CS61B Lecture #18: Assorted Topics

- Views
- Maps
- More partial implementations
- Array vs. linked: tradeoffs
- Sentinels
- Specialized sequences: stacks, queues, deques
- Circular buffering
- Recursion and stacks
- Adapters
Views

New Concept: A view is an alternative presentation of (interface to) an existing object.

- For example, the sublist method is supposed to yield a “view of” part of an existing list:

```
List<String> L = new ArrayList<String>();
L.add("at"); L.add("ax"); ...
List<String> SL = L.sublist(1, 4);
```

- Example: after `L.set(2, "bag")`, value of `SL.get(1)` is "bag", and after `SL.set(1, "bad")`, value of `L.get(2)` is "bad".

- Example: after `SL.clear()`, L will contain only "at" and "cat".

- Small challenge: “How do they do that?!”
Maps

• A Map is a kind of “modifiable function:”

```java
package java.util;
public interface Map<Key, Value> {
    Value get(Object key); // Value at KEY.
    Object put(Key key, Value value); // Set get(KEY) -> VALUE
}
```

```java
Map<String, String> f = new TreeMap<String, String>();
f.put("Paul", "George"); f.put("George", "Martin");
f.put("Dana", "John");
// Now f.get("Paul").equals("George")
//   f.get("Dana").equals("John")
//   f.get("Tom") == null
```
Map Views

public interface Map<Key, Value> {
    // Continuation

    /* Views of Maps */

    /** The set of all keys. */
    Set<Key> keySet();

    /** The multiset of all values that can be returned by get.
     * (A multiset is a collection that may have duplicates). */
    Collection<Value> values();

    /** The set of all (key, value) pairs */
    Set<Map.Entry<Key, Value>> entrySet();
}
View Examples

Using example from a previous slide:

```java
Map<String, String> f = new TreeMap<String, String>();
f.put("Paul", "George"); f.put("George", "Martin");
f.put("Dana", "John");
```

we can take various views of `f`:

```java
for (Iterator<String> i = f.keySet().iterator(); i.hasNext();)
    i.next() ===> Dana, George, Paul
// or, more succinctly:
for (String name : f.keySet())
    name ===> Dana, George, Paul

for (String parent : f.values())
    parent ===> John, Martin, George

for (Map.Entry<String, String> pair : f.entrySet())
    pair ===> (Dana,John), (George,Martin), (Paul,George)

f.keySet().remove("Dana");  // Now f.get("Dana") == null
```
Simple Banking I: Accounts

Problem: Want a simple banking system. Can look up accounts by name or number, deposit or withdraw, print.

Account Structure

class Account {
    Account(String name, String number, int init) {
        this.name = name; this.number = number;
        this.balance = init;
    }
    /** Account-holder’s name */
    final String name;
    /** Account number */
    final String number;
    /** Current balance */
    int balance;

    /** Print THIS on STR in some useful format. */
    void print(PrintStream str) { ... } }
Simple Banking II: Banks

class Bank {
    /* These variables maintain mappings of String -> Account. They keep
    * the set of keys (Strings) in "compareTo" order, and the set of
    * values (Accounts) is ordered according to the corresponding keys. */
    SortedMap<String, Account> accounts = new TreeMap<String, Account>();
    SortedMap<String, Account> names = new TreeMap<String, Account>();

    void openAccount(String name, int initBalance) {
        Account acc =
            new Account(name, chooseNumber(), initBalance);
        accounts.put(acc.number, acc);
        names.put(name, acc);
    }

    void deposit(String number, int amount) {
        Account acc = accounts.get(number);
        if (acc == null) ERROR(...);
        acc.balance += amount;
    }
    // Likewise for withdraw.
}
Banks (continued): Iterating

**Printing out Account Data**

/** Print out all accounts sorted by number on STR. */
void printByAccount(PrintStream str) {
    // accounts.values() is the set of mapped-to values. Its
    // iterator produces elements in order of the corresponding keys.
    for (Account account : accounts.values())
        account.print(str);
}

/** Print out all bank accounts sorted by name on STR. */
void printByName(PrintStream str) {
    for (Account account : names.values())
        account.print(str);
}

A Design Question: What would be an appropriate representation for keeping a record of all transactions (deposits and withdrawals) against each account?
Partial Implementations

• Besides interfaces (like List) and concrete types (like LinkedList), Java library provides abstract classes such as AbstractList.

