

inst.eecs.berkeley.edu/~cs61c  
**CS61C : Machine Structures**

**Lecture 37**  
**I/O : Networks**

2007-04-20



**TA Alex Kronrod**

`inst.eecs/~cs61c-te`

**Bad software upgrade causes  
service outage for Blackberry  
users**

**Software upgrade to improve Cache  
performance triggered a problem at  
RIM's hub for North American traffic...  
affecting more than five million users**



# I/O Review

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- I/O gives computers their **5 senses**
- I/O speed range is **12.5-million to one**
- Differences in processor and I/O speed → synchronize with I/O devices before use
- **Polling** works, but expensive
  - processor *repeatedly queries devices*
- **Interrupts** works, more complex
  - device *causes an exception*, causing OS to run and deal with the device
- I/O control leads to **Operating Systems**



# Why Networks?

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- Originally *sharing I/O devices* between **computers**  
ex: printers
- Then *communicating* between **computers**  
ex: file transfer protocol
- Then *communicating* between **people**  
ex: e-mail
- Then *communicating* between **networks of computers**  
ex: file sharing, www, ...



# How Big is the Network (2007)?

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**~30 in 273 Soda**

**~525 in inst.cs.berkeley.edu**

**~6,400 in eecs & cs .berkeley.edu**

**(1999) ~50,000 in berkeley.edu**

**~10,000,000 in .edu** (2005: ~9,000,000)

**~258,941,310 in US** (2005: ~217,000,000, 2006: ~286.5E6)  
(.net .com .edu .arpa .us .mil .org .gov)

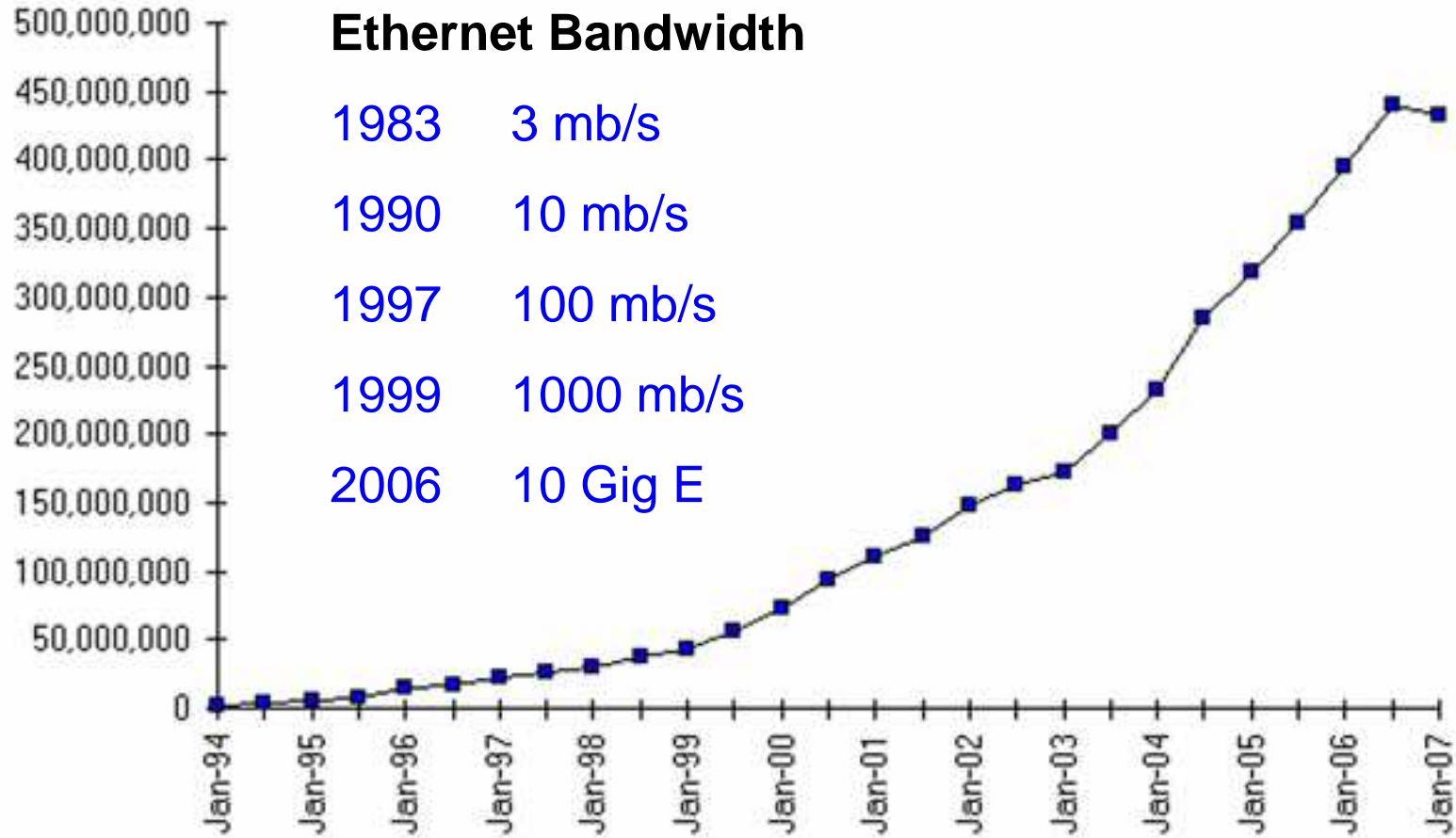
**~433,190,000 in the world**

(2005:~317,000,000, 2006: ~439,000,000)



