

## Lecture 3 – Introduction to the C Programming Language (pt 1)



**2007-01-22**

There is one handout  
today at the front and  
back of the room!

**Lecturer SOE Dan Garcia**

[www.cs.berkeley.edu/~ddgarcia](http://www.cs.berkeley.edu/~ddgarcia)

**HP overcomes Moore's Law? ⇒**

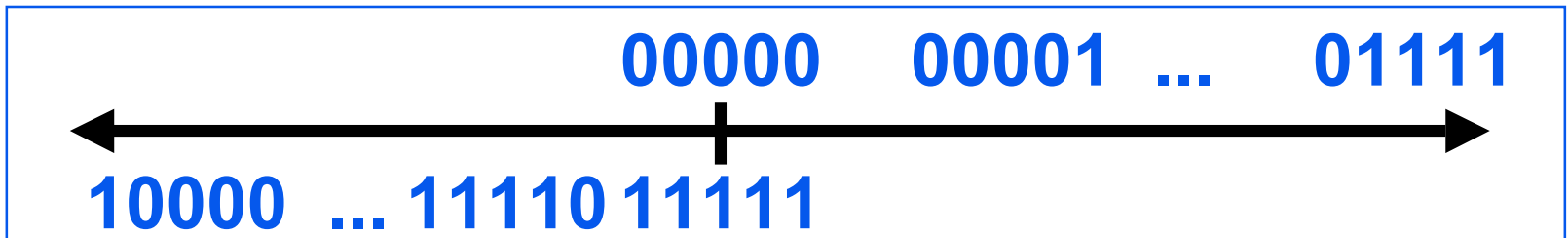
**Their design: “field programmable nanowire interconnect (FPNI)” uses a technique that “will enable chip makers to pack eight times as many transistors as is currently possible on a standard 45nm field programmable gate array (FPGA) chip.”**



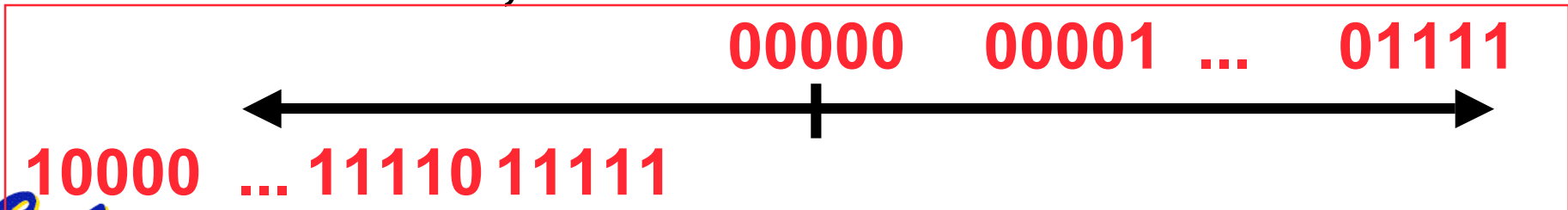
[hardware.slashdot.org/hardware/07/01/17/1333232.shtml](http://hardware.slashdot.org/hardware/07/01/17/1333232.shtml)

# Number review...

- We represent “things” in computers as particular bit patterns:  $N \text{ bits} \Rightarrow 2^N$
- Decimal for human calculations, binary for computers, hex to write binary more easily
- **1's complement** - mostly abandoned



- **2's complement** universal in computing: cannot avoid, so learn



• **Overflow: numbers  $\infty$ ; computers finite, errors!**

# Two's Complement shortcut: Negation

\*Check out [www.cs.berkeley.edu/~dsw/twos\\_complement.html](http://www.cs.berkeley.edu/~dsw/twos_complement.html)

- Change every 0 to 1 and 1 to 0 (invert or complement), then add 1 to the result

- Proof\*: Sum of number and its (one's) complement must be  $111\dots111_{\text{two}}$

However,  $111\dots111_{\text{two}} = -1_{\text{ten}}$

Let  $x' \Rightarrow$  one's complement representation of  $x$

$$\text{Then } x + x' = -1 \Rightarrow x + x' + 1 = 0 \Rightarrow \boxed{-x = x' + 1}$$

- Example: -3 to +3 to -3

x :	1111	1111	1111	1111	1111	1111	1111	1101	<small>two</small>
x' :	0000	0000	0000	0000	0000	0000	0000	0010	<small>two</small>
+1 :	0000	0000	0000	0000	0000	0000	0000	0011	<small>two</small>
( )' :	1111	1111	1111	1111	1111	1111	1111	1100	<small>two</small>
+1 :	1111	1111	1111	1111	1111	1111	1111	1101	<small>two</small>



You should be able to do this in your head...

