

inst.eecs.berkeley.edu/~cs61c
CS61C : Machine Structures

Lecture #2 – Number Representation

2007-01-19

There is one handout today at the front and back of the room!



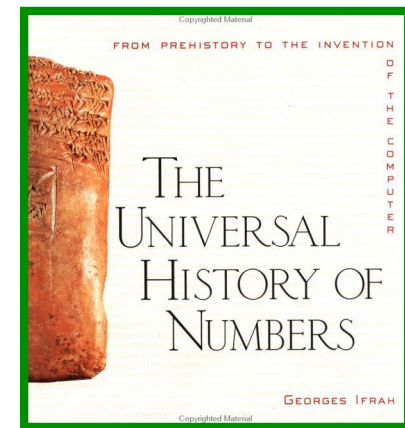
Lecturer SOE Dan Garcia

www.cs.berkeley.edu/~ddgarcia

Great book ⇒
**The Universal History
of Numbers**



by **Georges Ifrah**



Great DeCal courses I supervise (2 units)

- **UCBUGG**

- UC Berkeley Undergraduate Graphics Group
- Thursdays 5:30-7:30pm in 310 Soda
- Learn to create a short 3D animation
- No prereqs (but they might have too many students, so admission not guaranteed)
- <http://ucbugg.berkeley.edu>

- **MS-DOS X**

- Macintosh Software Developers for OS X
- Thursdays 5-7pm in 320 Soda
- Learn to program the Macintosh and write an awesome GUI application
- No prereqs (other than interest)
- <http://msdosx.berkeley.edu>



Review

- **Continued rapid improvement in computing**
 - **2X every 2.0 years in memory size;**
every 1.5 years in processor speed;
every 1.0 year in disk capacity;
 - **Moore's Law enables processor**
(2X transistors/chip ~1.5 yrs)
- **5 classic components of all computers**

Control Datapath Memory Input Output



Processor



My goal as an instructor

- **To make your experience in CS61C as enjoyable & informative as possible**
 - Humor, enthusiasm, graphics & technology-in-the-news in lecture
 - Fun, challenging projects & HW
 - Pro-student policies (exam clobbering)
- **To maintain Cal & EECS standards of excellence**
 - Your projects & exams will be just as rigorous as every year. Overall : B- avg
- **To be an HKN “7.0” man**
 - I know I speak fast when I get excited about material. I’m told every semester. Help me slow down when I go toooo fast.
 - Please give me feedback so I improve! Why am I not 7.0 for you? I will listen!!



Putting it all in perspective...

“If the automobile had followed the same development cycle as the computer,

– Robert X. Cringely



Decimal Numbers: Base 10

Digits: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9

Example:

3271 =

$$(3 \times 10^3) + (2 \times 10^2) + (7 \times 10^1) + (1 \times 10^0)$$



Numbers: positional notation

- **Number Base B \Rightarrow B symbols per digit:**

- **Base 10 (Decimal):** 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
- **Base 2 (Binary):** 0, 1

- **Number representation:**

- $d_{31}d_{30} \dots d_1d_0$ is a 32 digit number

- $\text{value} = d_{31} \times B^{31} + d_{30} \times B^{30} + \dots + d_1 \times B^1 + d_0 \times B^0$

- **Binary:** 0,1 (In binary digits called “bits”)



- $0b11010 = 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$
 $= 16 + 8 + 2$
#s often written $= 26$

0b... • Here 5 digit binary # turns into a 2 digit decimal #

- **Can we find a base that converts to binary easily?**



Hexadecimal Numbers: Base 16

- **Hexadecimal:**
0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F
 - Normal digits + 6 more from the alphabet
 - In C, written as **0x...** (e.g., 0xFAB5)
- **Conversion: Binary \Leftrightarrow Hex**
 - 1 hex digit represents 16 decimal values
 - 4 binary digits represent 16 decimal values
 - \Rightarrow 1 hex digit replaces 4 binary digits
- One hex digit is a “**nibble**”. Two is a “**byte**”
- **Example:**
 - 1010 1100 0011 (binary) = **0x_____** ?