# Growth Rate

Internet Domain Survey Host Count



Source: Internet Systems Consortium ([www.isc.org](http://www.isc.org))



[en.wikipedia.org/wiki/10\\_gigabit\\_ethernet](http://en.wikipedia.org/wiki/10_gigabit_ethernet)

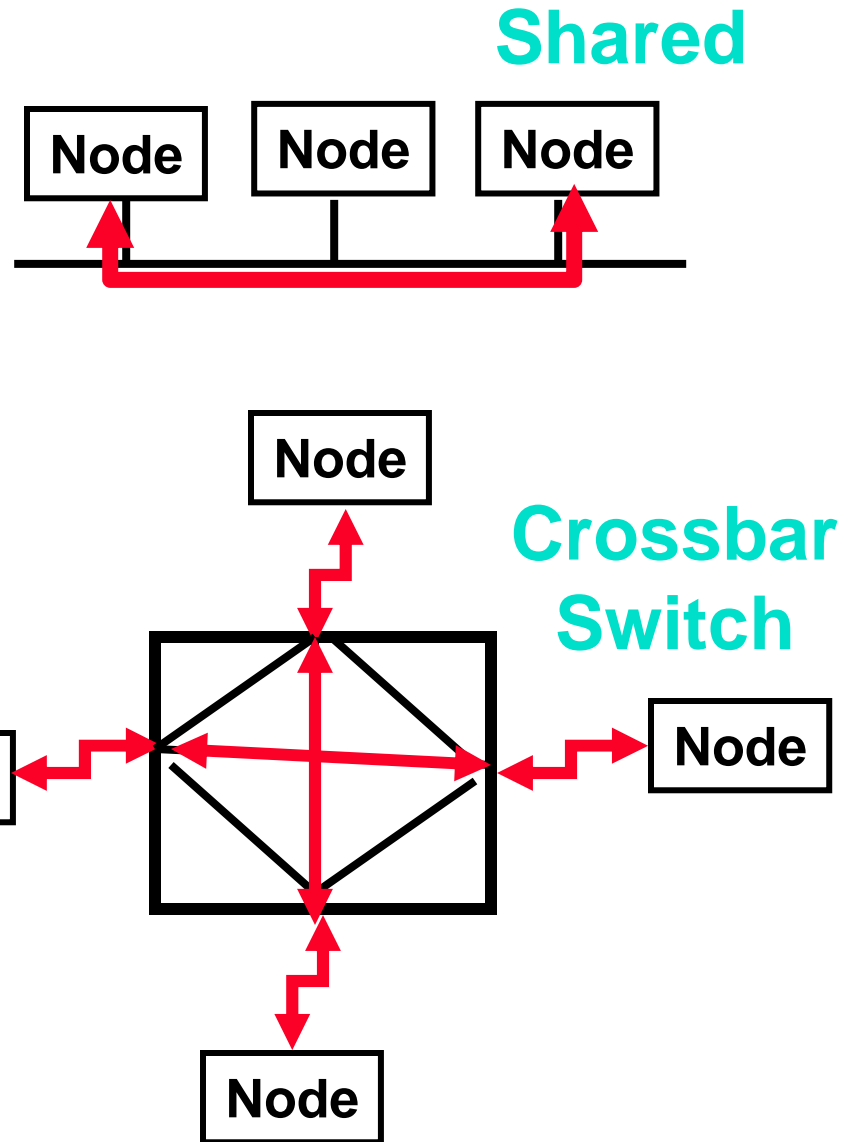
# Shared vs. Switched Based Networks

- **Shared vs. Switched:**

- **Switched:** pairs (“[point-to-point](#)” connections) communicate at same time
- **Shared:** 1 at a time (CSMA/CD)

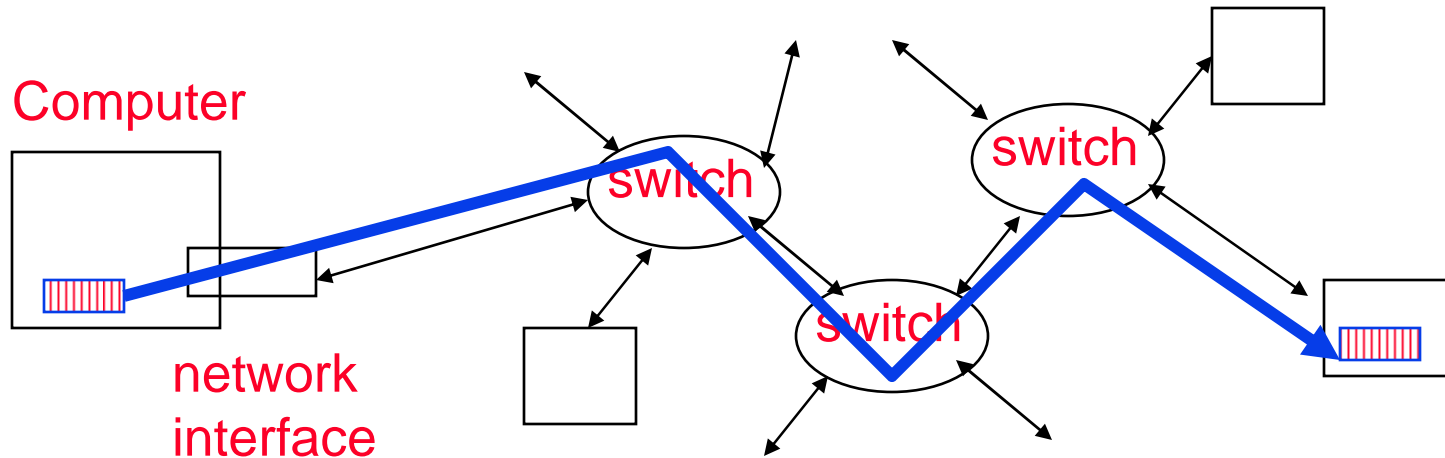
- **Aggregate bandwidth (BW) in switched network is many times shared:**