# Two's comp. shortcut: Sign extension

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- Convert 2's complement number rep. using  $n$  bits to more than  $n$  bits
- Simply **replicate** the most significant bit (sign bit) of smaller to fill new bits
  - 2's comp. positive number has infinite 0s
  - 2's comp. negative number has infinite 1s
  - Binary representation hides leading bits; sign extension restores some of them
  - 16-bit  $-4_{\text{ten}}$  to 32-bit:

1111 1111 1111 1100<sub>two</sub>

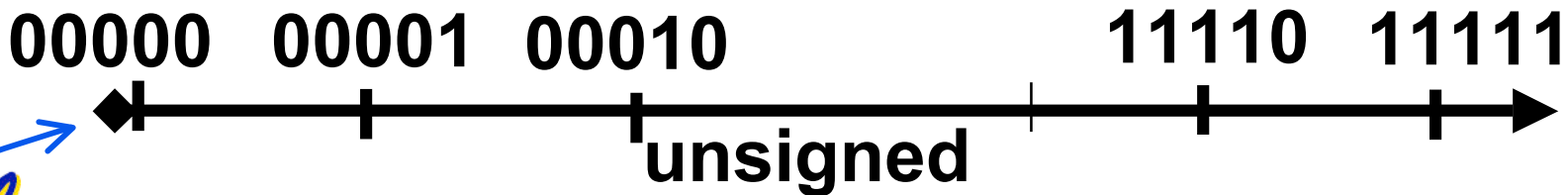
1111 1111 1111 1111 1111 1111 1111 1100<sub>two</sub>



# What if too big?

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- Binary bit patterns above are simply **representatives** of numbers. Strictly speaking they are called “numerals”.
- Numbers really have an  $\infty$  number of digits
  - with almost all being same (00...0 or 11...1) except for a few of the rightmost digits
  - Just don't normally show leading digits
- If result of add (or -, \*, /) cannot be represented by these rightmost HW bits, **overflow** is said to have occurred.



# Peer Instruction Question

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$X = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1100_{\text{two}}$

$Y = 0011\ 1011\ 1001\ 1010\ 1000\ 1010\ 0000\ 0000_{\text{two}}$

- A.  $X > Y$  (if **signed**)
- B.  $X > Y$  (if **unsigned**)
- C. An encoding for Babylonians could have  $2^N$  non-negative numbers w/N bits!

	ABC
0 :	<b>FFF</b>
1 :	<b>FFT</b>
2 :	<b>FTF</b>
3 :	<b>FTT</b>
4 :	<b>TFF</b>
5 :	<b>TFT</b>
6 :	<b>TF</b> <b>F</b>
7 :	<b>TTT</b>



# Introduction to C

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# Has there been an update to ANSI C?

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- **Yes! It's called the "C99" or "C9x" std**
  - You need "gcc -std=c99" to compile

- **References**

<http://en.wikipedia.org/wiki/C99>

[http://home.tiscalinet.ch/t\\_wolf/tw/c/c9x\\_changes.html](http://home.tiscalinet.ch/t_wolf/tw/c/c9x_changes.html)

- **Highlights**

- **Declarations anywhere, like Java (#15)**
- **Java-like // comments (to end of line) (#10)**
- **Variable-length non-global arrays (#33)**
- **<inttypes.h>: explicit integer types (#38)**
- **<stdbool.h> for boolean logic def's (#35)**
- **restrict keyword for optimizations (#30)**





# Disclaimer

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- **Important:** You will not learn how to fully code in C in these lectures! You'll still need your C reference for this course.
  - **K&R is a must-have reference**
    - Check online for more sources
  - **“JAVA in a Nutshell,” O'Reilly.**
    - Chapter 2, “How Java Differs from C”
  - **Brian Harvey's course notes**
    - On class website



# Compilation : Overview

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**C compilers** take C and convert it into an **architecture specific** machine code (string of 1s and 0s).

