

Lecture 37
I/O : Networks

2007-04-20



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inst.eecs/~cs61c-te

Bad software upgrade causes service outage for BlackBerry users

Software upgrade to improve Cache performance triggered a problem at RIM's hub for North American traffic... effecting more than five million users

<http://www.nytimes.com/2007/04/20/technology/>



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I/O Review

- I/O gives computers their **5 senses**
- I/O speed range is **12.5-million to one**
- Differences in processor and I/O speed → synchronize with I/O devices before use
- **Polling** works, but expensive
 - processor *repeatedly queries devices*
- **Interrupts** works, more complex
 - device *causes an exception*, causing OS to run and deal with the device
- I/O control leads to **Operating Systems**



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Why Networks?

- Originally *sharing I/O devices* between **computers**
 ex: printers
- Then *communicating* between **computers**
 ex: file transfer protocol
- Then *communicating* between **people**
 ex: e-mail
- Then *communicating* between **networks of computers**
 ex: file sharing, www, ...



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How Big is the Network (2007)?

- ~30 in 273 Soda
- ~525 in inst.cs.berkeley.edu
- ~6,400 in eecs & cs .berkeley.edu
- (1999) ~50,000 in berkeley.edu
- ~10,000,000 in .edu (2005: ~9,000,000)
- ~258,941,310 in US (2005: ~217,000,000, 2006: ~286.5E6)
 (.net .com .edu .arpa .us .mil .org .gov)
- ~433,190,000 in the world
 (2005: ~317,000,000, 2006: ~439,000,000)

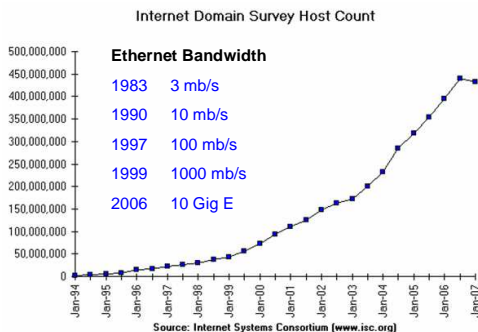


Source: Internet Software Consortium: www.isc.org

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Growth Rate



en.wikipedia.org/wiki/10_gigabit_ethernet

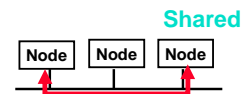
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Shared vs. Switched Based Networks

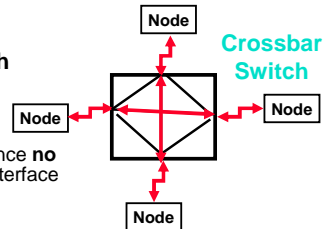
Shared vs. Switched:

- **Switched:** pairs ("point-to-point" connections) communicate at same time
- **Shared:** 1 at a time (CSMA/CD)



Aggregate bandwidth (BW) in switched network is many times shared:

- point-to-point faster since **no arbitration**, simpler interface

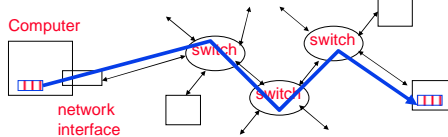


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What makes networks work?

- links connecting switches to each other and to computers or devices



- ability to name the components and to route packets of information - messages - from a source to a destination

- Layering, redundancy, protocols, and encapsulation as means of abstraction (61C big idea)



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Typical Types of Networks

- **Local Area Network (Ethernet)**
 - Inside a building: Up to 1 km
 - (peak) Data Rate: 10 Mbits/sec, 100 Mbits/sec, 1000 Mbits/sec (1.25, 12.5, 125 MBytes/s)
 - Run, installed by network administrators
- **Wide Area Network**
 - Across a continent (10km to 10000 km)
 - (peak) Data Rate: 1.5 Mb/s to 10000 Mb/s
 - Run, installed by telecommunications companies (Sprint, UUNet[MCI], AT&T)

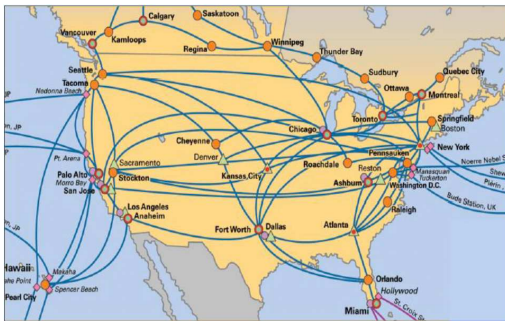


Wireless Networks (LAN), ...