- point-to-point faster since **no arbitration**, simpler interface



# What makes networks work?

- **links** connecting **switches** to each other and to computers or devices



- ability to **name** the components and to **route** packets of information - messages - from a source to a destination



- Layering, redundancy, protocols, and encapsulation as means of **abstraction** (61C big idea)



# Typical Types of Networks

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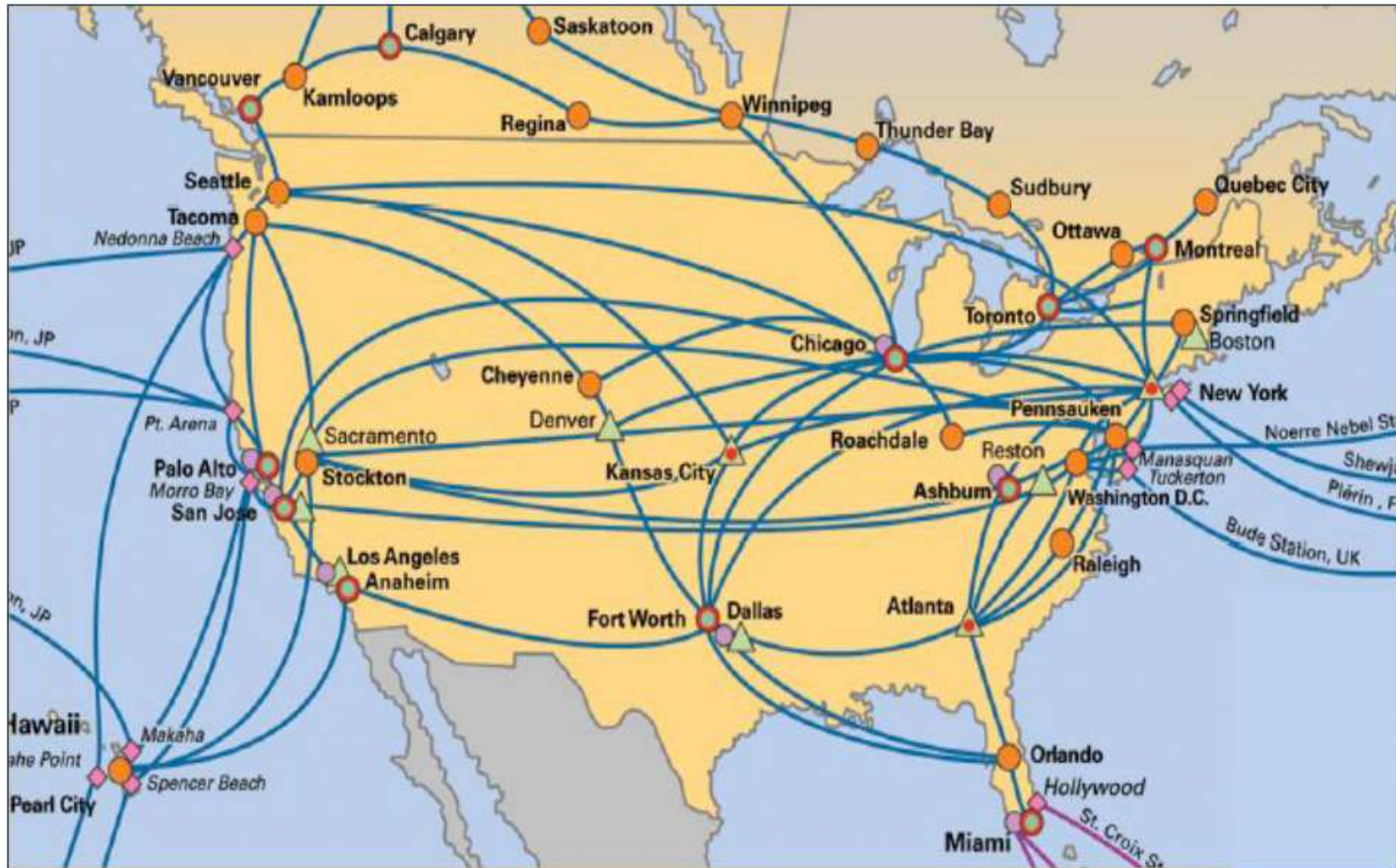
- **Local Area Network (Ethernet)**
  - Inside a building: Up to 1 km
  - (peak) Data Rate: 10 Mbits/sec, 100 Mbits/sec, 1000 Mbits/sec (1.25, 12.5, 125 MBytes/s)
  - Run, installed by network administrators
- **Wide Area Network**
  - Across a continent (10km to 10000 km)
  - (peak) Data Rate: 1.5 Mb/s to 10000 Mb/s
  - Run, installed by telecommunications companies (Sprint, UUNet[MCI], AT&T)