- Unlike Java which converts to **architecture independent** bytecode.
- Unlike most Scheme environments which interpret the code.
- These differ mainly in **when** your program is converted to machine instructions.
- For C, generally a 2 part process of **compiling** .c files to .o files, then **linking** the .o files into executables



# Compilation : Advantages

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- **Great run-time performance:** generally much faster than Scheme or Java for comparable code (because it optimizes for a given architecture)
- **OK compilation time:** enhancements in compilation procedure (`Makefiles`) allow only modified files to be recompiled



# Compilation : Disadvantages

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- All compiled files (including the executable) are **architecture specific**, depending on *both* the CPU type and the operating system.
- Executable must be **rebuilt** on each new system.
  - Called “**porting your code**” to a new architecture.
- The “change→compile→run [repeat]” iteration cycle is slow



# C vs. Java™ Overview (1/2)

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## Java

- Object-oriented (OOP)
- “Methods”
- Class libraries of data structures
- Automatic memory management

## C

- No built-in object abstraction. Data separate from methods.
- “Functions”
- C libraries are lower-level
- **Manual** memory management
- **Pointers**



# C vs. Java™ Overview (2/2)

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## Java

- **High** memory overhead from class libraries
- **Relatively Slow**
- Arrays initialize to **zero**
- **Syntax:**

```
/* comment */  
// comment  
System.out.print
```

## C

- **Low** memory overhead
- **Relatively Fast**
- Arrays initialize to **garbage**
- **Syntax: \***

```
/* comment */  
// comment  
printf
```

**\*You need newer C compilers to allow Java style comments, or just use C99**



# C Syntax: Variable Declarations

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- Very similar to Java, but with a few minor but important differences
- All variable declarations must go before they are used (at the beginning of the block)\*
- A variable may be initialized in its declaration.
- Examples of declarations:

- **correct:** {

```
int a = 0, b = 10;
```

```
...
```

- **Incorrect:**\* `for (int i = 0; i < 10; i++)`



\*C99 overcomes these limitations

# C Syntax: True or False?

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- **What evaluates to FALSE in C?**
  - 0 (integer)
  - NULL (pointer: more on this later)
  - no such thing as a Boolean\*
- **What evaluates to TRUE in C?**
  - **everything else...**
  - (same idea as in scheme: only #f is false, everything else is true!)



\*Boolean types provided by C99's `stdbool.h`



# C syntax : flow control

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- Within a function, remarkably **close to Java** constructs in methods (shows its legacy) in terms of flow control
  - `if-else`
  - `switch`
  - `while` and `for`
  - `do-while`



## C Syntax: main

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- To get the main function to accept arguments, use this:

```
int main (int argc, char *argv[])
```

- What does this mean?
  - `argc` will contain the number of strings on the command line (the executable counts as one, plus one for each argument).
    - Example: `unix% sort myFile`
  - `argv` is a pointer to an array containing the arguments as strings (more on pointers later).



# Administrivia

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- **Upcoming lectures**
  - C pointers and arrays in detail
- **HW**
  - HW0 due in discussion next week
  - HW1 due next Wed @ 23:59 PST
  - HW2 due following Wed @ 23:59 PST
- **Reading**
  - K&R Chapters 1-5 (lots, get started now!)
  - First quiz due Sun
- **Email me Ki - Me - Gi - ... mnemonics!**
  - The subject should be “kibi mebi gibi acronym”



# Address vs. Value

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- **Consider memory to be a single huge array:**
  - Each cell of the array has an address associated with it.
  - Each cell also stores some value.
  - Do you think they use signed or unsigned numbers? Negative address?!
- Don't confuse the **address** referring to a memory location with the **value** stored in that location.

