

`inst.eecs.berkeley.edu/~cs61c`  
**UCB CS61C : Machine Structures**



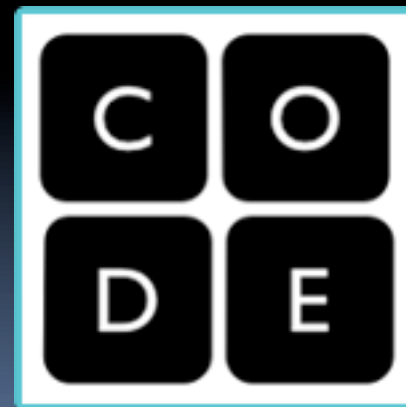
**Sr Lecturer SOE**  
**Dan Garcia**

**Lecture 07**  
**Introduction to MIPS : Decisions II**

**2013-02-06**

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# Review

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- Memory is **byte**-addressable, but **lw** and **sw** access one **word** at a time.
- A pointer (used by **lw** and **sw**) is just a memory address, so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using **conditional statements** within **if, while, do while, for**.
- MIPS Decision making instructions are the **conditional branches: beq and bne**.
- New Instructions:

**lw, sw, beq, bne, j**



# Last time: Loading, Storing bytes 1/2

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- In addition to word data transfers (**lw**, **sw**), MIPS has **byte** data transfers:
  - load byte: **lb**
  - store byte: **sb**
- same format as **lw**, **sw**
- E.g., **lb \$s0, 3(\$s1)**
  - *contents of memory location with address = sum of "3" + contents of register s1 is copied to the low byte position of register s0.*



# Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
  - lb: sign extends to fill upper 24 bits



- Normally don't want to sign extend chars
- MIPS instruction that doesn't sign extend when loading bytes:
  - load byte unsigned: **lbu**



# Overflow in Arithmetic (1/2)

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- Reminder: Overflow occurs when there is a mistake in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):

$$\begin{array}{r} 15 \\ + 3 \\ \hline 18 \end{array}$$

$$\begin{array}{r} 1111 \\ + 0011 \\ \hline 10010 \end{array}$$

- But we don't have room for 5-bit solution, so the solution would be **0010**, which is **+2**, and wrong.



# Overflow in Arithmetic (2/2)

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- Some languages detect overflow (Ada), some don't (C)
- MIPS solution is 2 kinds of arithmetic instructions:
  - These cause overflow to be detected
    - add (**add**)
    - add immediate (**addi**)
    - subtract (**sub**)
  - These do not cause overflow detection
    - add unsigned (**addu**)
    - add immediate unsigned (**addiu**)
    - subtract unsigned (**subu**)
- Compiler selects appropriate arithmetic
  - MIPS C compilers produce **addu, addiu, subu**



# Two “Logic” Instructions

- Here are 2 more new instructions
- Shift Left: `sll $s1,$s2,2` #s1=s2<<2
  - Store in \$s1 the value from \$s2 shifted 2 bits to the left (they fall off end), inserting 0's on right; << in C.
  - Before: `0000 0002hex`  
`0000 0000 0000 0000 0000 0000 0000 0010two`
  - After: `0000 0008hex`  
`0000 0000 0000 0000 0000 0000 0000 1000two`
  - What arithmetic effect does shift left have?
- Shift Right: `srl` is opposite shift; >>



# Loops in C/Assembly (1/3)

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- Simple loop in C; **A[ ]** is an array of ints

```
do { g = g + A[i];  
    i = i + j;  
} while (i != h);
```

- Rewrite this as:

```
Loop: g = g + A[i];  
      i = i + j;  
      if (i != h) goto Loop;
```

- Use this mapping:

<b>g</b> ,	<b>h</b> ,	<b>i</b> ,	<b>j</b> ,	<b>base of A</b>
<b>\$s1</b> ,	<b>\$s2</b> ,	<b>\$s3</b> ,	<b>\$s4</b> ,	<b>\$s5</b>