## Wireless Networks (LAN), ...



# The Sprint U.S. Topology (2001)



# Administrivia

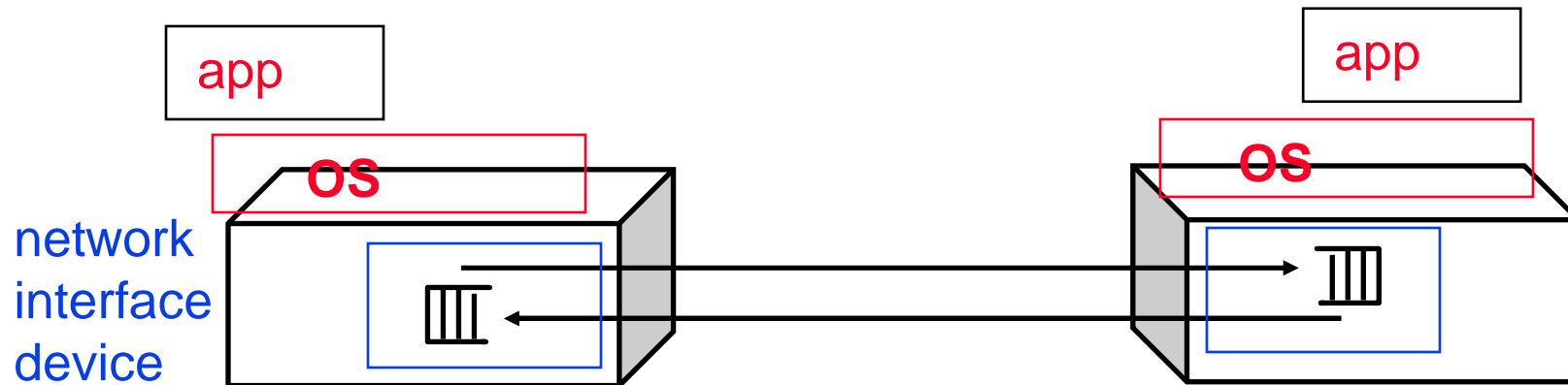
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- **Crunch time**
  - **Last Lecture and Course Surveys on Monday 2007-05-07**
  - **Final Review Session on Wed 2007-05-09**
  - **3 weeks + 1 day until the Final ...**
    - **Final Exam on Saturday 2007-05-12**  
**12:30-3:30 @ 2050 VLSB. IS MANDATORY!**
- **Project 4 is out!**
  - **May work in pairs.**



# ABCs of Networks: 2 Computers

- **Starting Point: Send bits between 2 computers**



- Queue (First In First Out) on each end
- Can send both ways (“**Full Duplex**”)
  - One-way information is called “**Half Duplex**”
- Information sent called a “**message**”
  - Note: Messages also called **packets**



# A Simple Example: 2 Computers

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- **What is Message Format?**
  - **Similar idea to Instruction Format**
  - **Fixed size? Number bits?**



- **Header (Trailer)**: information to deliver message
- **Payload**: data in message
- **What can be in the data?**
  - anything that you can represent as bits
  - values, chars, commands, addresses...



# Questions About Simple Example

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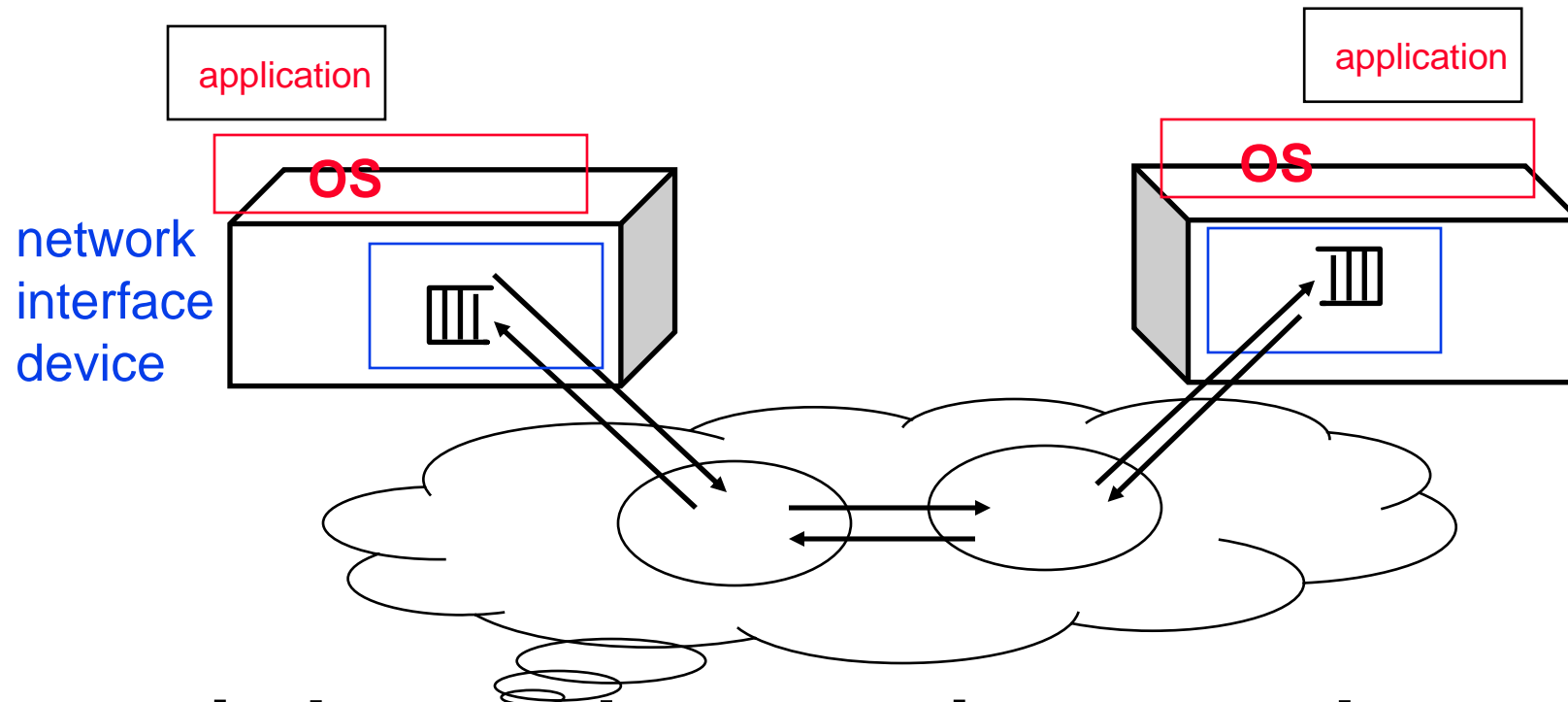
- What if more than 2 computers want to communicate?
  - Need computer “address field” in packet to know:
    - which computer should receive it (destination)
    - which computer to reply to (source)
  - Just like envelopes!