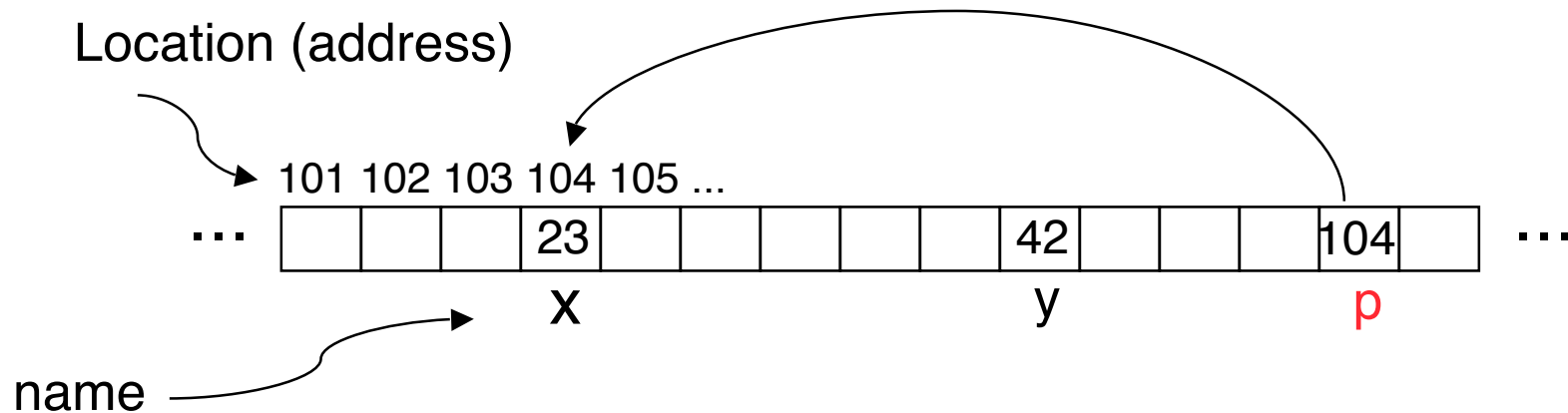
101 102 103 104 105 ...



# Pointers

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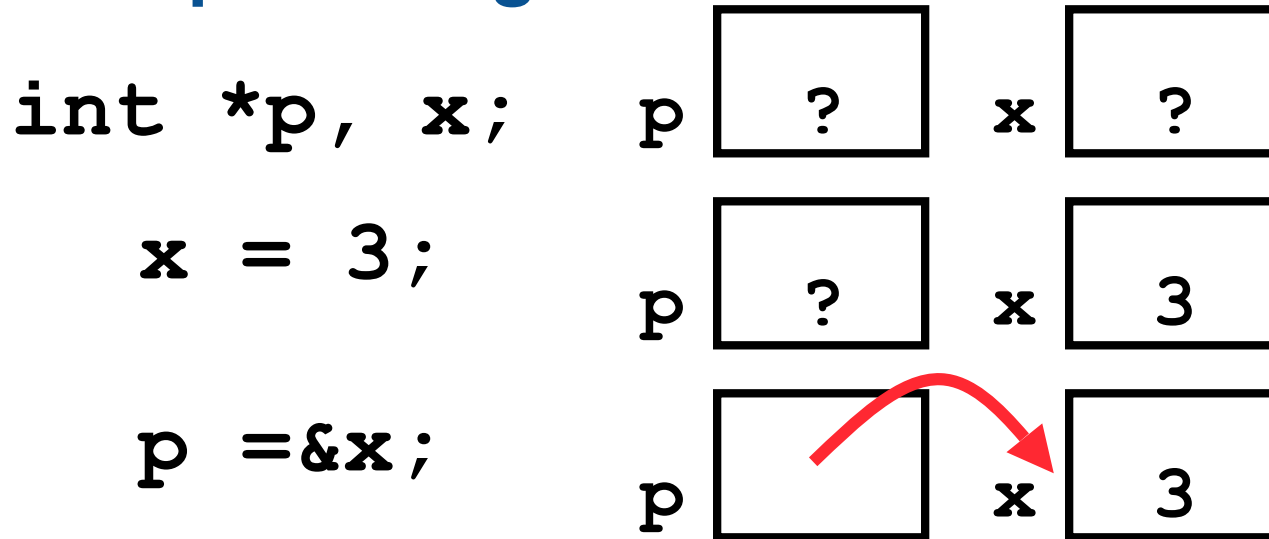
- An address refers to a particular memory location. In other words, it points to a memory location.
- **Pointer**: A variable that contains the address of a variable.