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The Sprint U.S. Topology (2001)



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Administrivia

- **Crunch time**
 - Last Lecture and Course Surveys on Monday 2007-05-07
 - Final Review Session on Wed 2007-05-09
 - 3 weeks + 1 day until the Final ...
 - Final Exam on Saturday 2007-05-12
12:30-3:30 @ 2050 VLSB. IS MANDATORY!
- **Project 4 is out!**
 - May work in pairs.

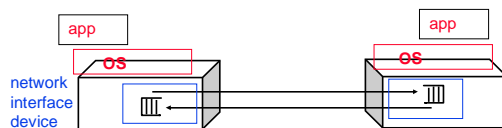


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ABCs of Networks: 2 Computers

- **Starting Point:** Send bits between 2 computers



- Queue (First In First Out) on each end
- Can send both ways ("Full Duplex")
 - One-way information is called "Half Duplex"
- Information sent called a "message"
 - Note: Messages also called packets

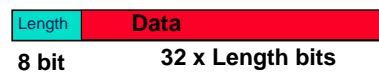


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A Simple Example: 2 Computers

- **What is Message Format?**
 - Similar idea to Instruction Format
 - Fixed size? Number bits?



- **Header (Trailer):** information to deliver message
- **Payload:** data in message
 - What can be in the data?
 - anything that you can represent as bits
 - values, chars, commands, addresses...



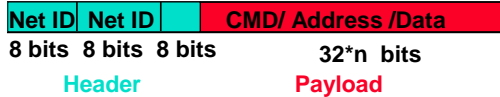
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Questions About Simple Example

- What if more than 2 computers want to communicate?
 - Need computer “address field” in packet to know:
 - which computer should receive it (destination)
 - which computer to reply to (source)
 - Just like envelopes!

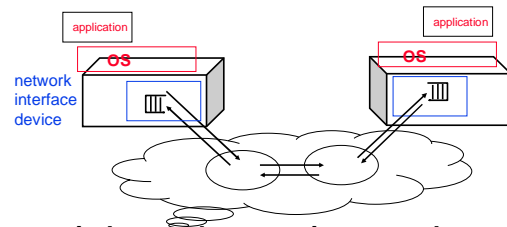
Dest. Source Len



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ABCs: many computers



- switches and routers interpret the header in order to deliver the packet
- source encodes and destination decodes content of the payload



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Questions About Simple Example

- What if message is garbled in transit?
- Add redundant information that is checked when message arrives to be sure it is OK
- 8-bit sum of other bytes: called “Check sum”; upon arrival compare check sum to sum of rest of information in message. XOR also popular.



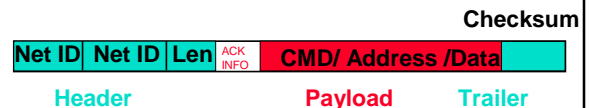
Learn about Checksums in Math 55/CS 70...

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Questions About Simple Example

- What if message never arrives?
 - Receiver tells sender when it arrives
 - Send an ACK (ACKnowledgement) [like registered mail]
 - Sender retries if waits too long
 - Don't discard message until it is ACK'ed
 - If check sum fails, don't send ACK



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Observations About Simple Example

- Simple questions (like those on the previous slides) lead to:
 - more complex procedures to send/receive message
 - more complex message formats
- **Protocol**: algorithm for properly sending and receiving messages (packets)
 - ...an agreement on how to communicate



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Software Protocol to Send and Receive

- SW Send steps
 - 1: Application copies data to OS buffer
 - 2: OS calculates checksum, starts timer
 - 3: OS sends data to network interface HW and says start
- SW Receive steps
 - 3: OS copies data from network interface HW to OS buffer
 - 2: OS calculates checksum, if OK, send ACK; if not, delete message (sender resends when timer expires)
 - 1: If OK, OS copies data to user address space, & signals application to continue



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Protocol for Networks of Networks?

- **Abstraction** to cope with **complexity of communication** (compare to Abstraction for complexity of computation)

• Networks are like onions

• Hierarchy of layers:

- Application (chat client, game, etc.)
- Transport (TCP, UDP)
- Network (IP)
- Physical Link (wired, wireless, etc.)