Dest. Source Len



# ABCs: many computers

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- **switches and routers interpret the header in order to deliver the packet**
- **source encodes and destination decodes content of the payload**

# Questions About Simple Example

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- What if message is garbled in transit?
- Add redundant information that is checked when message arrives to be sure it is OK
- 8-bit sum of other bytes: called “**Checksum**”; upon arrival compare checksum to sum of rest of information in message. **xor** also popular.

Checksum



Header

Payload

Trailer



Learn about Checksums in Math 55/CS 70...

# Questions About Simple Example

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- What if message never arrives?
- Receiver tells sender when it arrives
  - Send an ACK (ACKnowledgement) [like registered mail]
  - Sender retries if waits too long
- Don't discard message until it is ACK'ed
- If check sum fails, don't send ACK

Checksum



Header

Payload

Trailer





# Observations About Simple Example

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- **Simple questions (like those on the previous slides) lead to:**
  - more complex procedures to send/receive message
  - more complex message formats
- **Protocol: algorithm for properly sending and receiving messages (packets)**
  - ...an agreement on how to communicate



# Software Protocol to Send and Receive

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- **SW Send steps**

- 1: Application copies data to OS buffer

- 2: OS calculates checksum, starts timer

- 3: OS sends data to network interface HW and says start

- **SW Receive steps**

- 3: OS copies data from network interface HW to OS buffer

- 2: OS calculates checksum, if OK, send ACK; if not, delete message (sender resends when timer expires)

- 1: If OK, OS copies data to user address space, & signals application to continue



# Protocol for Networks of Networks?

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- **Abstraction** to cope with **complexity of communication** (compare to Abstraction for complexity of computation)

- **Networks are like onions**

- **Hierarchy of layers:**

- **Application (chat client, game, etc.)**
- **Transport (TCP, UDP)**
- **Network (IP)**
- **Physical Link (wired, wireless, etc.)**



Networks are like onions.

They stink?

Yes. No!

Oh, they make you cry.

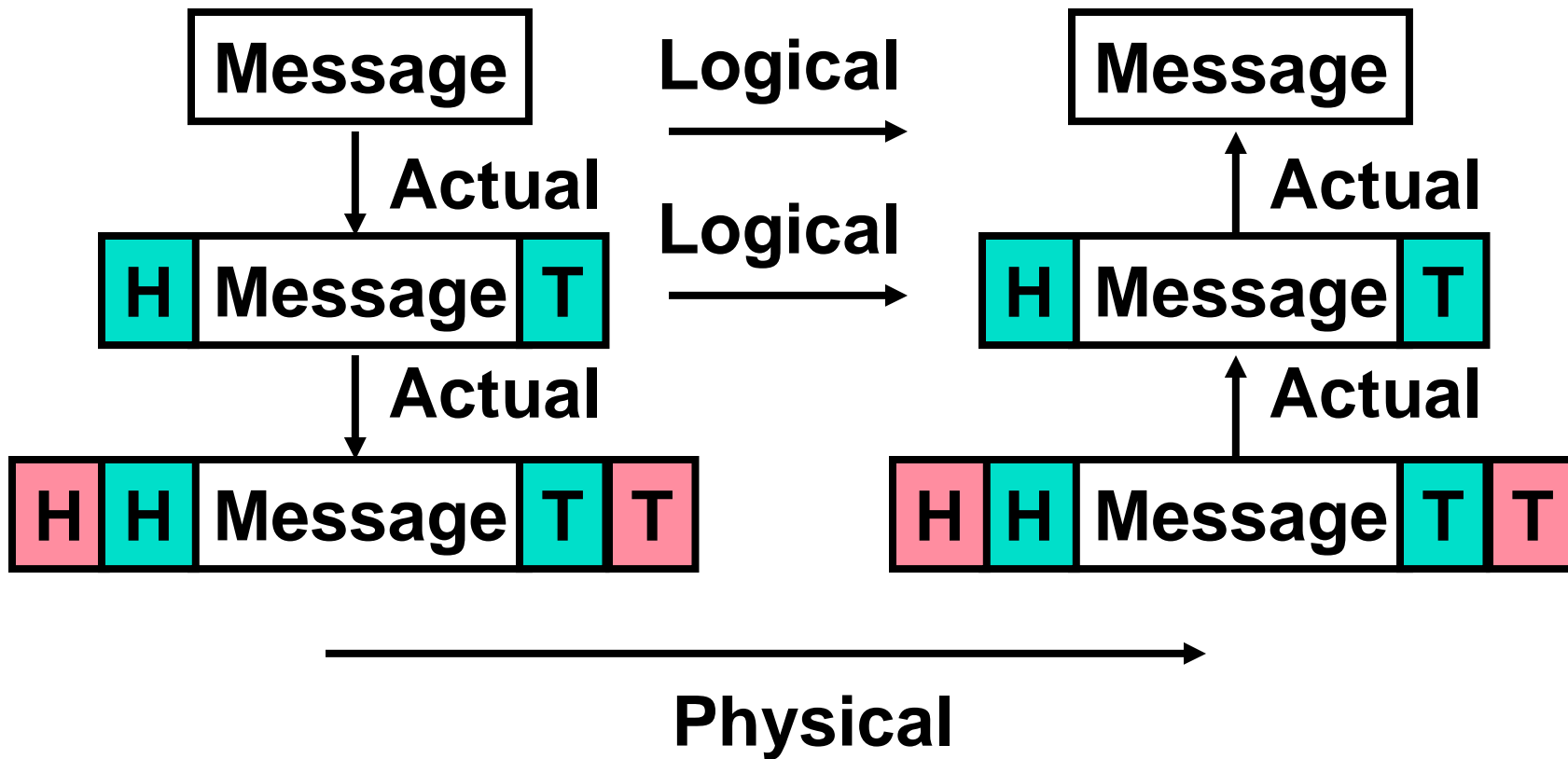
No!... Layers. Onions have layers.