Decimal vs. Hexadecimal vs. Binary

Examples:

1010 1100 0011 (binary)
= 0xAC3

10111 (binary)
= 0001 0111 (binary)
= 0x17

0x3F9
= 11 1111 1001 (binary)

*How do we convert between
hex and Decimal?*

00	0	0000
01	1	0001
02	2	0010
03	3	0011
04	4	0100
05	5	0101
06	6	0110
07	7	0111
08	8	1000
09	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111

MEMORIZE!



Kilo, Mega, Giga, Tera, Peta, Exa, Zetta, Yotta

physics.nist.gov/cuu/Units/binary.html

- Common use prefixes (all SI, except K [= k in SI])

Name	Abbr	Factor	SI size
Kilo	K	$2^{10} = 1,024$	$10^3 = 1,000$
Mega	M	$2^{20} = 1,048,576$	$10^6 = 1,000,000$
Giga	G	$2^{30} = 1,073,741,824$	$10^9 = 1,000,000,000$
Tera	T	$2^{40} = 1,099,511,627,776$	$10^{12} = 1,000,000,000,000$
Peta	P	$2^{50} = 1,125,899,906,842,624$	$10^{15} = 1,000,000,000,000,000$
Exa	E	$2^{60} = 1,152,921,504,606,846,976$	$10^{18} = 1,000,000,000,000,000,000$
Zetta	Z	$2^{70} = 1,180,591,620,717,411,303,424$	$10^{21} = 1,000,000,000,000,000,000,000$
Yotta	Y	$2^{80} = 1,208,925,819,614,629,174,706,176$	$10^{24} = 1,000,000,000,000,000,000,000,000$

- Confusing! Common usage of “kilobyte” means 1024 bytes, but the “correct” SI value is 1000 bytes
- **Hard Disk** manufacturers & **Telecommunications** are the only computing groups that use SI factors, so what is advertised as a 30 GB drive will actually only hold about 28×2^{30} bytes, and a 1 Mbit/s connection transfers 10^6 bps.



kibi, mebi, gibi, tebi, pebi, exbi, zebi, yobi

en.wikipedia.org/wiki/Binary_prefix

- **New IEC Standard Prefixes [only to exbi officially]**

Name	Abbr	Factor
kibi	Ki	$2^{10} = 1,024$
mebi	Mi	$2^{20} = 1,048,576$
gibi	Gi	$2^{30} = 1,073,741,824$
tebi	Ti	$2^{40} = 1,099,511,627,776$
pebi	Pi	$2^{50} = 1,125,899,906,842,624$
exbi	Ei	$2^{60} = 1,152,921,504,606,846,976$
zebi	Zi	$2^{70} = 1,180,591,620,717,411,303,424$
yobi	Yi	$2^{80} = 1,208,925,819,614,629,174,706,176$

As of this writing, this proposal has yet to gain widespread use...

- **International Electrotechnical Commission (IEC) in 1999 introduced these to specify binary quantities.**
 - **Names come from shortened versions of the original SI prefixes (same pronunciation) and *bi* is short for “binary”, but pronounced “bee” :-)**
 - **Now SI prefixes only have their base-10 meaning and never have a base-2 meaning.**

