Announcements

- HW9 out tonight, due 4/3
- Ants extra credit due 4/3
  - See Piazza for submission instructions
Data Structure Applications

The data structures we cover in 61A are used everywhere in CS

More about data structures in 61B

Example: recursive lists (also called linked lists)

- Operating systems
- Interpreters and compilers
- Anything that uses a queue

The Scheme programming language, which we will learn soon, uses recursive lists as its primary data structure
Example: Environments

```
def curry(fn):
    def outer(x):
        def inner(y):
            return fn(x, y)
        return inner
    return outer

from operator import add
curry(add)(3)(4)
```
Trees with Internal Node Values

Trees can have values at internal nodes as well as their leaves.

```python
class Tree(object):
    def __init__(self, entry, left=None, right=None):
        self.entry = entry
        self.left = left
        self.right = right

def fib_tree(n):
    if n == 1:
        return Tree(0)
    if n == 2:
        return Tree(1)
    left = fib_tree(n - 2)
    right = fib_tree(n - 1)
    return Tree(left.entry + right.entry, left, right)
```
Implementing Sets

What we should be able to do with a set:

- Membership testing: Is a value an element of a set?
- Union: Return a set with all elements in set1 or set2
- Intersection: Return a set with any elements in set1 and set2
- Adjunction: Return a set with all elements in s and a value v
Sets as Unordered Sequences

Proposal 1: A set is represented by a recursive list that contains no duplicate items

This is how we implemented dictionaries

```python
def empty(s):
    return s is Rlist.empty

def set_contains(s, v):
    if empty(s):
        return False
    elif s.first == v:
        return True
    return set_contains(s.rest, v)
```
### Sets as Unordered Sequences

**def adjoin_set(s, v):**
- If `set_contains(s, v):`
  - Return `s`
- Return `Rlist(v, s)`

**def intersect_set(set1, set2):**
- `f = lambda v: set_contains(set2, v)`
- Return `filter_rlist(set1, f)`

**def union_set(set1, set2):**
- `f = lambda v: not set_contains(set2, v)`
- `set1_not_set2 = filter_rlist(set1, f)`
- Return `extend_rlist(set1_not_set2, set2)`

**Time order of growth**
- $\Theta(n)$
- The size of the set
- $\Theta(n^2)$
- Assume sets are the same size

- $\Theta(n^2)$
Sets as Ordered Sequences

**Proposal 2:** A set is represented by a recursive list with unique elements ordered from least to greatest

```python
def set_contains2(s, v):
    if empty(s) or s.first > v:
        return False
    elif s.first == v:
        return True
    return set_contains(s.rest, v)
```

Order of growth? $\Theta(n)$
This algorithm assumes that elements are in order.

```python
def intersect_set2(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    e1, e2 = set1.first, set2.first
    if e1 == e2:
        rest = intersect_set2(set1.rest, set2.rest)
        return Rlist(e1, rest)
    elif e1 < e2:
        return intersect_set2(set1.rest, set2)
    elif e2 < e1:
        return intersect_set2(set1, set2.rest)
```

Order of growth? $\Theta(n)$
Proposal 3: A set is represented as a Tree. Each entry is:

- Larger than all entries in its left branch and
- Smaller than all entries in its right branch
Membership in Tree Sets

Set membership tests traverse the tree

• The element is either in the left or right sub-branch
• By focusing on one branch, we reduce the set by about half

```python
def set_contains3(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains3(s.right, v)
    elif s.entry > v:
        return setContains3(s.left, v)
```

If 9 is in the set, it is in this branch

Order of growth?
Adjoining to a Tree Set

Right!  Left!  Right!  Stop!

1  7  11
3  7  11
1  7  11

3  9

5

4

9

8

7

9

7

8

11

8

None

None

None

Cal
What Did I Leave Out?

Sets as ordered sequences:
- Adjoining an element to a set
- Union of two sets

Sets as binary trees:
- Intersection of two sets
- Union of two sets

That's homework 9!
Social Implications / Programming Practices

- Why things go wrong
- What can we do about this
Therac-25 Case Study

- Medical imaging device

Figure 9  Typical Therac-25 Facility
Therac-25 Case Study

- What happened?
- 6 serious injuries
- 4 deaths
- Otherwise effective – saved hundreds of lives
Lesson to be learned

- Social responsibility in engineering
- First real incident of fatal software failure
- Bigger issue
  - No bad guys
  - Honestly believed there was nothing wrong
“Software Rot”

- Other engineering fields: clear sense of degradation and decay
- Can software become brittle or fractured?
A bigger picture

- **All software is part of a bigger system**
  - Software degrades because:
    - Other piece of software changes
    - Hardware changes
    - Environment changes
Ex: Compatibility Issues

This program has known compatibility issues

Check online to see if solutions are available from the Microsoft website. If solutions are found, Windows will automatically display a website that lists steps you can take.

Publisher: Symantec
Location: C:\Users\japtop\AppData\Local\Temp\NIS15.5.0.23\Setup.exe

Norton Internet Security ver. 2007 and 2008 is incompatible with this version of Windows. For more information, contact Symantec.

Toggle Hide details
Check for solutions online
Run program
Cancel

Don't show this message again
A bigger issue

- The makers of the Therac did not fully understand the complexity of their software.
- Complexity of constructs in other fields more apparent.
A “simple” program

- This program can delete any file you can
Complexity in the Therac-25

- Abundant user interface issues
- Cursor position and field entry
- Default values
- Too many error messages
Too many error messages
Too many error messages

Internet Explorer

When you see a dialog box like this, click 'Yes' to make it go away. If available, click the checkbox first to avoid being bothered by it again.

- In the future, do not show this message.

[Yes]  [No]
(More) Complexity in the Therac-25

- No atomic test-and-set
- No hardware interlocks
How can we solve these things?

- Know your user
- Fail-Soft (or Fail-Safe)
- Audit Trail
- Correctness from the start
- Redundancy
def mutable_rlist():
    def dispatch(message, value=None):
        nonlocal contents
        if message == 'first':
            return first(contents)
        if message == 'rest':
            return rest(contents)
        if message == 'len':
            return len_rlist(contents)
        ...
        else:
            print('Unknown message')
    return dispatch
Correctness from the start

- Edsger Dijkstra: “On the Cruelty of Really Teaching Computing Sciences”
- CS students shouldn’t use computers
- Rigorously prove correctness of their programs
- Correctness proofs
- Compilation (pre-execution) analysis
On debugging

- Black box debugging
- Glass box debugging
- Don’t break what works
- Golden rule of debugging...
Golden rule of debugging

☐ “Debug by subtraction, not by addition”

☐ Prof. Brian Harvey