Networks have layers.



# Protocol Family Concept

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# Protocol Family Concept

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- Key to **protocol families** is that communication occurs **logically** at the same level of the protocol, called **peer-to-peer**...  
  
...but is **implemented via services at the next lower level**
- **Encapsulation**: carry higher level information within lower level “envelope”
- **Fragmentation**: break packet into multiple smaller packets and reassemble



# Protocol for Network of Networks

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- **IP: Best-Effort Packet Delivery**  
**(Network Layer)**
- Packet switching
  - Send data in packets
  - Header with source & destination address
- “Best effort” delivery
  - Packets may be lost
  - Packets may be corrupted
  - Packets may be delivered out of order



# Protocol for Network of Networks

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- Transmission Control Protocol/Internet Protocol (TCP/IP)

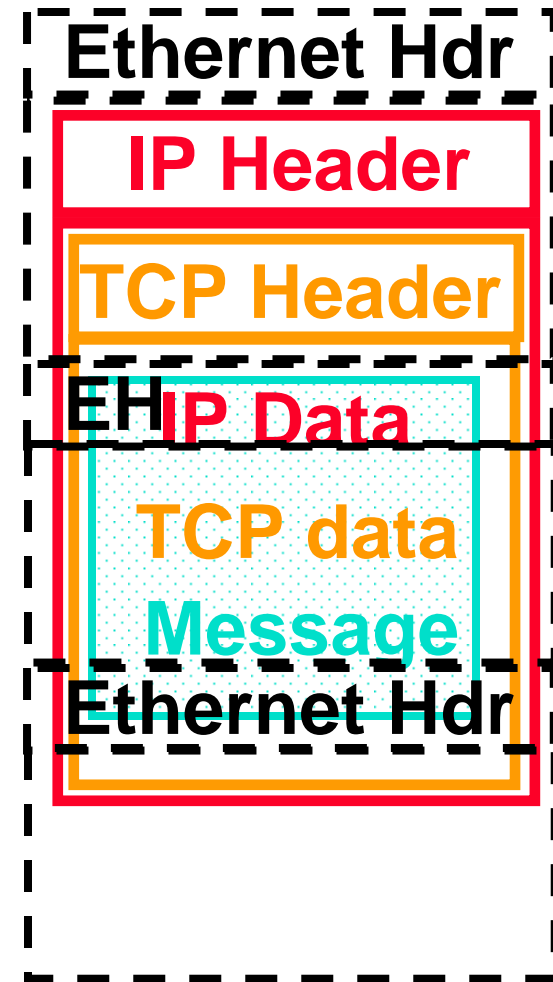
(TCP :: a Transport Layer)

- This protocol family is the **basis of the Internet**, a WAN protocol
- IP makes best effort to deliver
- TCP guarantees delivery
- TCP/IP so popular it is used even when communicating locally: even across homogeneous LAN