# Pointers

- How to create a pointer:

**& operator: get address of a variable**



Note the “\*” gets used 2 different ways in this example. In the declaration to indicate that `p` is going to be a pointer, and in the `printf` to get the value pointed to by `p`.

- How get a value pointed to?

\* “dereference operator”: get value pointed to

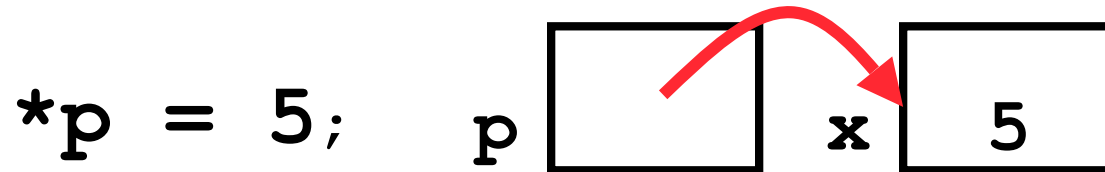
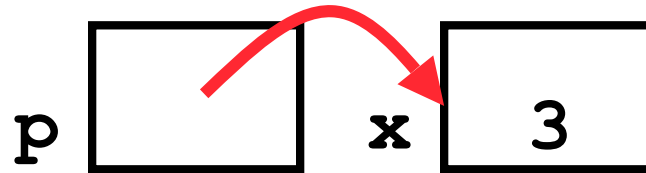
```
printf("p points to %d\n", *p);
```



# Pointers

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- How to change a variable pointed to?
  - Use dereference `*` operator on left of `=`



# Pointers and Parameter Passing

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- **Java and C pass parameters “by value”**
  - **procedure/function/method gets a copy of the parameter, so changing the copy cannot change the original**

```
void addOne (int x) {  
    x = x + 1;  
}
```

```
int y = 3;  
addOne (y) ;
```

**y is still = 3**





# Pointers and Parameter Passing

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- How to get a function to change a value?

```
void addOne (int *p) {  
    *p = *p + 1;  
}
```

```
int y = 3;
```

```
addOne (&y) ;
```

**y is now = 4**



# Pointers

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- Pointers are used to point to **any** data type (`int`, `char`, a `struct`, etc.).
- Normally a pointer can only point to one type (`int`, `char`, a `struct`, etc.).
  - `void *` is a type that can point to anything (generic pointer)
  - Use sparingly to help avoid program bugs... and security issues... and a lot of other bad things!



# Peer Instruction Question

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```
void main(); {
    int *p, x=5, y; // init
    y = *(p = &x) + 10;
    int z;
    flip-sign(p);
    printf("x=%d,y=%d,p=%d\n", x, y, p);
}
flip-sign(int *n){*n = -(*n)}
```

How many syntax/logic errors?

<u>#Errors</u>
0
1
2
3
4
5
6
7
8
9



# Peer Instruction Answer

---

```
void main(); {
    int *p, x=5, y; // init
    y = *(p = &x) + 10;
    int z;
    flip-sign(p);
    printf("x=%d,y=%d,p=%d\n", x, y, *p);
}
flip-sign(int *n){*n = -(*n);}
```

How many syntax/logic errors? I get **5**.  
(signed printing of pointer illogical)

#Errors

0  
1  
2  
3  
4  
5  
6  
7  
8  
9



## And in conclusion...

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- All declarations go at the beginning of each function.
- Only 0 and NULL evaluate to FALSE.
- All data is in memory. Each memory location has an address to use to refer to it and a value stored in it.
- A **pointer** is a C version of the address.
  - \* “follows” a pointer to its value
  - & gets the address of a value