# Loops in C/Assembly (2/3)

- Final compiled MIPS code:

```
Loop: sll  $t1, $s3, 2      # $t1 = 4*I
      addu $t1, $t1, $s5    # $t1 = addr A+4i
      lw   $t1, 0($t1)     # $t1 = A[i]
      addu $s1, $s1, $t1   # g = g + A[i]
      addu $s3, $s3, $s4   # i = i + j
      bne  $s3, $s2, Loop  # goto Loop
                          # if i != h
```

- Original code:

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```



# Loops in C/Assembly (3/3)

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- There are three types of loops in C:
  - **while**
  - **do ... while**
  - **for**
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch



# Administrivia

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- Any administrivia?



# Inequalities in MIPS (1/4)

- Until now, we've only tested equalities (`==` and `!=` in C). General programs need to test `<` and `>` as well.
- Introduce MIPS Inequality Instruction:
  - "Set on Less Than"
  - Syntax: `slt reg1,reg2,reg3`
  - Meaning:

```
if (reg2 < reg3)
    reg1 = 1;
else reg1 = 0;
```

Same thing...

"set" means "change to 1",  
"reset" means "change to 0".



# Inequalities in MIPS (2/4)

- How do we use this? Compile by hand:

```
if (g < h) goto Less; #g:$s0, h:$s1
```

- Answer: compiled MIPS code...

```
slt $t0,$s0,$s1 # $t0 = 1 if g<h
bne $t0,$0,Less # goto Less
                # if $t0!=0
                # (if (g<h)) Less:
```

- Register \$0 always contains the value 0, so **bne** and **beq** often use it for comparison after an **slt** instruction.
- A **slt** → **bne** pair means **if (... < ...) goto...**



# Inequalities in MIPS (3/4)

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- Now we can implement  $<$ ,  
but how do we implement  $>$ ,  $\leq$  and  $\geq$  ?
- We could add 3 more instructions, but:
  - MIPS goal: **Simpler is Better**
- Can we implement  $\leq$  in one or more instructions using just **slt** and **branches**?
  - What about  $>$ ?
  - What about  $\geq$ ?



# Inequalities in MIPS (4/4)

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```
# a:$s0, b:$s1
slt $t0,$s0,$s1 # $t0 = 1 if a<b
beq $t0,$0,skip # skip if a >= b
    <stuff>     # do if a<b
skip:
```

Two independent variations possible:

Use `slt $t0,$s1,$s0` instead of

`slt $t0,$s0,$s1`

Use `bne` instead of `beq`



# Immediates in Inequalities

- There is also an immediate version of **slt** to test against constants: **slti**
  - Helpful in **for** loops

**C**      **if (g >= 1) goto Loop**

**MIPS**      **Loop:**      . . .

**I**      **slti \$t0,\$s0,1**      # \$t0 = 1 if

**P**                     # \$s0 < 1 (g < 1)

**S**      **beq \$t0,\$0,Loop**      # goto Loop

                    # if \$t0 == 0

                    # (if (g >= 1))



An **slt** → **beq** pair means **if (... ≥ ...) goto...**



# What about unsigned numbers?

- Also **unsigned** inequality instructions:

**sltu, sltiu**

...which sets result to **1** or **0** depending on unsigned comparisons

- What is value of **\$t0**, **\$t1**?

**(\$s0 = FFFF FFFA<sub>hex</sub>, \$s1 = 0000 FFFA<sub>hex</sub>)**

**slt \$t0, \$s0, \$s1**

**sltu \$t1, \$s0, \$s1**



# MIPS Signed vs. Unsigned – diff meanings!

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- MIPS terms Signed/Unsigned “overloaded”:
  - Do/Don't sign extend
    - (lb, lbu)
  - Do/Don't overflow
    - (add, addi, sub, mult, div)
    - (addu, addiu, subu, multu, divu)
  - Do signed/unsigned compare
    - (slt, slti/sltu, sltiu)



# Peer Instruction

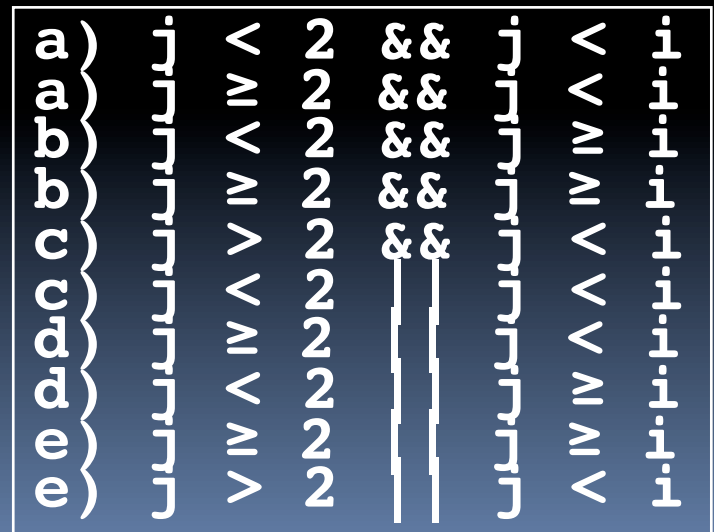
```