# TCP/IP packet, Ethernet packet, protocols

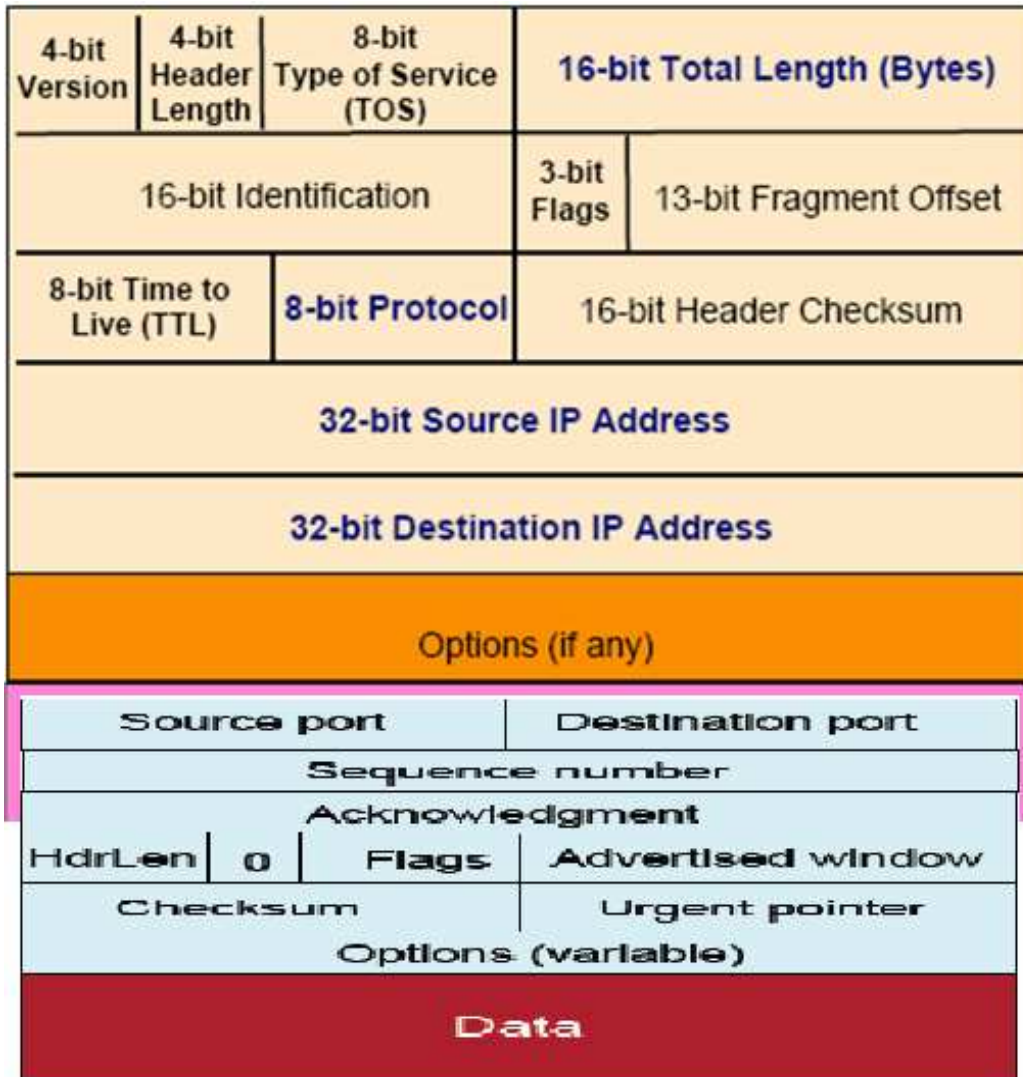
- Application sends message
- TCP breaks into 64KiB segments, adds 20B header
- IP adds 20B header, sends to network
- If Ethernet, broken into 1500B packets with headers, trailers (24B)
- All Headers, trailers have length field, destination,





# TCP/IP in action

Creating a Packet:



TCP  
TCP Header

IP Header  
IP



# Overhead vs. Bandwidth

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- Networks are typically advertised using peak bandwidth of network link: e.g., 100 Mbits/sec Ethernet (“100 base T”)
- Software overhead to put message into network or get message out of network often limits useful bandwidth
- Assume overhead to send and receive = 320 microseconds ( $\mu\text{s}$ ), want to send 1000 Bytes over “100 Mbit/s” Ethernet
  - Network transmission time:  
 $1000\text{B} \times 8\text{b/B} / 100\text{Mb/s}$   
 $= 8000\text{b} / (100\text{b}/\mu\text{s}) = 80 \mu\text{s}$
  - Effective bandwidth:  $8000\text{b} / (320 + 80)\mu\text{s} = 20 \text{ Mb/s}$



## And in conclusion...

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- **Protocol suites allow networking of heterogeneous components**
  - Another form of principle of abstraction
  - Protocols  $\Rightarrow$  operation in presence of failures
  - Standardization key for LAN, WAN
- **Integrated circuit (“Moore’s Law”) revolutionizing network switches as well as processors**
  - Switch just a specialized computer
- **Trend from shared to switched networks to get faster links and scalable bandwidth**
- **Interested?**



- EE122 (CS-based in Fall, EE –based in Spring)

CS61C L37 I/O : Networks (27)