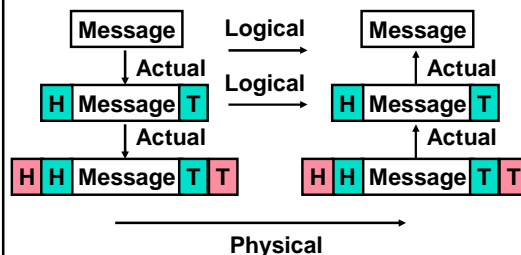
Networks are like onions. They stink? Yes. No! Oh, they make you cry. No!... Layers. Onions have layers. Networks have layers.



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Protocol Family Concept



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Protocol Family Concept

- Key to **protocol families** is that communication occurs **logically** at the same level of the protocol, called **peer-to-peer**...

...but is **implemented via services at the next lower level**

- **Encapsulation**: carry higher level information within lower level "envelope"
- **Fragmentation**: break packet into multiple smaller packets and reassemble



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Protocol for Network of Networks

• IP: **Best-Effort Packet Delivery** (Network Layer)

- Packet switching
 - Send data in packets
 - Header with source & destination address
- "Best effort" delivery
 - Packets may be lost
 - Packets may be corrupted
 - Packets may be delivered out of order



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Protocol for Network of Networks

• **Transmission Control Protocol/Internet Protocol (TCP/IP)** (TCP :: a Transport Layer)

- This protocol family is the **basis of the Internet**, a WAN protocol
- IP makes best effort to deliver
- TCP guarantees delivery
- TCP/IP so popular it is used even when communicating locally: even across homogeneous LAN

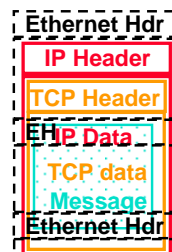


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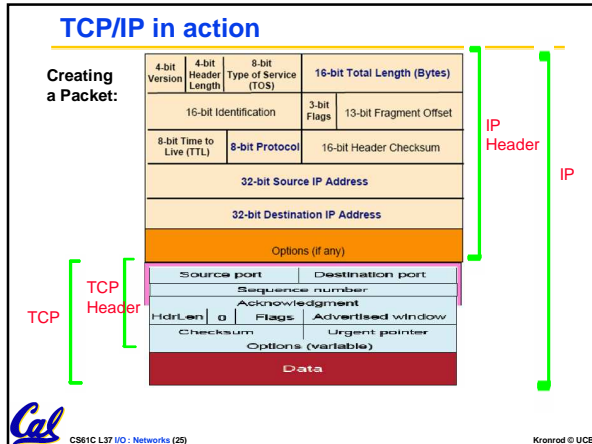
TCP/IP packet, Ethernet packet, protocols

- Application sends message
- TCP breaks into 64KiB segments, adds 20B header
- IP adds 20B header, sends to network
- If Ethernet, broken into 1500B packets with headers, trailers (24B)
- All Headers, trailers have length field, destination,



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Overhead vs. Bandwidth

- Networks are typically advertised using peak bandwidth of network link: e.g., 100 Mbits/sec Ethernet (“100 base T”)
- Software overhead to put message into network or get message out of network often limits useful bandwidth
- Assume overhead to send and receive = 320 microseconds (μ s), want to send 1000 Bytes over “100 Mbit/s” Ethernet
 - Network transmission time: $1000\text{B} \times 8\text{b/B} / 100\text{Mb/s} = 8000\text{b} / (100\text{b}/\mu\text{s}) = 80 \mu\text{s}$
 - Effective bandwidth: $8000\text{b} / (320+80)\mu\text{s} = 20 \text{ Mb/s}$

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And in conclusion...

- Protocol suites allow networking of heterogeneous components
 - Another form of principle of abstraction
 - Protocols \Rightarrow operation in presence of failures
 - Standardization key for LAN, WAN
- Integrated circuit (“Moore’s Law”) revolutionizing network switches as well as processors
 - Switch just a specialized computer
- Trend from shared to switched networks to get faster links and scalable bandwidth
- Interested?
 - EE122 (CS-based in Fall, EE –based in Spring)

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[Bonus] Example: Network Media

Twisted Pair (“Cat 5”): Copper, 1mm thick, twisted to avoid antenna effect

Light: 3 parts are cable, light source, light detector

Fiber Optics

Transmitter Is L.E.D or Laser Diode, light source, Buffer, Cladding, Total internal reflection, Receiver detector, Photodiode, Silica: glass or plastic; actually < 1/10 diameter of copper

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