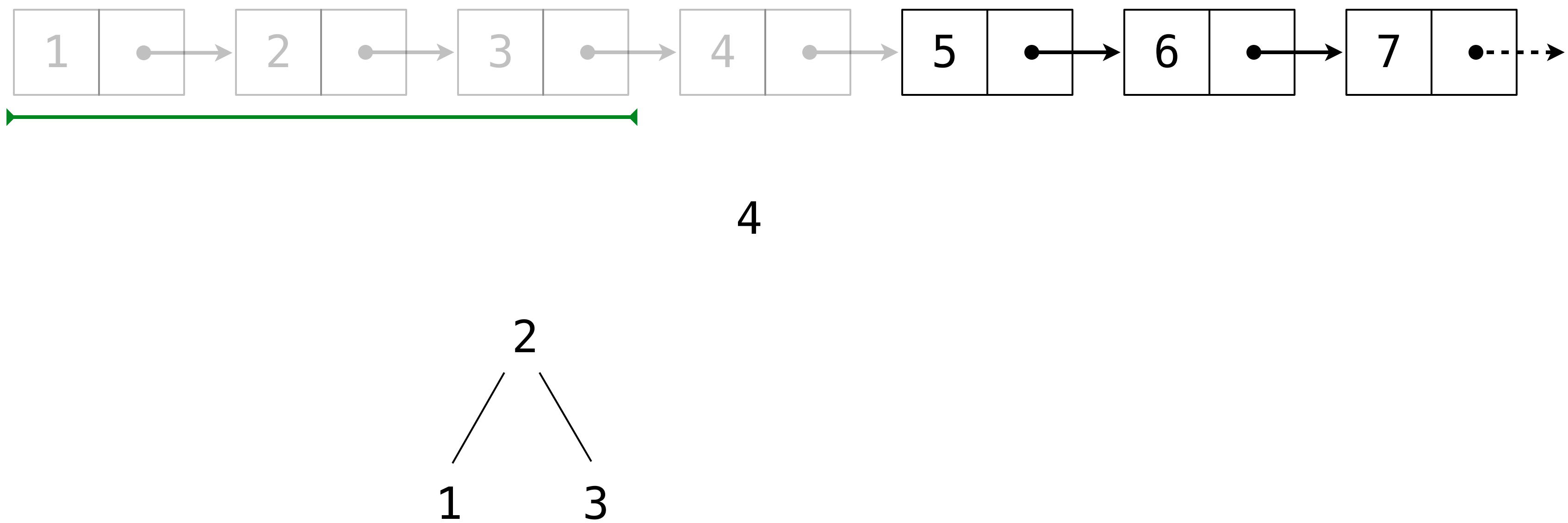

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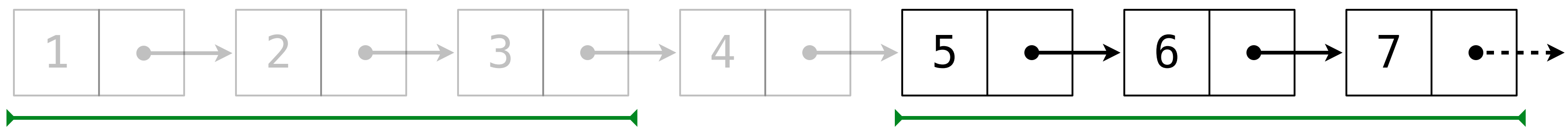