Kronrod © UCB

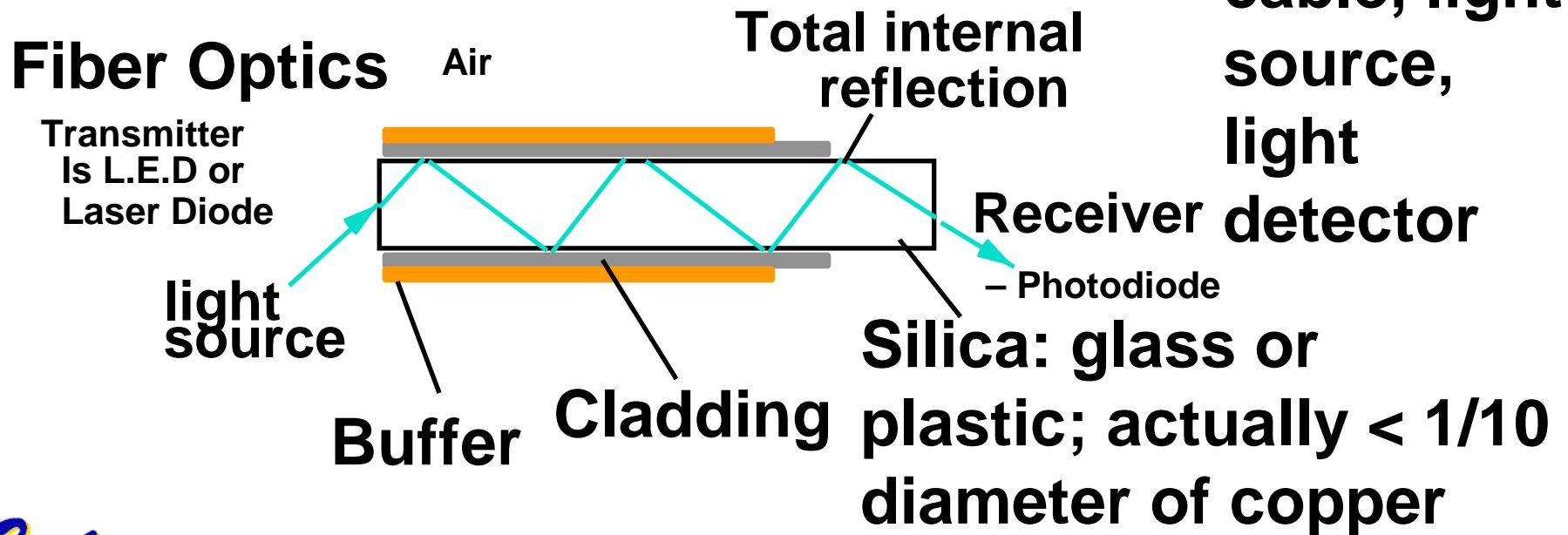
# [Bonus] Example: Network Media

**Twisted Pair  
("Cat 5"):**

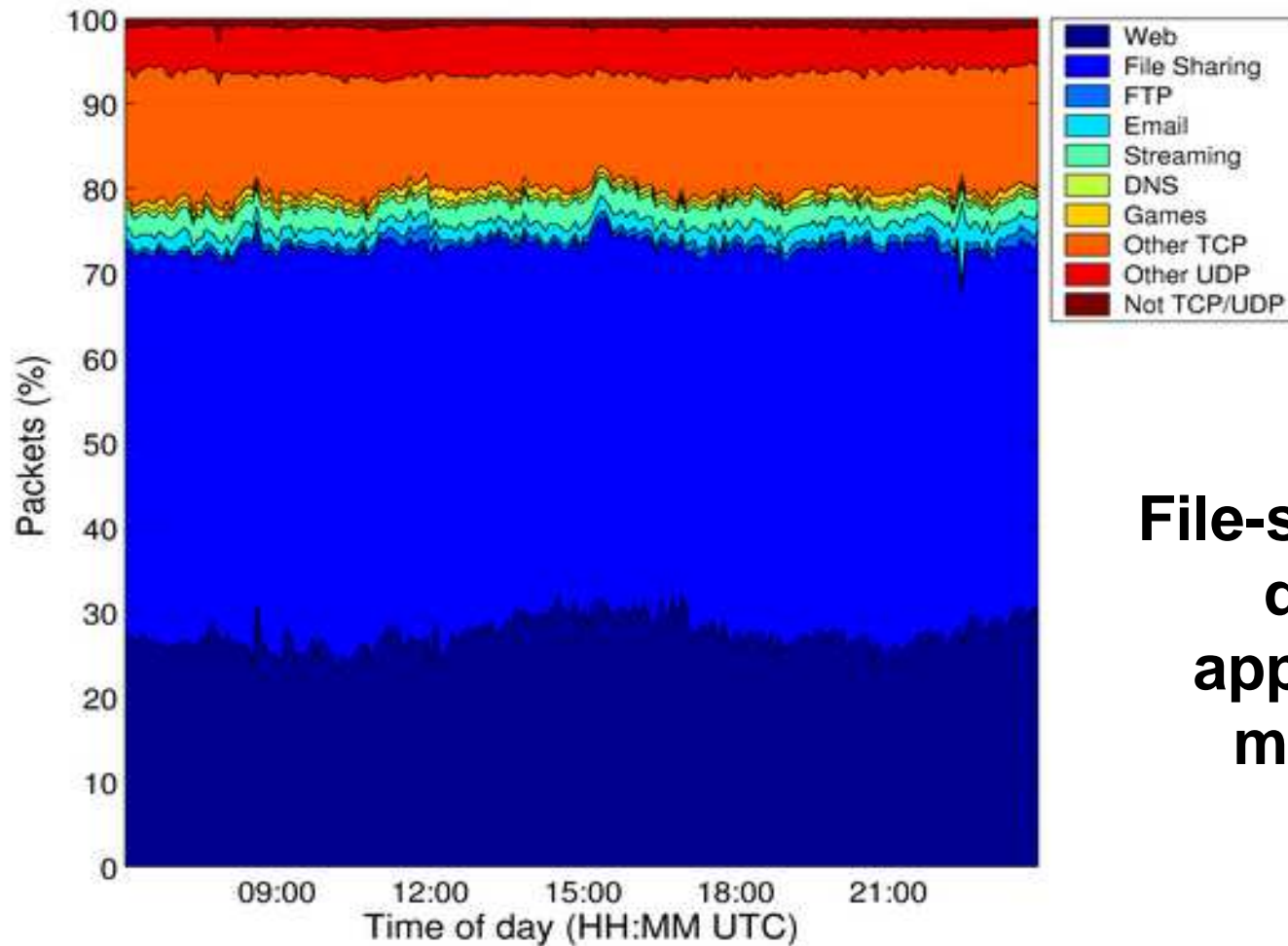


**Copper, 1mm thick, twisted to avoid antenna effect**

**Light:  
3 parts are  
cable, light  
source,  
light**



# [Bonus] Backbone Link App Composition



**File-sharing is the dominant application on many links!**

