Researchers at Microsoft and UW are working on a system that uses the fact that your body can act as an antenna and notes how ambient electric fields change to figure out what your position or motion was. The idea is you don’t need a camera or Wiimote to interact with it!

Review

- Memory is byte-addressable, but lw and sw access one word at a time.
- A pointer (used by lw and sw) is just a memory address, so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using conditional statements within if, while, do while, for.
- MIPS Decision making instructions are the conditional branches: beq and bne.
- New Instructions:
  
  lw, sw, beq, bne, j
Last time: Loading, Storing bytes 1/2

- In addition to word data transfers (lw, sw), MIPS has byte data transfers:
  - load byte: lb
  - store byte: sb
- same format as lw, sw
- E.g., lb $s0, 3($s1)
  - contents of memory location with address = sum of “3” + contents of register s1 is copied to the low byte position of register s0.
Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
  - `lb`: sign extends to fill upper 24 bits

```
xxxx xxxz xxxz xxxz
```

...is copied to "sign-extend"

- Normally don’t want to sign extend chars
- MIPS instruction that doesn’t sign extend when loading bytes:
  - `load byte unsigned: lbu`
Overflow in Arithmetic (1/2)

- Reminder: Overflow occurs when there is a mistake in arithmetic due to the limited precision in computers.

- Example (4-bit unsigned numbers):

  \[
  \begin{array}{c}
  15 \\
  + 3 \\
  \hline
  18
  \end{array}
  \quad
  \begin{array}{c}
  1111 \\
  + 0011 \\
  \hline
  10010
  \end{array}
  \]

- But we don’t have room for 5-bit solution, so the solution would be 0010, which is +2, and wrong.
Overflow in Arithmetic (2/2)

- Some languages detect overflow (Ada), some don’t (C)

- MIPS solution is 2 kinds of arithmetic instructs:
  - These **cause overflow to be detected**
    - add (add)
    - add immediate (addi)
    - subtract (sub)
  - These **do not cause overflow detection**
    - add unsigned (addu)
    - add immediate unsigned (addiu)
    - subtract unsigned (subu)

- Compiler selects appropriate arithmetic
  - MIPS C compilers produce **addu, addiu, subu**
Two "Logic" Instructions

- Here are 2 more new instructions

- Shift Left: `s1l $s1,$s2,2` # $s1 = $s2 << 2
  - Store in $s1 the value from $s2 shifted 2 bits to the left (they fall off end), inserting 0's on right; `<<` in C.
  - Before: `0000 0002`_{hex}
    - `0000 0000 0000 0000 0000 0000 0000 0010`_{two}
  - After: `0000 0008`_{hex}
    - `0000 0000 0000 0000 0000 0000 0000 1000`_{two}
  - What arithmetic effect does shift left have?

- Shift Right: `srl` is opposite shift; `>>`
Loops in C/Assembly (1/3)

- Simple loop in C; \( A[\] \) is an array of \textit{ints}

  \[
  \text{do} \{ \quad g = g + A[i]; \\
  \quad i = i + j; \\
  \} \text{ while } (i != h);
  \]

- Rewrite this as:

  \[
  \text{Loop: } g = g + A[i]; \\
  \quad i = i + j; \\
  \quad \text{if } (i != h) \text{ goto Loop;}
  \]

- Use this mapping:

  \[
  g, \quad h, \quad i, \quad j, \quad \text{base of } A \\
  \quad $s1, \quad $s2, \quad $s3, \quad $s4, \quad $s5
  \]
Loops in C/Assembly (2/3)

- Final compiled MIPS code:

  ```
  Loop:  sll  $t1,$s3,2   # $t1= 4*I
         addu $t1,$t1,$s5   # $t1=addr A+4i
         lw   $t1,0($t1)    # $t1=A[i]
         addu $s1,$s1,$t1   # g=g+A[i]
         addu $s3,$s3,$s4   # i=i+j
         bne  $s3,$s2,Loop  # goto Loop
                     # if i!=h
  ```

- Original code:

  ```
  Loop: g = g + A[i];
         i = i + j;
         if (i != h) goto Loop;
  ```
Loops in C/Assembly (3/3)

- There are three types of loops in C:
  - `while`
  - `do ... while`
  - `for`

- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.

- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch
Administrivia

- The schedule through week 7 has been determined
  - Midterm 7-9pm on 2011-10-06
- Other administrivia?
Inequalities in MIPS (1/4)

- Until now, we’ve only tested equalities (== and != in C). General programs need to test < and > as well.
- Introduce MIPS Inequality Instruction:
  - “Set on Less Than”
  - Syntax: `slt reg1,reg2,reg3`
  - Meaning: `reg1 = (reg2 < reg3);`

```c
if (reg2 < reg3)
    reg1 = 1;
else reg1 = 0;
```

“set” means “change to 1”, “reset” means “change to 0”. Same thing...
Inequalities in MIPS (2/4)

- How do we use this? Compile by hand:
  \[
  \text{if (} g < h \text{) goto } \text{Less; } \#g:$s0, h:$s1
  \]

- Answer: compiled MIPS code…
  \[
  \text{slt } $t0,$s0,$s1 \quad \# \quad t0 = 1 \text{ if } g<h
  \\
  \text{bne } $t0,$0,Less \quad \# \quad \text{goto } \text{Less}
  \quad \# \quad \text{if } t0!=0
  \quad \# \quad (\text{if } (g<h)) \text{ Less:}
  \]

- Register $0$ always contains the value 0, so \text{bne} and \text{beq} often use it for comparison after an \text{slt} instruction.