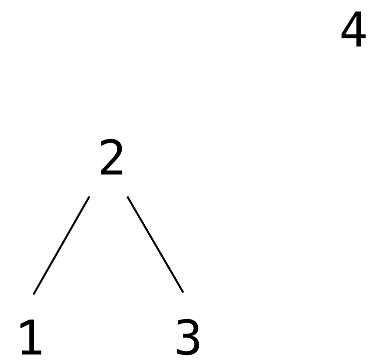

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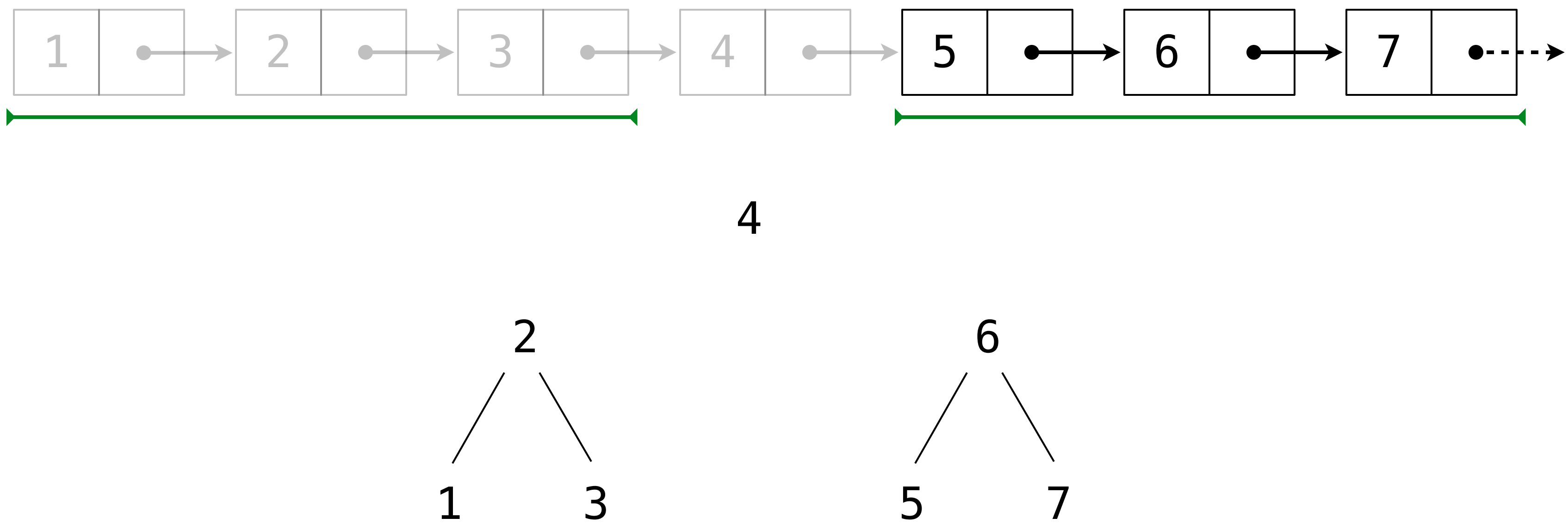

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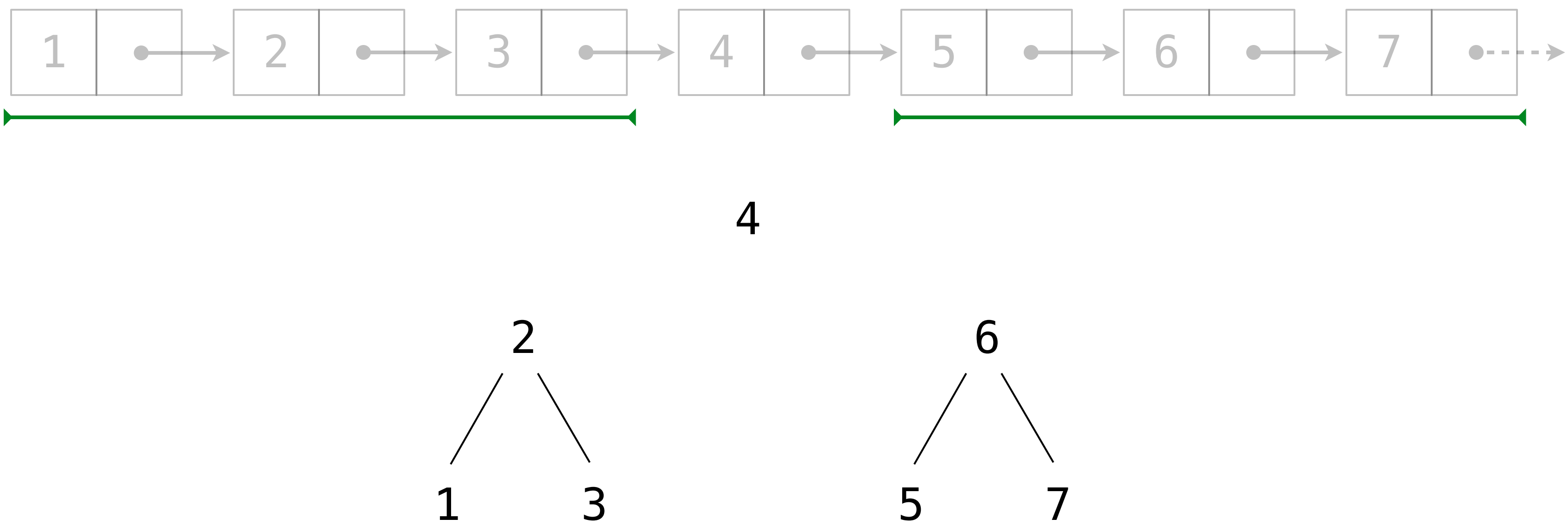

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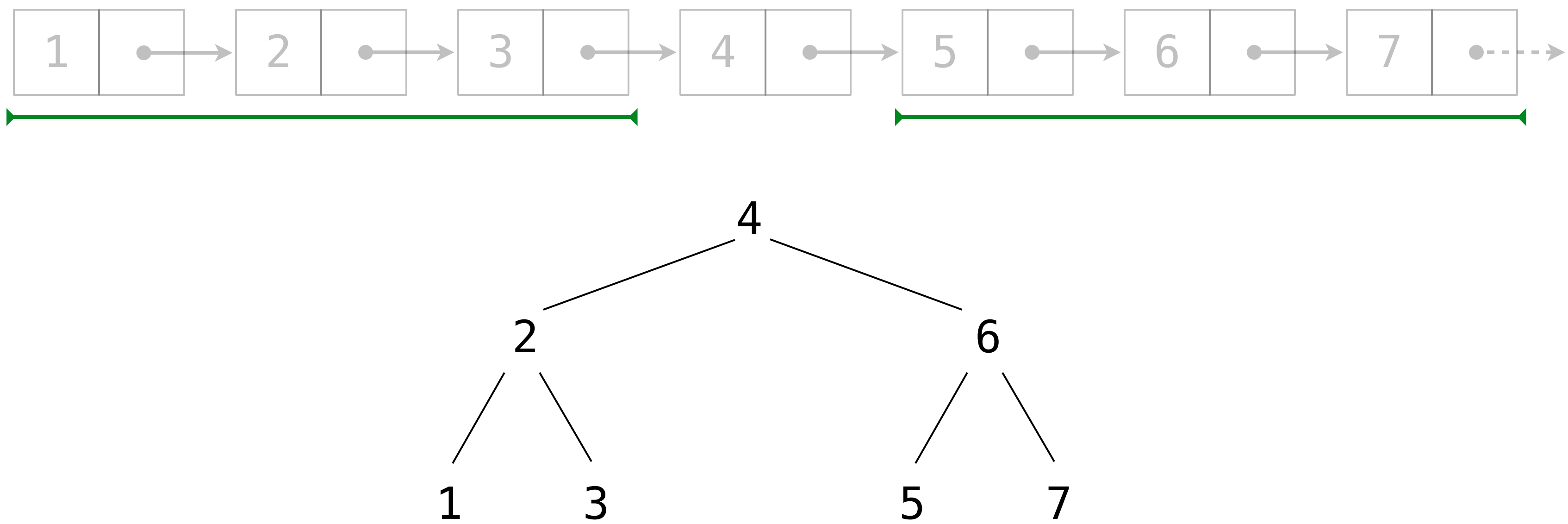

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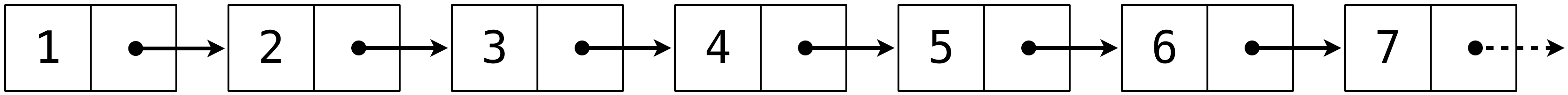

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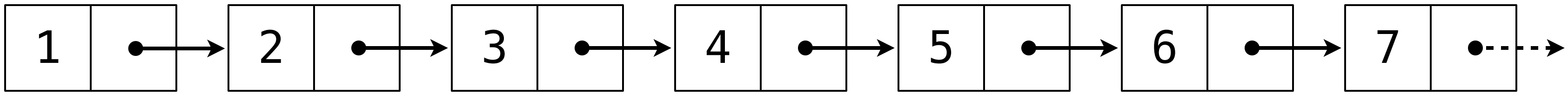


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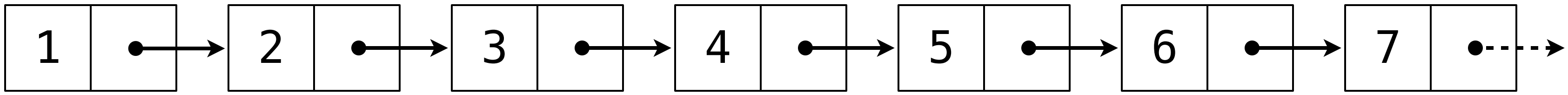

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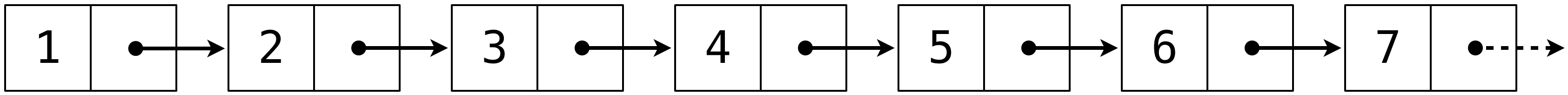


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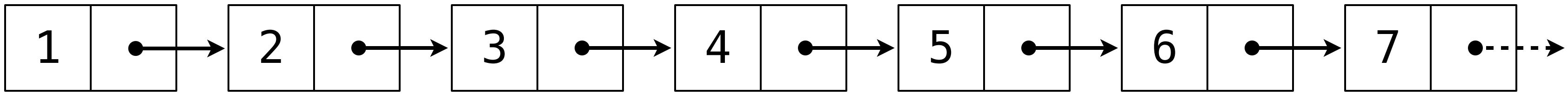


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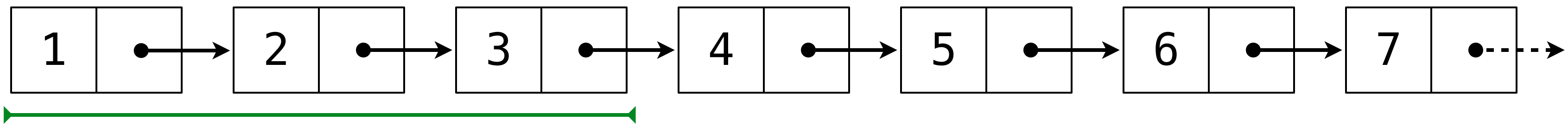


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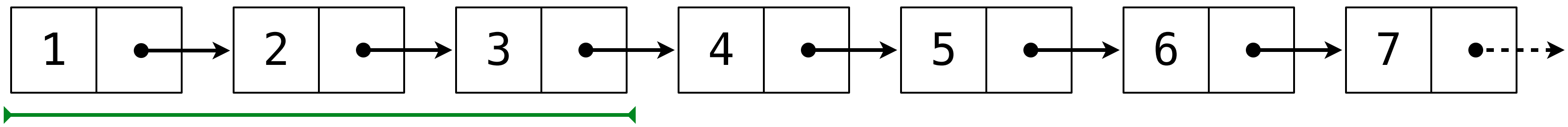


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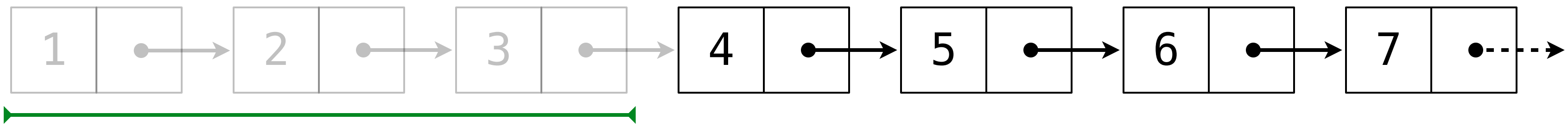


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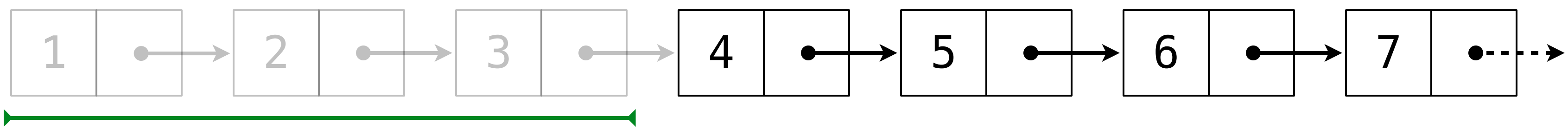
```

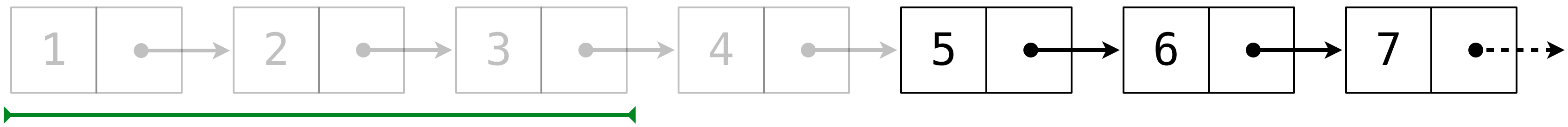
The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
```

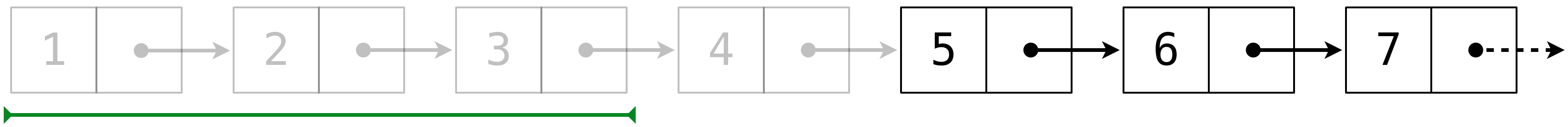
The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
```

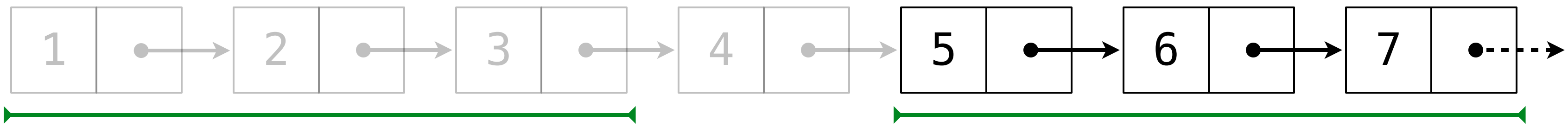
The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
                  (right-result (partial-tree (cdr non-left-elts)
                                              right-size)))
```

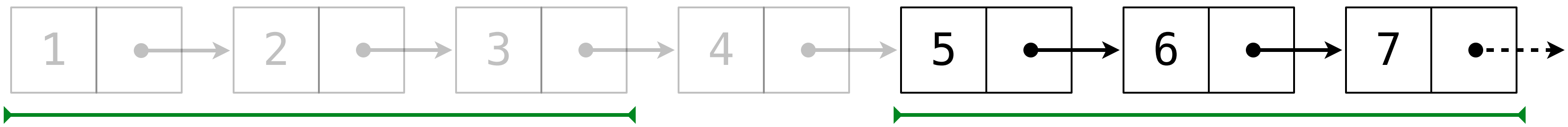
The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
                  (right-result (partial-tree (cdr non-left-elts)
                                              right-size)))
```

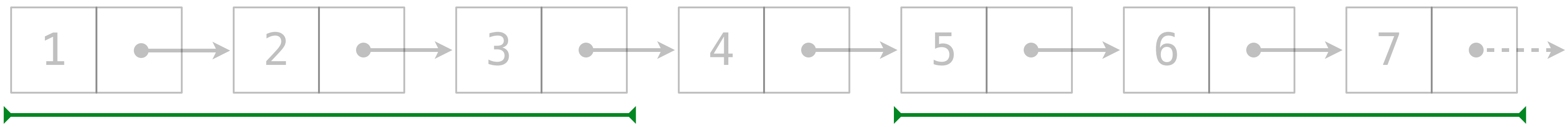
The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
                  (right-result (partial-tree (cdr non-left-elts)
                                              right-size)))
              (let ((right-tree (car right-result))
                    (remaining-elts (cdr right-result)))
```

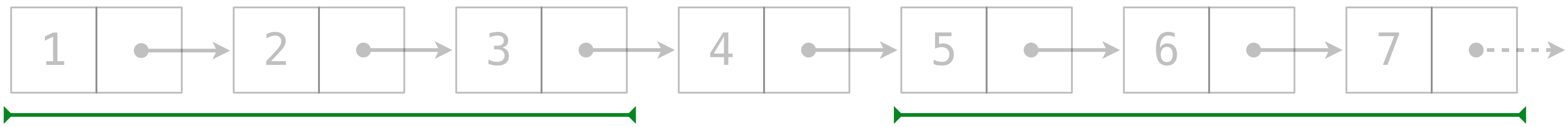
The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
                  (right-result (partial-tree (cdr non-left-elts)
                                              right-size)))
              (let ((right-tree (car right-result))
                    (remaining-elts (cdr right-result)))
```

The Let Special Form




```
(define (entry tree) ...)
(define (left-branch tree) ...)
(define (right-branch tree) ...)
(define (make-tree entry left right) ...)
(define (list->tree elements)
  (car (partial-tree elements (length elements))))

(define (partial-tree elts n)
  (if (= n 0)
      (cons nil elts)
      (let ((left-size (quotient (- n 1) 2)))
        (let ((left-result (partial-tree elts left-size)))
          (let ((left-tree (car left-result))
                (non-left-elts (cdr left-result))
                (right-size (- n (+ left-size 1))))
            (let ((this-entry (car non-left-elts))
                  (right-result (partial-tree (cdr non-left-elts)
                                              right-size)))
              (let ((right-tree (car right-result))
                    (remaining-elts (cdr right-result)))
                (cons (make-tree this-entry left-tree right-tree)
                      remaining-elts))))))))
```

The Begin Special Form

```
(begin <exp1> <exp2> ... <expn>)
```

The Begin Special Form

```
(begin <exp1> <exp2> ... <expn>)
```

Demo