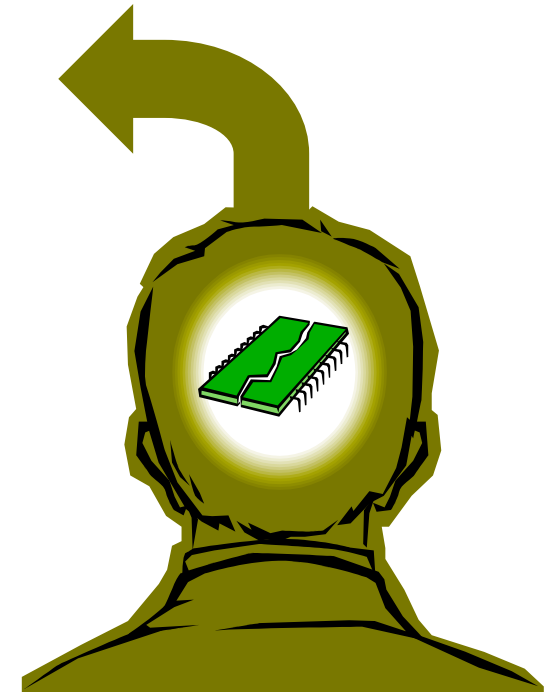


The way to remember #s

- What is 2^{34} ? How many bits addresses (i.e., what's $\text{ceil } \log_2 = \lg$ of) 2.5 TiB?

- Answer! 2^{XY} means...

X=0	⇒	---	Y=0	⇒	1
X=1	⇒	kibi $\sim 10^3$	Y=1	⇒	2
X=2	⇒	mebi $\sim 10^6$	Y=2	⇒	4
X=3	⇒	gibi $\sim 10^9$	Y=3	⇒	8
X=4	⇒	tebi $\sim 10^{12}$	Y=4	⇒	16
X=5	⇒	pebi $\sim 10^{15}$	Y=5	⇒	32
X=6	⇒	exbi $\sim 10^{18}$	Y=6	⇒	64
X=7	⇒	zebi $\sim 10^{21}$	Y=7	⇒	128
X=8	⇒	yobi $\sim 10^{24}$	Y=8	⇒	256
			Y=9	⇒	512



MEMORIZE!



What to do with representations of numbers?

- **Just what we do with numbers!**

- Add them
- Subtract them
- Multiply them
- Divide them
- Compare them

$$\begin{array}{r} 11 \\ 1010 \\ + 0111 \\ \hline 10001 \end{array}$$

- **Example: $10 + 7 = 17$**

- ...so simple to add in binary that we can build circuits to do it!
- subtraction just as you would in decimal
- Comparison: How do you tell if $X > Y$?



Which base do we use?

- **Decimal:** great for humans, especially when doing arithmetic
- **Hex:** if human looking at long strings of binary numbers, its much easier to convert to hex and look 4 bits/symbol
 - Terrible for arithmetic on paper
- **Binary:** what computers use; you will learn how computers do +, -, *, /
 - To a computer, numbers always binary
 - Regardless of how number is written:
 - $32_{\text{ten}} == 32_{10} == 0x20 == 100000_2 == 0b100000$
 - Use subscripts “ten”, “hex”, “two” in book, slides when might be confusing



BIG IDEA: Bits can represent anything!!

- **Characters?**

- 26 letters \Rightarrow 5 bits ($2^5 = 32$)
- upper/lower case + punctuation \Rightarrow 7 bits (in 8) (“ASCII”)
- standard code to cover all the world’s languages \Rightarrow 8,16,32 bits (“Unicode”) www.unicode.com



- **Logical values?**

- 0 \Rightarrow False, 1 \Rightarrow True

- colors ? Ex: Red (00) Green (01) Blue (11)

- locations / addresses? commands?

- **MEMORIZE:** N bits \Leftrightarrow at most 2^N things



How to Represent Negative Numbers?

- So far, **unsigned numbers**
- Obvious solution: define leftmost bit to be sign!
 - $0 \Rightarrow +, 1 \Rightarrow -$
 - Rest of bits can be numerical value of number
- Representation called **sign and magnitude**
- MIPS uses 32-bit integers. $+1_{\text{ten}}$ would be:
0000 0000 0000 0000 0000 0000 0000 0001
- And -1_{ten} in sign and magnitude would be:
1000 0000 0000 0000 0000 0000 0000 0001



Shortcomings of sign and magnitude?