- A \text{slt} \rightarrow \text{bne} pair means \text{if(... < ... goto...}
Inequalities in MIPS (3/4)

- Now we can implement $<$, but how do we implement $>$, $\leq$ and $\geq$?
- We could add 3 more instructions, but:
  - MIPS goal: Simpler is Better
- Can we implement $\leq$ in one or more instructions using just `slt` and branches?
  - What about $>$?
  - What about $\geq$?
# a:$s0, b:$s1

slt $t0,$s0,$s1  # $t0 = 1 if a<b
beq $t0,$0,skip  # skip if a >= b
<stuff>

# do if a<b

skip:

Two independent variations possible:

Use `slt $t0,$s1,$s0` instead of `slt $t0,$s0,$s1`

Use `bne` instead of `beq`
Immediates in Inequalities

- There is also an immediate version of `slt` to test against constants: `slti`
  - Helpful in `for` loops

```c
if (g >= 1) goto Loop
```

```mips
slti $t0,$s0,1  # $t0 = 1 if $s0<1 (g<1)
beq   $t0,$0,Loop  # goto Loop
       # if $t0==0 (if (g>=1))
```

An `slt` ➔ `beq` pair means if(... ≥ ...) goto...
What about unsigned numbers?

- Also **unsigned** inequality instructions: sltu, sltiu

  ...which sets result to 1 or 0 depending on unsigned comparisons

- What is value of $t0, t1$?

  ($s0 = \text{FFFF FFFA}_{\text{hex}}, s1 = \text{0000 FFFA}_{\text{hex}}$)

  slt $t0, s0, s1$
  sltu $t1, s0, s1$
MIPS Signed vs. Unsigned – diff meanings!

- MIPS terms Signed/Unsigned “overloaded”:
  - Do/Don’t sign extend
    - (lb, lbu)
  - Do/Don’t overflow
    - (add, addi, sub, mult, div)
    - (addu, addiu, subu, multu, divu)
  - Do signed/unsigned compare
    - (slt, slti/sltu, sltiu)
Peer Instruction

Loop:

```assembly
addi $s0, $s0, -1       # i = i - 1
slt  $t0, $s1, 2        # $t0 = (j < 2)
beq  $t0, $0 , Loop    # goto Loop if $t0 == 0
slt  $t0, $s1, $s0      # $t0 = (j < i)
bne  $t0, $0 , Loop    # goto Loop if $t0 != 0
```

($s0=i, $s1=j)

What C code properly fills in the blank in loop below?

```c
do {i--;} while(___);
```
“And in conclusion…”

- To help the conditional branches make decisions concerning inequalities, we introduce: “Set on Less Than” called `slt, slti, sltu, sltiu`
- One can store and load (signed and unsigned) bytes as well as words with `lb, lbu`
- Unsigned add/sub don’t cause overflow
- New MIPS Instructions:
  - `sll, srl, lb, lbu`
  - `slt, slti, sltu, sltiu`
  - `addu, addiu, subu`
Example: The C Switch Statement (1/3)

- Choose among four alternatives depending on whether \( k \) has the value 0, 1, 2 or 3. Compile this C code:

```c
switch (k) {
    case 0: f=i+j; break; /* k=0 */
    case 1: f=g+h; break; /* k=1 */
    case 2: f=g-h; break; /* k=2 */
    case 3: f=i-j; break; /* k=3 */
}
```
Example: The C Switch Statement (2/3)

- This is complicated, so simplify.
- Rewrite it as a chain of if-else statements, which we already know how to compile:

  ```c
  if(k==0) f=i+j;
      else if(k==1) f=g+h;
          else if(k==2) f=g-h;
              else if(k==3) f=i-j;
  ```

- Use this mapping:
  
  ```
  f:$s0, g:$s1, h:$s2,
  i:$s3, j:$s4, k:$s5
  ```
Example: The C Switch Statement (3/3)

- Final compiled MIPS code:

```
bne $s5,$0,L1       # branch k!=0
add $s0,$s3,$s4    # k==0 so f=i+j
j Exit            # end of case so Exit
L1: addi $t0,$s5,-1 # $t0=k-1
bne $t0,$0,L2      # branch k!=1
add $s0,$s1,$s2    # k==1 so f=g+h
j Exit            # end of case so Exit
L2: addi $t0,$s5,-2 # $t0=k-2
bne $t0,$0,L3      # branch k!=2
sub $s0,$s1,$s2    # k==2 so f=g-h
j Exit            # end of case so Exit
L3: addi $t0,$s5,-3 # $t0=k-3
bne $t0,$0,Exit    # branch k!=3
sub $s0,$0,$s3,$s4 # k==3 so f=i-j
Exit:
```