Loop: addi $s0, $s0, -1      # i = i - 1
      slti $t0, $s1, 2      # $t0 = (j < 2)
      beq  $t0, $0, Loop    # goto Loop if $t0 == 0
      slt  $t0, $s1, $s0    # $t0 = (j < i)
      bne  $t0, $0, Loop    # goto Loop if $t0 != 0
  
```

( $\$s0=i$ ,  $\$s1=j$ )

What C code properly fills in the blank in loop below?

do { $i--$ ; } while(\_\_\_\_);



# “And in conclusion...”

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- To help the **conditional branches** make decisions concerning inequalities, we introduce: “Set on Less Than” called

**slt, slti, sltu, sltiu**

- One can store and load (signed and unsigned) **bytes** as well as words with **lb, lbu**

- Unsigned add/sub **don't cause overflow**

- New MIPS Instructions:

**sll, srl, lb, lbu**

**slt, slti, sltu, sltiu**

**addu, addiu, subu**



# Bonus Slides



# Example: The C Switch Statement (1/3)

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- Choose among four alternatives depending on whether `k` has the value 0, 1, 2 or 3. Compile this C code:

```
switch (k) {  
    case 0: f=i+j; break; /* k=0 */  
    case 1: f=g+h; break; /* k=1 */  
    case 2: f=g-h; break; /* k=2 */  
    case 3: f=i-j; break; /* k=3 */  
}
```



# Example: The C Switch Statement (2/3)

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- This is complicated, so **simplify**.
- Rewrite it as a chain of if-else statements, which we already know how to compile:

```
if(k==0) f=i+j;
    else if(k==1) f=g+h;
        else if(k==2) f=g-h;
            else if(k==3) f=i-j;
```

- Use this mapping:

```
f:$s0, g:$s1, h:$s2,
i:$s3, j:$s4, k:$s5
```



# Example: The C Switch Statement (3/3)

- Final compiled MIPS code:

```
    bne $s5,$0,L1      # branch k!=0
    add $s0,$s3,$s4    #k==0 so f=i+j
    j   Exit          # end of case so Exit
L1:  addi $t0,$s5,-1   # $t0=k-1
    bne $t0,$0,L2     # branch k!=1
    add $s0,$s1,$s2    #k==1 so f=g+h
    j   Exit          # end of case so Exit
L2:  addi $t0,$s5,-2   # $t0=k-2
    bne $t0,$0,L3     # branch k!=2
    sub $s0,$s1,$s2    #k==2 so f=g-h
    j   Exit          # end of case so Exit
L3:  addi $t0,$s5,-3   # $t0=k-3
    bne $t0,$0,Exit    # branch k!=3
    sub $s0,$s3,$s4    # k==3 so f=i-j
```

Exit:

