Announcements
Distributed Computing
Distributed Computing

A distributed computing application consists of multiple programs running on multiple computers that together coordinate to perform some task
• Computation is performed in parallel by many computers
• Information can be restricted to certain computers
• Redundancy and geographic diversity improve reliability

Characteristics of distributed computing:
• Computers are independent — they do not share memory
• Coordination is enabled by messages passed across a network
• Individual programs have differentiating roles

Distributed computing for large-scale data processing:
• Databases respond to queries over a network
• Data sets can be partitioned across multiple machines (next lecture)
Network Messages

Computers communicate via messages: sequences of bytes transmitted over a network

Messages can serve many purposes:
• Send data to another computer
• Request data from another computer
• Instruct a program to call a function on some arguments
• Transfer a program to be executed by another computer

Messages conform to a message protocol adopted by both the sender (to encode the message) & receiver (to interpret the message)
• For example, bits at fixed positions may have fixed meanings
• Components of a message may be separated by delimiters
• Protocols are designed to be implemented by many different programming languages on many different types of machines
Internet Protocol
The Internet Protocol

The Internet Protocol (IP) specifies how to transfer packets of data among networks
• Networks are inherently unreliable at any point
• The structure of a network is dynamic, not fixed
• No system exists to monitor or track communications

Packets are forwarded toward their destination on a best effort basis
Programs that use IP typically need a policy for handling lost packets

All machines know IPv4
Decremented on forwarding

Max length: 216 = 65,536
E.g., 192.168.1.1

Packets can't survive forever

The packet knows its size
Where to send the packet
Where to send error reports

E.g., 192.168.1.1

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Transmission Control Protocol
Transmission Control Protocol

The design of the Internet Protocol (IPv4) imposes constraints:

- Packets are limited to 65,535 bytes each
- Packets may arrive in a different order than they were sent
- Packets may be duplicated or lost

The Transmission Control Protocol (TCP) improves reliability:

- Ordered, reliable transmission of arbitrary byte streams
- Implemented using the IP. Every TCP connection involves sending IP packets
- Each packet in a TCP session has a sequence number:
  - The receiver can correctly order packets that arrive out of order
  - The receiver can ignore duplicate packets
- All received packets are acknowledged; both parties know that transmission succeeded
- Packets that aren't acknowledged are sent repeatedly

The socket module in Python implements the TCP
TCP Handshakes

All TCP connections begin with a sequence of messages called a "handshake" which verifies that communication is possible

"Can you hear me now?"  *Let's design a handshake protocol*

Handshake Goals:
- Computer A knows that it can send data to and receive data from Computer B
- Computer B knows that it can send data to and receive data from Computer A
- Lots of separate connections can exist without any confusion
- The number of required messages is minimized

Communication Rules:
- Computer A can send an initial message to Computer B requesting a new connection
- Computer B can respond to messages from Computer A
- Computer A can respond to messages from Computer B
Message Sequence of a TCP Connection

- **Computer A**
- Establishes packet numbering system
- Synchronization request
  - Acknowledgement & synchronization request
  - Acknowledgement
  - Data message from A to B
    - Acknowledgement
    - Data message from B to A
      - Acknowledgement
      - Termination signal
        - Acknowledgement & termination signal
        - Acknowledgement
  - Termination signal
    - Acknowledgement
Client/Server Architecture
The Client/Server Architecture

One server provides information to multiple clients through request and response messages

**Server role:** Respond to service requests with requested information

**Client role:** Request information and make use of the response

**Abstraction:** The client knows what service a server provides, but not how it is provided
Client/Server Example: The World Wide Web

The client is a web browser (e.g., Firefox):
- Request content for a location
- Interpret the content for the user

The server is a web server:
- Interpret requests and respond with content
The Hypertext Transfer Protocol

The Hypertext Transfer Protocol (HTTP) is a protocol designed to implement a Client/Server architecture

Browser issues a GET request to a server at www.nytimes.com for the content (resource) at location "pages/todayspaper"

Server response contains more than just the resource itself:
- Status code, e.g. 200 OK, 404 Not Found, 403 Forbidden, etc.
- Date of response; type of server responding
- Last-modified time of the resource
- Type of content and length of content
Properties of a Client/Server Architecture

Benefits:
• Creates a separation of concerns among components
• Enforces an abstraction barrier between client and server
• A centralized server can reuse computation across clients

Liabilities:
• A single point of failure: the server
• Computing resources become scarce when demand increases

Common use cases:
• Databases — The database serves responses to query requests
• Open Graphics Library (OpenGL) — A graphics processing unit (GPU) serves images to a central processing unit (CPU)
• Internet file and resource transfer: HTTP, FTP, email, etc.
Peer-to-Peer Architecture
The Peer-to-Peer Architecture

All participants in a distributed application contribute computational resources: processing, storage, and network capacity.

Messages are relayed through a network of participants.

Each participant has only partial knowledge of the network.

Network Structure Concerns

Some data transfers on the Internet are faster than others

The time required to transfer a message through a peer-to-peer network depends on the route chosen
Example: Skype

Skype is a Voice Over IP (VOIP) system that uses a hybrid peer-to-peer architecture.

Login & contacts are handled via a centralized server.

Conversations between two computers that cannot send messages to each other directly are relayed through supernodes.

Any Skype client with its own IP address may be a supernode.