Announcements
- Guerrilla Section 5 this weekend on Scheme & functional programming
- Sunday 4/5 12:00pm – 2:30pm in 271 Soda
- Homework 7 due Wednesday 4/8 @ 11:59pm
- Homework party Tuesday 4/7 5pm-6:30pm in 2050 VLSB
- Quiz 3 released Tuesday 4/7 & due Thursday 4/9 @ 11:59pm
- Please check your grades on glookup and request regrades for mistakes
  - http://cs61a.org/regrades.html

Programming Languages

A computer typically executes programs written in many different programming languages

- Machine languages: statements are interpreted by the hardware itself
  - A fixed set of instructions invoke operations implemented by the circuitry of the central processing unit (CPU)
  - Operations refer to specific hardware memory addresses; no abstraction mechanisms

- High-level languages: statements and expressions are interpreted by another program or compiled (translated) into another language
  - Provide means of abstraction such as naming, function definition, and objects
  - Abstract away system details to be independent of hardware and operating system

```
Python 3 Byte Code
LOAD_FAST                0 (x)
BINARY_MULTIPLY
RETURN_VALUE
```

Metalinguistic Abstraction

A powerful form of abstraction is to define a new language that is tailored to a particular type of application or problem domain

- Type of application: Erlang was designed for concurrent programs. It has built-in elements for expressing concurrent communication. It is used, for example, to create chat servers with many simultaneous connections
- Problem domain: The MediaWiki mark-up language was designed for generating static web pages. It has built-in elements for text formatting and cross-page linking. It is used, for example, to create Wikipedia pages

A programming language has:
- Syntax: The legal statements and expressions in the language
- Semantics: The execution/evaluation rule for those statements and expressions

To create a new programming language, you either need a:
- Specification: A document describes the precise syntax and semantics of the language
- Canonical Implementation: An interpreter or compiler for the language

Parsing

A Parser takes text and returns an expression

<table>
<thead>
<tr>
<th>Text</th>
<th>Lexical analysis</th>
<th>Tokens</th>
<th>Syntactic analysis</th>
<th>Expression</th>
</tr>
</thead>
<tbody>
<tr>
<td>'(+ 1'</td>
<td>'(', '+', 1</td>
<td>('(', '+', 1, ')')</td>
<td>Pair('(', Pair('+', 1), ')')</td>
<td></td>
</tr>
<tr>
<td>'(- 23)'</td>
<td>'(', '-', 23</td>
<td>('(', '-', 23, ')')</td>
<td>Pair('(', Pair('-', 23), ')')</td>
<td></td>
</tr>
<tr>
<td>'(* 4 5.6)'</td>
<td>'(', '*', 4, 5.6</td>
<td>('(', '*', 4, 5.6, ')')</td>
<td>Pair('(', Pair('*', 4, 5.6), ')')</td>
<td></td>
</tr>
</tbody>
</table>

Recursive Syntactic Analysis

A predictive recursive descent parser inspects only k tokens to decide how to proceed, for some fixed k

Can English be parsed via predictive recursive descent?

The horse raced past the barn fell.

Iterative process
- Checks for malformed tokens
- Determines types of tokens
- Processes one line at a time

Tree-recursive process
- Balances parentheses
- Returns tree structure
- Processes multiple lines

Sentence subject

The horse raced past the barn fell.
Reading Scheme Lists

A Scheme list is written as elements in parentheses:

\((\text{element}_0\ \text{element}_1\ \ldots \text{element}_n)\)

Each \text{element} can be a combination or primitive:

\((+\ (*\ (3\ (+\ (2\ 4)\ (+\ (3\ 5)))\ (+\ (-10\ 7)\ 6)))\)

The task of parsing a language involves coercing a string representation of an expression to the expression itself.

Parsers must validate that expressions are well-formed.

(Demo)

http://composingprograms.com/examples/scalc/scheme_reader.py.html

Syntactic Analysis

Syntactic analysis identifies the hierarchical structure of an expression, which may be nested.

Recursive call: scheme_read sub-expressions and combine them.

(Demo)

The Pair Class

The Pair class represents Scheme pairs and lists. A list is a pair whose second element is either a list or nil.

The Pair class contains instance attributes:

- self
- first
- second

Some methods only apply to well-formed lists.

Scheme expressions are represented as Scheme lists! Source code is data.

(Demo)

Calculator Syntax

The Calculator language has primitive expressions and call expressions. (That's it!)

A primitive expression is a number: 2, -4, 5.6

A call expression is a combination that begins with an operator (+, -, *, /) followed by 0 or more expressions: (+ 1 2 3), (/ 3 (+ 4 5))

Expressions are represented as Scheme Lists (Pair instances) that encode tree structures.

(Demo)

Calculator Semantics

The value of a calculator expression is defined recursively.

Primitive: A number evaluates to itself.

Call: A call expression evaluates to its argument values combined by an operator. Depending on the arguments:

- Product of the arguments
- If one argument, negate it. If more than one, subtract the rest from the first.
- If one argument, invert it. If more than one, divide the rest from the first.

The Eval Function

The eval function computes the value of an expression, which is always a number.

It is a generic function that dispatches on the type of the expression (primitive or call).

(Demo)


Evaluation

Recursive call
returns a number for each operand.

A number evaluates...

A call expression evaluates...

to its argument values combined by an operator.
Applying Built-in Operators

The apply function applies some operation to a (Scheme) list of argument values.

In calculator, all operations are named by built-in operators: +, -, *, /.

```
def calc_apply(operator, args):
    if operator == '+':
        return reduce(add, args, 0)
    elif operator == '-':
        ...
    elif operator == '*':
        ...
    else:
        raise TypeError
```

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Interactive Interpreters

Read-Eval-Print Loop

The user interface for many programming languages is an interactive interpreter.

1. Print a prompt
2. Read text input from the user
3. Parse the text input into an expression
4. Evaluate the expression
5. If any errors occur, report those errors, otherwise
6. Print the value of the expression and repeat

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Raising Exceptions

Exceptions are raised within lexical analysis, syntactic analysis, eval, and apply.

Example exceptions
- **Lexical analysis**: The token 2.3.4 raises ValueError("invalid numeral")
- **Syntactic analysis**: An extra ) raises SyntaxError("unexpected token")
- **Eval**: An empty combination raises TypeError("() is not a number or call expression")
- **Apply**: No arguments to - raises TypeError("- requires at least 1 argument")

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Handling Exceptions

An interactive interpreter prints information about each error.

A well-designed interactive interpreter should not halt completely on an error, so that the user has an opportunity to try again in the current environment.