- **Arithmetic circuit complicated**
 - Special steps depending whether signs are the same or not
- **Also, two zeros**
 - $0x00000000 = +0_{\text{ten}}$
 - $0x80000000 = -0_{\text{ten}}$
 - What would two 0s mean for programming?
- **Therefore sign and magnitude abandoned**



Administrivia

- **Upcoming lectures**
 - **Next three lectures: Introduction to C**
- **Lab overcrowding**
 - **Remember, you can go to ANY discussion (none, or one that doesn't match with lab, or even more than one if you want)**
 - **Overcrowded labs - consider finishing at home and getting checkoffs in lab, or bringing laptop to lab**
- **HW**
 - **HW0 due in discussion next week**
 - **HW1 due this Wed @ 23:59 PST**
 - **HW2 due following Wed @ 23:59 PST**
- **Reading**
 - **K&R Chapters 1-6 (lots, get started now!); 1st quiz due Sun!**
- **Soda locks doors @ 6:30pm & on weekends**
- **Look at class website, newsgroup often!**

[http://inst.eecs.berkeley.edu/~cs61c/
ucb.class.cs61c](http://inst.eecs.berkeley.edu/~cs61c/ucb.class.cs61c)



Shortcomings of One's complement?

- **Arithmetic still a somewhat complicated.**
- **Still two zeros**
 - $0x00000000 = +0_{\text{ten}}$
 - $0xFFFFFFFF = -0_{\text{ten}}$
- **Although used for awhile on some computer products, one's complement was eventually abandoned because another solution was better.**

