

**Disclaimer:** This discussion worksheet is fairly long and is not designed to be finished in a single section. Some of these questions are of the level that you might see on an exam and are meant to provide extra practice with asymptotic analysis.

## 1 More Running Time

Give the worst case and best case running time in  $\Theta(\cdot)$  notation in terms of  $M$  and  $N$ .

(a) Assume that `slam()`  $\in \Theta(1)$  and returns a boolean.

```
1 public void comeon(int M, int N) {
2     int j = 0;
3     for (int i = 0; i < N; i += 1) {
4         for (; j < M; j += 1) {
5             if (slam(i, j))
6                 break;
7         }
8     }
9
10    for (int k = 0; k < 1000 * N; k += 1) {
11        System.out.println("space jam");
12    }
13 }
```

(b) *Exam Practice:* Give the worst case and best case running time in  $\Theta(\cdot)$  notation in terms of  $N$  for `find`.

```
1 public static boolean find(int tgt, int[] arr) {
2     int N = arr.length;
3     return find(tgt, arr, 0, N);
4 }
5 private static boolean find(int tgt, int[] arr, int lo, int hi) {
6     if (lo == hi || lo + 1 == hi) {
7         return arr[lo] == tgt;
8     }
9     int mid = (lo + hi) / 2;
10    for (int i = 0; i < mid; i += 1) {
11        System.out.println(arr[i]);
12    }
13    return arr[mid] == tgt || find(tgt, arr, lo, mid)
14        || find(tgt, arr, mid, hi);
15 }
```

## 2 Recursive Running Time

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For the following recursive functions, give the worst case and best case running time in  $\Theta(\cdot)$  notation.

(a) Give the running time in terms of  $N$ .

```
1 public void andslam(int N) {
2     if (N > 0) {
3         for (int i = 0; i < N; i += 1) {
4             System.out.println("bigballer.jpg");
5         }
6         andslam(N / 2);
7     }
8 }
```

(b) Give the running time for `andwelcome(arr, 0, N)` where  $N$  is the length of the input array `arr`.

```
1 public static void andwelcome(int[] arr, int low, int high) {
2     System.out.print("[ ");
3     for (int i = low; i < high; i += 1) {
4         System.out.print("loyal ");
5     }
6     System.out.println("]");
7     if (high - low > 0) {
8         double coin = Math.random();
9         if (coin > 0.5) {
10            andwelcome(arr, low, low + (high - low) / 2);
11        } else {
12            andwelcome(arr, low, low + (high - low) / 2);
13            andwelcome(arr, low + (high - low) / 2, high);
14        }
15    }
16 }
```

(c) Give the running time in terms of  $N$ .

```
1 public int tothe(int N) {
2     if (N <= 1) {
3         return N;
4     }
5     return tothe(N - 1) + tothe(N - 1) + tothe(N - 1);
6 }
```

(d) *Exam Practice:* Give the running time in terms of  $N$

```
1 public static void spacejam(int N) {
2     if (N == 1) {
3         return;
4     }
5     for (int i = 0; i < N; i += 1) {
6         spacejam(N-1);
7     }
8 }
```

### 3 Hey you watchu gon do?

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For each example below, there are two algorithms solving the same problem. Given the asymptotic runtimes for each, is one of the algorithms **guaranteed** to be faster? If so, which? And if neither is always faster, explain why. Assume the algorithms have very large input (so  $N$  is very large).

- (a) Algorithm 1:  $\Theta(N)$ , Algorithm 2:  $\Theta(N^2)$
- (b) Algorithm 1:  $\Omega(N)$ , Algorithm 2:  $\Omega(N^2)$
- (c) Algorithm 1:  $O(N)$ , Algorithm 2:  $O(N^2)$
- (d) Algorithm 1:  $\Theta(N^2)$ , Algorithm 2:  $O(\log N)$
- (e) Algorithm 1:  $O(N \log N)$ , Algorithm 2:  $\Omega(N \log N)$

Why did we need to assume that  $N$  was large?

### 4 Big Ballin' Bounds

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1. Prove the following bounds by finding some constant  $M > 0$  and input  $N > 0$  for  $M \in \mathbb{R}, N \in \mathbb{N}$  such that  $f$  and  $g$  satisfy the relationship.
  - (a)  $f \in O(g)$  for  $f = 2n, g = n^2$
  - (b)  $f \in \Omega(g)$  for  $f = 0.1n, g = 40$
  - (c)  $f \in \Theta(g)$  for  $f = \log(n), g = \log(n^a)$ , for  $a > 0$ .
2. Answer the following claims with true or false. If false, provide a counterexample.
  - (a) If  $f(n) \in O(g(n))$ , then  $500f(n) \in O(g(n))$ .
  - (b) If  $f(n) \in \Theta(g(n))$ , then  $2^{f(n)} \in \Theta(2^{g(n)})$ .