• Idea is to take advantage of the fact that operations are related to each other.

• Example: once you know how to do get(k) and size() for an implementation of List, you can implement all the other methods needed for a read-only list (and its iterators).

• Now throw in add(k,x) and you have all you need for the additional operations of a growable list.

• Add set(k,x) and remove(k) and you can implement everything else.
Example: The java.util.AbstractList helper class

```java
public abstract class AbstractList<Item> implements List<Item> {
    /** Inherited from List */
    // public abstract int size();
    // public abstract Item get(int k);
    public boolean contains(Object x) {
        for (int i = 0; i < size(); i += 1) {
            if ((x == null && get(i) == null) ||
                (x != null && x.equals(get(i))))
                return true;
        }
        return false;
    }
    return false;
}
/* OPTIONAL: Throws exception; override to do more. */
void add(int k, Item x) {
    throw new UnsupportedOperation();
}
Likewise for remove, set
```
Example, continued: AListIterator

// Continuing abstract class AbstractList<Item>:
public Iterator<Item> iterator() { return listIterator(); }
public ListIterator<Item> listIterator() {
    return new AListIterator(this);
}

private static class AListIterator implements ListIterator<Item> {
    AbstractList<Item> myList;
    AListIterator(AbstractList<Item> L) { myList = L; }
    /** Current position in our list. */
    int where = 0;

    public boolean hasNext() { return where < myList.size(); }
    public Item next() { where += 1; return myList.get(where-1); }
    public void add(Item x) { myList.add(where, x); where += 1; }
    ... previous, remove, set, etc.
}
...
Aside: Another way to do AListIterator

It's also possible to make the nested class non-static:

```java
public Iterator<Item> iterator() { return listIterator(); }
public ListIterator<Item> listIterator() { return this.new AListIterator(); }
```

```java
private class AListIterator implements ListIterator<Item> {
    /** Current position in our list. */
    int where = 0;

    public boolean hasNext() { return where < AbstractList.this.size(); }
    public Item next() { where += 1; return AbstractList.this.get(where-1); }
    public void add(Item x) { AbstractList.this.add(where, x); where += 1; }
    ... previous, remove, set, etc.
}
```

• Here, AbstractList.this means “the AbstractList I am attached to” and X.new AListIterator means "create a new AListIterator that is attached to X."

• In this case you can abbreviate this.new as new and can leave off some AbstractList.this parts, since meaning is unambiguous.
Example: Using AbstractList

Problem: Want to create a reversed view of an existing List (same elements in reverse order). Operations on the original list affect the view, and vice-versa.

```java
public class ReverseList<Item> extends AbstractList<Item> {
    private final List<Item> L;

    public ReverseList(List<Item> L) { this.L = L; }

    public int size() { return L.size(); }

    public Item get(int k) { return L.get(L.size() - k - 1); }

    public void add(int k, Item x) { L.add(L.size() - k, x); }

    public Item set(int k, Item x) { return L.set(L.size() - k - 1, x); }

    public Item remove(int k) { return L.remove(L.size() - k - 1); }
}
```
Getting a View: Sublists

Problem: \( L . \text{sublist}(\text{start}, \text{end}) \) is a List that gives a view of part of an existing list. Changes in one must affect the other. How?

```java
// Continuation of class AbstractList. Error checks not shown.
List<Item> sublist(int start, int end) {
    return this.new Sublist(start, end);
}

private class Sublist extends AbstractList<Item> {
    private int start, end;
    Sublist(int start, int end) { obvious }

    public int size() { return end-start; }
    public Item get(int k) { return AbstractList.this.get(start+k); }

    public void add(int k, Item x) {
        AbstractList.this.add(start+k, x); end += 1;
    }
    ...
}
```
What Does a Sublist Look Like?

• Consider \( SL = L\text{.sublist}(3, 5); \)
Arrays and Links

• Two main ways to represent a sequence: array and linked list
• In Java Library: ArrayList and Vector vs. LinkedList.

• Array:
  - Advantages: compact, fast ($\Theta(1)$) random access (indexing).
  - Disadvantages: insertion, deletion can be slow ($\Theta(N)$)

• Linked list:
  - Advantages: insertion, deletion fast once position found.
  - Disadvantages: space (link overhead), random access slow.
Implementing with Arrays

• Biggest problem using arrays is insertion/deletion in the middle of a list (must shove things over).