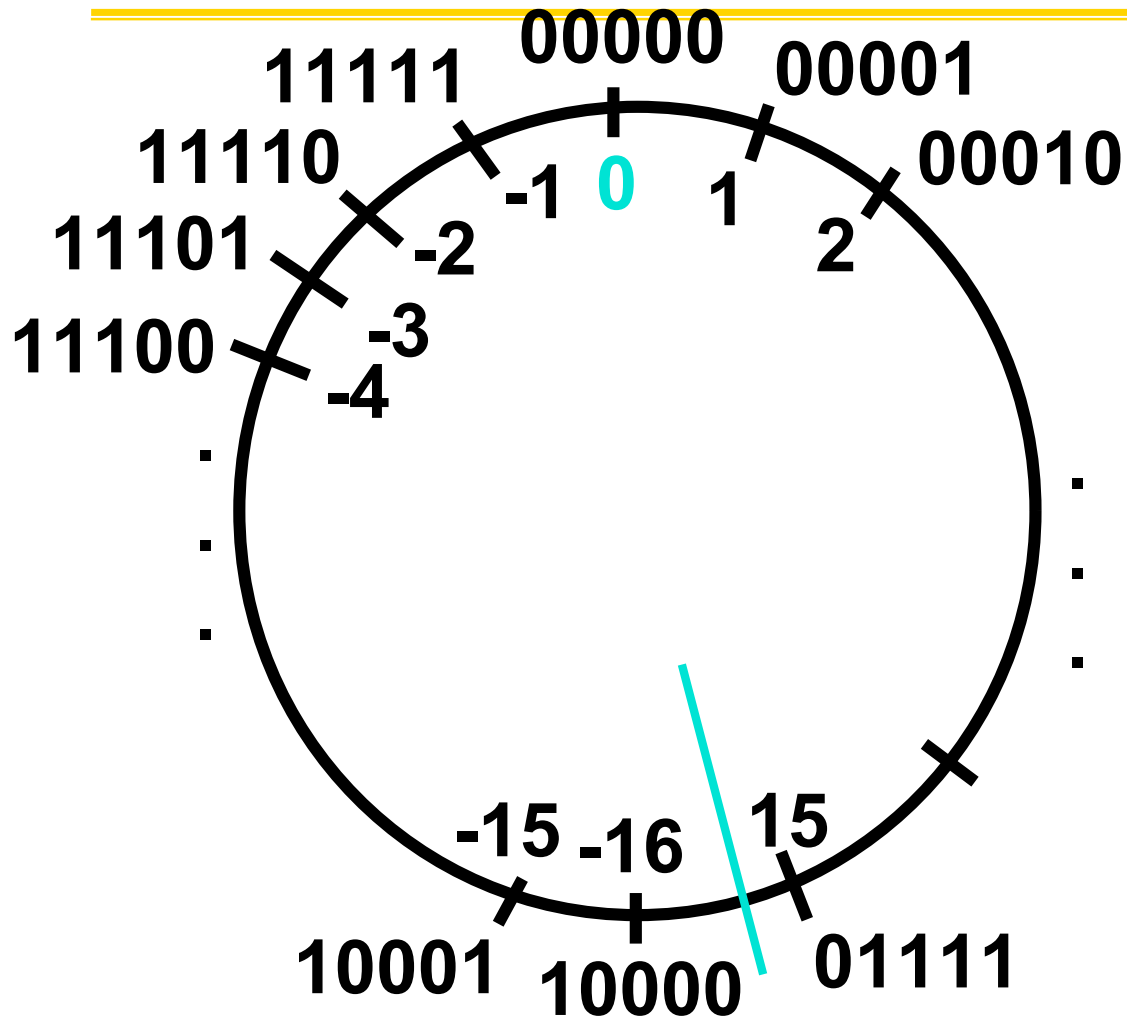


Standard Negative Number Representation

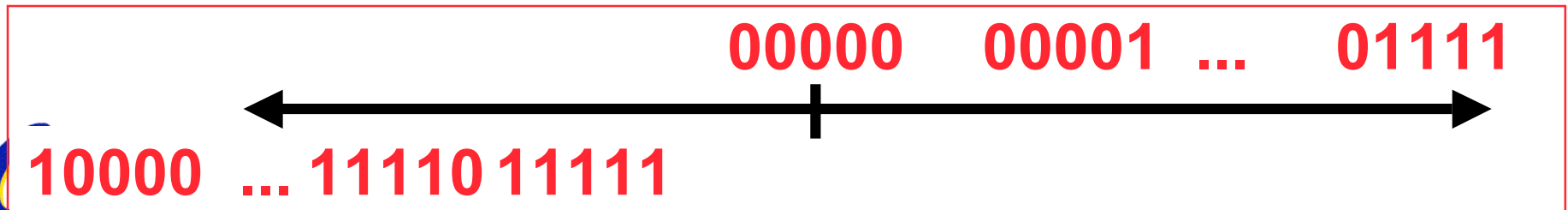
- What is result for unsigned numbers if tried to subtract large number from a small one?
 - Would try to borrow from string of leading 0s, so result would have a string of leading 1s
 - $3 - 4 \Rightarrow 00\dots0011 - 00\dots0100 = 11\dots1111$
 - With no obvious better alternative, pick representation that made the hardware simple
 - As with sign and magnitude, leading 0s \Rightarrow positive, leading 1s \Rightarrow negative
 - $000000\dots xxx$ is ≥ 0 , $111111\dots xxx$ is < 0
 - except $1\dots1111$ is -1, not -0 (as in sign & mag.)
- This representation is Two's Complement



2's Complement Number "line": N = 5



- 2^{N-1} non-negatives
- 2^{N-1} negatives
- **one zero**
- how many positives?



