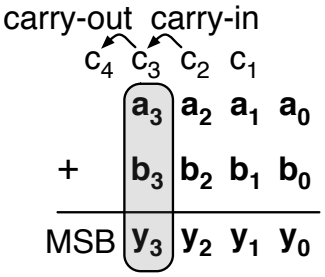
CS 61C Spring 2010 — Signed Arithmetic & MIPS

Scott Beamer (cs61c-ta)

Week 5 (2/16)

Overflow for Signed Numbers

- *Overflow* - if the result of an arithmetic operation is incorrect due to not enough bits in the result
- *Carry-out* - if the carry-out of the MSB is 1
- Overflow is impossible if the sign of the two inputs is different
- Can you think of a way to detect overflow from the carry-in and carry-out of the MSB?



Unsigned Instructions

- As mentioned in lecture, the term 'unsigned' is overloaded in MIPS
- Do/don't sign extend loaded byte - lb / lbu
- Do/don't detect overflow - add, addi, sub, mult, div / addu, addiu, multu, divu
- Do signed / unsigned compare - slt, slti / sltu, sltiu

Decisions

- We can make any comparison (<, >, <=, >=, ==, !=) using slt and beq/bne
- Can reverse the arguments to slt to get sget
- Fill in the gaps in table:

C	MIPS

`int total = 0;`
`int *end = start + n;`
`while(start != end) {`
`  total += start[0];`
`  start++;`
`}`
`return total;`
	`    sll  $a1, $a1, 2`
	`    add  $a1, $a1, $a0`
	`    add  $v0, $0, $0`
	`    beq  $a1, $a0, L2`
	`L1: lw   $t0, 0($a0)`
	`    add  $v0, $v0, $t0`
	`    addi $a0, $a0, 4`
	`    bne  $a1, $a0, L1`
	`L2: jr   $ra`

`switch(op) {`
`  case '+':`
`  result = x + y; break;`
`  case '-':`
`  result = x - y; break;`
`  case '&':`
`  result = x & y; break;`
`  case '|':`
`  result = x | y; break;`
`}`
`return result;`
 	
`// Mappings`
`result→$v0  x→$a0  y→$a1 op→$a2`
	`      addi $t0, $0, 43`
	`      addi $t1, $0, 45`
	`      addi $t2, $0, 38`
	`      addi $t3, $0, 124`
	`      beq  $a2, $t0, plus`
	`      beq  $a2, $t1, sub`
	`      beq  $a2, $t2, and`
	`      beq  $a2, $t3, or`
	`plus: add  $v0, $a0, $a1`
	`      j    done`
	`sub:  sub  $v0, $a0, $a1`
	`      j    done`
	`and:  and  $v0, $a0, $a1`
	`      j    done`
	`or:   or   $v0, $a0, $a1`
	`      j    done`
	`done: jr   $ra`

MIPS Register Conventions

- *Caller* - the calling function
- *Callee* - the function being called
- *Saved* - $s0-s7, $sp
- *Volatile* - $ra, $v0-v1, $a0-a3, $t0-t9, $at
- *Responsibility of Caller* - save any volatile registers it will want after the function call
- *Responsibility of Callee* - save any saved registers it wants to use during function call

Steps to Save Registers

- *Prologue* - move $sp down, store registers into memory offset from $sp
- *Body* - do whatever required saving
- *Epilogue* - load registers from memory offset from $sp, move $sp back up

MIPS Registers (registers in italics aren't needed for CS 61C)

Register #	Register Name	Use
$0	$zero	constant source of 0
$1	$at	used by the assembler
$2-3	$v0-v1	used to return values from a function
$4-7	$a0-a3	used to pass argument into a function
$8-15	$t0-t7	store *temporary* values
$16-23	$s0-s7	store *saved* values
$24-25	$t8-t9	store *temporary* values
$26-27	*$k0-k1*	*used by the kernel*
$28	*$gp*	*global pointer*
$29	$sp	stack pointer
$30	*$fp*	*frame pointer*
$31	$ra	return address

MIPS Register Saving Practice

- Fill in the prologue and epilogue of the caller and callee
- Caller uses: $t0, $s2, $a0, ($ra)
- Callee uses: $s0, $s1, $t2, $v0

Caller	Callee

`# Prologue`
`addi $sp, $sp, -12`
`sw $t0, 8($sp)`
`sw $a0, 4($sp)`
`sw $ra, 0($sp)`
`# Body`
`jal CALLEE`
`# Epilogue`
`lw $t0, 8($sp)`
`lw $a0, 4($sp)`
`lw $ra, 0($sp)`
`addi $sp, $sp, 12`
	`# Prologue`
	`addi $sp, $sp, -8`
	`sw $s0, 4($sp)`
	`sw $s1, 0($sp)`
	
	`# Body`
	`# Epilogue`
	`lw $s0, 4($sp)`
	`lw $s1, 0($sp)`
	`addi $sp, $sp, 8`