• Adding/deleting from ends can be made fast:
  – Double array size to grow; amortized cost constant (Lecture #15).
  – Growth at one end really easy; classical stack implementation:

    S.push("X");
    S.push("Y");
    S.push("Z");

  – To allow growth at either end, use circular buffering:

  – Random access still fast.
Essentials of linking should now be familiar

Used in Java LinkedList. One possible representation for linked list and an iterator object over it:

```java
L = new LinkedList<String>();
L.add("axolotl");
L.add("kludge");
L.add("xerophyte");

I = L.listIterator();
I.next();
```

Diagram:

```
L: | 3 |
   └── I: |
          | 1 |
             | lastReturned
          | here
             | nextIndex

L = new LinkedList<String>();
I = L.listIterator();
I.next();
```
Clever trick: Sentinels

- A sentinel is a dummy object containing no useful data except links.
- Used to eliminate special cases and to provide a fixed object to point to in order to access a data structure.
- Avoids special cases ('if' statements) by ensuring that the first and last item of a list always have (non-null) nodes—possibly sentinels—before and after them:

```c
// To delete list node at p:
p.next.prev = p.prev;
p.prev.next = p.next;

// To add new node N before p:
N.prev = p.prev; N.next = p;
p.prev.next = N;
p.prev = N;
```
Specialization

- Traditional special cases of general list:
  - **Stack**: Add and delete from one end (LIFO).
  - **Queue**: Add at end, delete from front (FIFO).
  - **Dequeue**: Add or delete at either end.

- All of these easily representable by either array (with circular buffering for queue or deque) or linked list.

- Java has the **List** types, which can act like any of these (although with non-traditional names for some of the operations).

- Also has `java.util.Stack`, a subtype of **List**, which gives traditional names ("push", "pop") to its operations. There is, however, no "stack" interface.
Stacks and Recursion

- Stacks related to recursion. In fact, can convert any recursive algorithm to stack-based (however, generally no great performance benefit):
  - Calls become “push current variables and parameters, set parameters to new values, and loop.”
  - Return becomes “pop to restore variables and parameters.”

```python
findExit(start):
    if isExit(start)
        FOUND
    else if (!isCrumb(start))
        leave crumb at start;
        for each square, x, adjacent to start:
            if legal(start,x) && !isCrumb(x)
                findExit(x)
```

Call: findExit((0,0))
Exit: (4, 2)

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findExit(start):
    S = new empty stack;
    push start on S;
    while S not empty:
        pop S into start;
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                    push x on S
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```
2, 3
3, 2
3, 1
1 2 5
4 3 6
```
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- Stacks related to recursion. In fact, can convert any recursive algorithm to stack-based (however, generally no great performance benefit):
  - Calls become "push current variables and parameters, set parameters to new values, and loop."
  - Return becomes "pop to restore variables and parameters."

```
findExit(start):
  if isExit(start)
      FOUND
  else if (!isCrumb(start))
      leave crumb at start;
      for each square, x,
          adjacent to start:
              if legal(start,x) && !isCrumb(x)
                  findExit(x)
```

Call: findExit((0,0))
Exit: (4, 2)

```
findExit(start):
  S = new empty stack;
  push start on S;
  while S not empty:
      pop S into start;
      if isExit(start)
          FOUND
      else if (!isCrumb(start))
          leave crumb at start;
          for each square, x,
              adjacent to start (in reverse):
                  if legal(start,x) && !isCrumb(x)
                      push x on S
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Design Choices: Extension, Delegation, Adaptation

- The standard `java.util.Stack` type `extends` Vector:

  ```java
class Stack<Item> extends Vector<Item> {
    void push(Item x) { add(x); } ...
  }
```

- Could instead have *delegated* to a field:

  ```java
class ArrayStack<Item> {
    private ArrayList<Item> repl = new ArrayList<Item>();
    void push(Item x) { repl.add(x); } ...
  }
```

- Or, could generalize, and define an *adapter*: a class used to make objects of one kind behave as another:

  ```java
public class StackAdapter<Item> {
  private List repl;
  /** A stack that uses REPL for its storage. */
  public StackAdapter(List<Item> repl) { this.repl = repl; }
  public void push(Item x) { repl.add(x); } ...
}
```

  ```java
class ArrayStack<Item> extends StackAdapter<Item> {
    ArrayStack() { super(new ArrayList<Item>()); }
  }
```