Two's Complement for N=32

0000 ... 0000 0000 0000 0000	$_{two} =$	0_{ten}
0000 ... 0000 0000 0000 0001	$_{two} =$	1_{ten}
0000 ... 0000 0000 0000 0010	$_{two} =$	2_{ten}
...		
0111 ... 1111 1111 1111 1101	$_{two} =$	$2,147,483,645_{ten}$
0111 ... 1111 1111 1111 1110	$_{two} =$	$2,147,483,646_{ten}$
0111 ... 1111 1111 1111 1111	$_{two} =$	$2,147,483,647_{ten}$
1000 ... 0000 0000 0000 0000	$_{two} =$	$-2,147,483,648_{ten}$
1000 ... 0000 0000 0000 0001	$_{two} =$	$-2,147,483,647_{ten}$
1000 ... 0000 0000 0000 0010	$_{two} =$	$-2,147,483,646_{ten}$
...		
1111 ... 1111 1111 1111 1101	$_{two} =$	-3_{ten}
1111 ... 1111 1111 1111 1110	$_{two} =$	-2_{ten}
1111 ... 1111 1111 1111 1111	$_{two} =$	-1_{ten}

- One zero; 1st bit called **sign bit**
- 1 “extra” negative: no positive $2,147,483,648_{ten}$



Two's Complement Formula

- Can represent positive and negative numbers in terms of the bit value times a power of 2:

$$d_{31} \times \underbrace{-(2^{31})}_{\text{red circle}} + d_{30} \times 2^{30} + \dots + d_2 \times 2^2 + d_1 \times 2^1 + d_0 \times 2^0$$

- Example: 1101_{two}

$$= 1 \times -(2^3) + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$$

$$= -2^3 + 2^2 + 0 + 2^0$$

$$= -8 + 4 + 0 + 1$$

$$= -8 + 5$$

$$= -3_{\text{ten}}$$



Two's Complement shortcut: Negation

*Check out www.cs.berkeley.edu/~dsw/twos_complement.html

- Change every 0 to 1 and 1 to 0 (invert or complement), then add 1 to the result

- Proof*: Sum of number and its (one's) complement must be $111\dots111_{\text{two}}$

However, $111\dots111_{\text{two}} = -1_{\text{ten}}$

Let $x' \Rightarrow$ one's complement representation of x

$$\text{Then } x + x' = -1 \Rightarrow x + x' + 1 = 0 \Rightarrow \boxed{-x = x' + 1}$$

- Example: -3 to +3 to -3

x :	1111	1111	1111	1111	1111	1111	1111	1101	<small>two</small>
x' :	0000	0000	0000	0000	0000	0000	0000	0010	<small>two</small>
+1 :	0000	0000	0000	0000	0000	0000	0000	0011	<small>two</small>
()' :	1111	1111	1111	1111	1111	1111	1111	1100	<small>two</small>
+1 :	1111	1111	1111	1111	1111	1111	1111	1101	<small>two</small>



You should be able to do this in your head...

Two's comp. shortcut: Sign extension

- Convert 2's complement number rep. using n bits to more than n bits
- Simply **replicate** the most significant bit (sign bit) of smaller to fill new bits
 - 2's comp. positive number has infinite 0s
 - 2's comp. negative number has infinite 1s
 - Binary representation hides leading bits; sign extension restores some of them
 - 16-bit -4_{ten} to 32-bit:

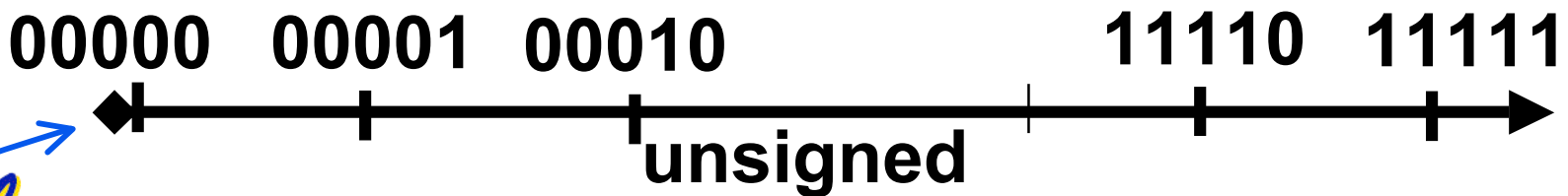
1111 1111 1111 1100_{two}

1111 1111 1111 1111 1111 1111 1111 1100_{two}



What if too big?

