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**UCB CS61C : Machine Structures**

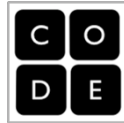
**Lecture 07**  
**Introduction to MIPS : Decisions II**

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**Review**

- Memory is byte-addressable, but **lw** and **sw** access one word at a time.
- A pointer (used by **lw** and **sw**) is just a memory address, so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using conditional statements within **if**, **while**, **do while**, **for**.
- MIPS Decision making instructions are the conditional branches: **beq** and **bne**.
- New Instructions:

**lw, sw, beq, bne, j**



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**Last time: Loading, Storing bytes 1/2**

- In addition to word data transfers (**lw, sw**), MIPS has byte data transfers:
  - load byte: **lb**
  - store byte: **sb**
- same format as **lw, sw**
- E.g., **lb \$s0, 3(\$s1)**
  - contents of memory location with address = sum of "3" + contents of register **s1** is copied to the low byte position of register **s0**.



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**Loading, Storing bytes 2/2**

- What do with other 24 bits in the 32 bit register?
  - lb: sign extends to fill upper 24 bits



- Normally don't want to sign extend chars
- MIPS instruction that doesn't sign extend when loading bytes:
  - load byte unsigned: **lbu**



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**Overflow in Arithmetic (1/2)**

- Reminder: Overflow occurs when there is a mistake in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):
 

15	1111
+ 3	+ 0011
18	10010

  - But we don't have room for 5-bit solution, so the solution would be **0010**, which is **+2**, and wrong.



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**Overflow in Arithmetic (2/2)**

- Some languages detect overflow (Ada), some don't (C)
- MIPS solution is 2 kinds of arithmetic instructs:
  - These cause overflow to be detected
    - add (**add**)
    - add immediate (**addi**)
    - subtract (**sub**)
  - These do not cause overflow detection
    - add unsigned (**addu**)
    - add immediate unsigned (**addiu**)
    - subtract unsigned (**subu**)
- Compiler selects appropriate arithmetic
  - MIPS C compilers produce **addu, addiu, subu**



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## Two "Logic" Instructions

- Here are 2 more new instructions
- Shift Left: `sll $s1,$s2,2 #s1=s2<<2`
  - Store in `$s1` the value from `$s2` shifted 2 bits to the left (they fall off end), inserting 0's on right; `<<` in C.
  - Before: `0000 0002hex`  
`0000 0000 0000 0000 0000 0000 0000 0010two`
  - After: `0000 0008hex`  
`0000 0000 0000 0000 0000 0000 0000 1000two`
  - What arithmetic effect does shift left have?
- Shift Right: `srl` is opposite shift; `>>`

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## Loops in C/Assembly (1/3)

- Simple loop in C; `A[]` is an array of ints
 

```
do { g = g + A[i];
    i = i + j;
  } while (i != h);
```
- Rewrite this as:
 

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```
- Use this mapping:
 

```
g, h, i, j, base of A
$s1, $s2, $s3, $s4, $s5
```

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## Loops in C/Assembly (2/3)

- Final compiled MIPS code:
 

```
Loop: sll $t1,$s3,2 # $t1= 4*I
      addu $t1,$t1,$s5 # $t1=addr A+4i
      lw $t1,0($t1) # $t1=A[i]
      addu $s1,$s1,$t1 # g=g+A[i]
      addu $s3,$s3,$s4 # i=i+j
      bne $s3,$s2,Loop # goto Loop
      # if i!=h
```
- Original code:
 

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```

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## Loops in C/Assembly (3/3)

- There are three types of loops in C:
  - `while`
  - `do ... while`
  - `for`
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch

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## Administrivia

- Any administrivia?

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## Inequalities in MIPS (1/4)

- Until now, we've only tested equalities (`==` and `!=` in C). General programs need to test `<` and `>` as well.
  - Introduce MIPS Inequality Instruction:
    - "Set on Less Than"
    - Syntax: `slt reg1,reg2,reg3`
    - Meaning:
 

```
if (reg2 < reg3)
    reg1 = 1;
else reg1 = 0;
```

Same thing...
- "set" means "change to 1",  
"reset" means "change to 0".

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## Inequalities in MIPS (2/4)

- How do we use this? Compile by hand:  
`if (g < h) goto Less; #g:$s0, h:$s1`
- Answer: compiled MIPS code...  
`slt $t0,$s0,$s1 # $t0 = 1 if g<h  
bne $t0,$0,Less # goto Less  
# if $t0!=0  
# (if (g<h)) Less:`
- Register \$0 always contains the value 0, so `bne` and `beq` often use it for comparison after an `slt` instruction.
- A `slt` → `bne` pair means `if(... < ...)goto...`

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## Inequalities in MIPS (3/4)

- Now we can implement `<`, but how do we implement `>`, `≤` and `≥` ?
- We could add 3 more instructions, but:
  - MIPS goal: Simpler is Better
- Can we implement `≤` in one or more instructions using just `slt` and branches?
  - What about `>`?
  - What about `≥`?

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## Inequalities in MIPS (4/4)

```
# a:$s0, b:$s1  
slt $t0,$s0,$s1 # $t0 = 1 if a<b  
beq $t0,$0,skip # skip if a >= b  
    <stuff>      # do if a<b  
skip:
```

Two independent variations possible:

Use `slt $t0,$s1,$s0` instead of  
`slt $t0,$s0,$s1`  
Use `bne` instead of `beq`

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## Immediates in Inequalities

- There is also an immediate version of `slt` to test against constants: `slti`
  - Helpful in `for` loops

```
C      if (g >= 1) goto Loop
```

```
M      Loop: . . .  
I      slti $t0,$s0,1    # $t0 = 1 if  
P      # $s0<1 (g<1)  
S      beq $t0,$0,Loop  # goto Loop  
                        # if $t0==0  
                        # (if (g>=1))
```

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An `slt` → `beq` pair means `if(... ≥ ...)goto...`

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## What about unsigned numbers?

- Also unsigned inequality instructions:  
`sltu, sltiu`  
... which sets result to **1** or **0** depending on unsigned comparisons
- What is value of `$t0, $t1`?  
(`$s0 = FFFF FFFAhex, $s1 = 0000 FFFAhex`)  
`slt $t0, $s0, $s1`  
`sltu $t1, $s0, $s1`

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## MIPS Signed vs. Unsigned – diff meanings!

- MIPS terms Signed/Unsigned “overloaded”:
  - Do/Don't sign extend
    - (`lb, lbu`)
  - Do/Don't overflow
    - (`add, addi, sub, mult, div`)
    - (`addu, addiu, subu, multu, divu`)
  - Do signed/unsigned compare
    - (`slt, slti/sltu, sltiu`)

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