- Binary bit patterns above are simply **representatives** of numbers. Strictly speaking they are called “numerals”.
- Numbers really have an ∞ number of digits
 - with almost all being same (00...0 or 11...1) except for a few of the rightmost digits
 - Just don't normally show leading digits
- If result of add (or -, *, /) cannot be represented by these rightmost HW bits, **overflow** is said to have occurred.



Peer Instruction Question

$X = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1100_{\text{two}}$

$Y = 0011\ 1011\ 1001\ 1010\ 1000\ 1010\ 0000\ 0000_{\text{two}}$

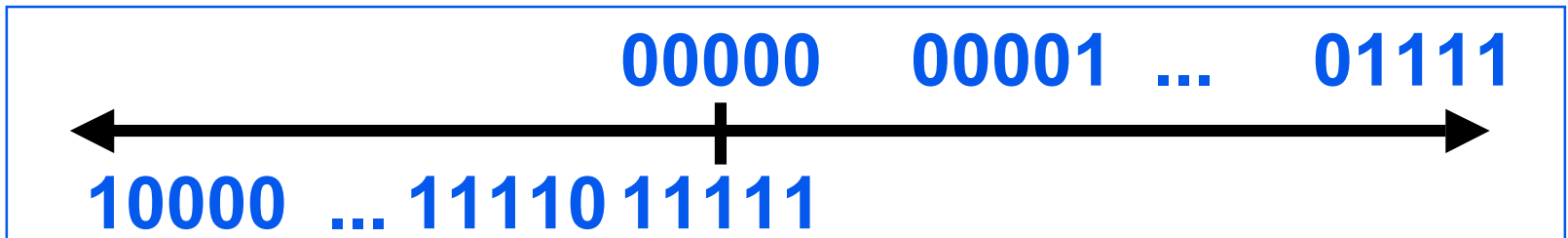
- A. $X > Y$ (if **signed**)
- B. $X > Y$ (if **unsigned**)
- C. An encoding for Babylonians could have 2^N non-negative numbers w/N bits!

	ABC
0:	FFF
1:	FFT
2:	FTF
3:	FTT
4:	TFF
5:	TFT
6:	TF
7:	TTT

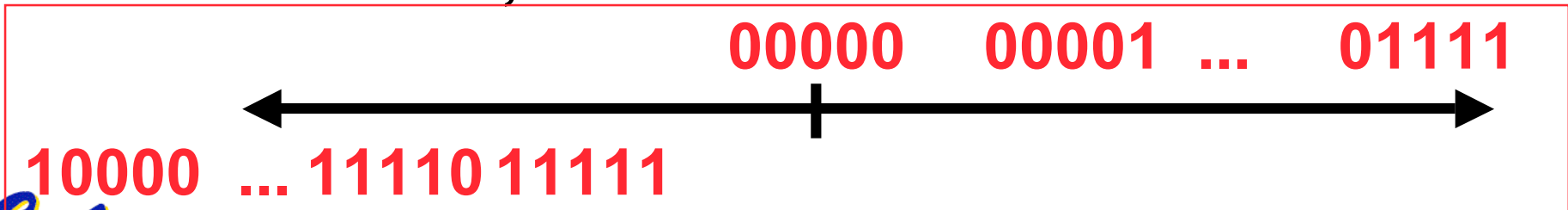


Number summary...

- We represent “things” in computers as particular bit patterns: $N \text{ bits} \Rightarrow 2^N$
- Decimal for human calculations, binary for computers, hex to write binary more easily
- **1's complement** - mostly abandoned



- **2's complement** universal in computing: cannot avoid, so learn



• **Overflow: numbers ∞ ; computers finite, errors!**