

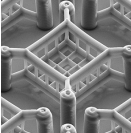
Lecture 26 – Combinational Logic Blocks



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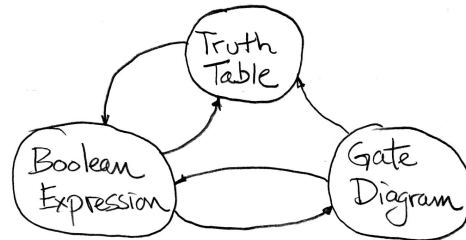
Very fast 3D Micro Printer =>

A new company called Nanoscribe has developed a fabrication device that can create structures like the one at the right at the micro scale in minutes (instead of hours). The idea is that "tiny, ultrashort pulses from a near-infrared laser on a light-sensitive material solidifies on spot. Mirrors not motors



Review

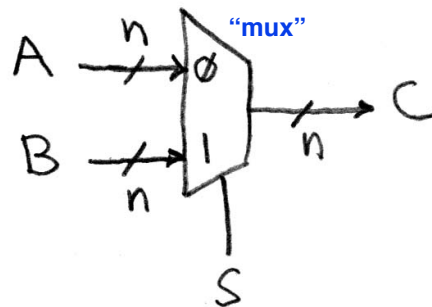
- Use this table and techniques we learned to transform from 1 to another



Today

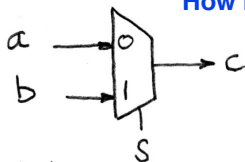
- Data Multiplexors
- Arithmetic and Logic Unit
- Adder/Subtractor

Data Multiplexor (here 2-to-1, n-bit-wide)



N instances of 1-bit-wide mux

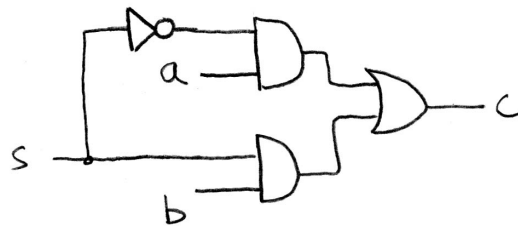
How many rows in TT?



$$\begin{aligned}
 c &= \bar{s}a\bar{b} + \bar{s}ab + s\bar{a}b + sab \\
 &= \bar{s}(a\bar{b} + ab) + s(\bar{a}b + ab) \\
 &= \bar{s}(a(\bar{b} + b)) + s((\bar{a} + a)b) \\
 &= \bar{s}(a(1)) + s((1)b) \\
 &= \bar{s}a + sb
 \end{aligned}$$

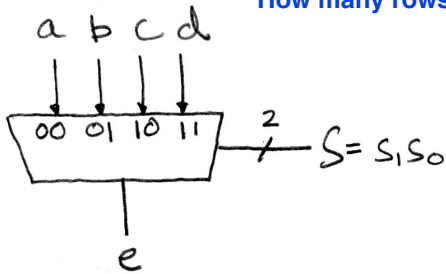
How do we build a 1-bit-wide mux?

$$\bar{s}a + sb$$



4-to-1 Multiplexor?

How many rows in TT?

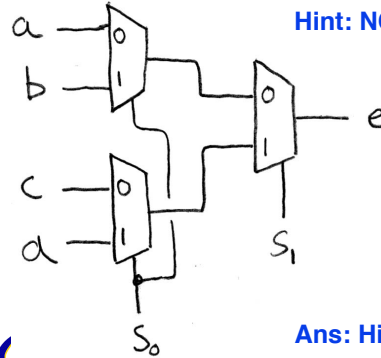


$$e = \bar{s}_1\bar{s}_0a + \bar{s}_1s_0b + s_1\bar{s}_0c + s_1s_0d$$



Is there any other way to do it?

Hint: NCAA tourney!

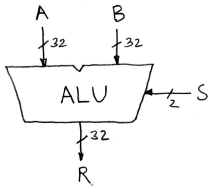


Ans: Hierarchically!



Arithmetic and Logic Unit

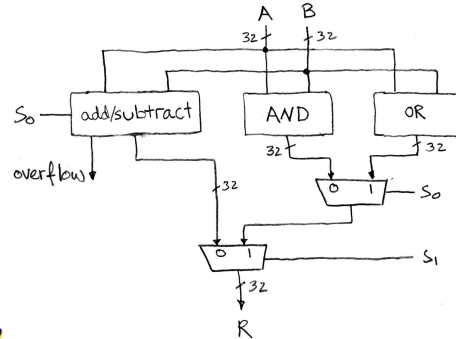
- Most processors contain a special logic block called "Arithmetic and Logic Unit" (ALU)
- We'll show you an easy one that does ADD, SUB, bitwise AND, bitwise OR



when $S=00$, $R=A+B$
 when $S=01$, $R=A-B$
 when $S=10$, $R=A \text{ AND } B$
 when $S=11$, $R=A \text{ OR } B$



Our simple ALU



Adder/Subtractor Design -- how?

- Truth-table, then determine canonical form, then minimize and implement as we've seen before
- Look at breaking the problem down into smaller pieces that we can cascade or hierarchically layer



Adder/Subtractor – One-bit adder LSB...

	a_3	a_2	a_1	a_0	a_0	b_0	s_0	c_1
					0	0	0	0
+					0	1	1	0
					1	0	1	0
					1	1	0	1
	s_3	s_2	s_1	s_0				

$$s_0 =$$

$$c_1 =$$



Adder/Subtractor – One-bit adder (1/2)...

	a_i	b_i	c_i	s_i	c_{i+1}
	0	0	0	0	0
	0	0	1	1	0
	0	1	0	1	0
	0	1	1	0	1
	1	0	0	1	0
	1	0	1	0	1
	1	1	0	0	1
	1	1	1	1	1

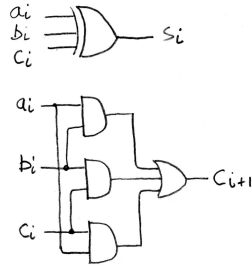
a_3	a_2	a_1	a_0
b_3	b_2	b_1	b_0
s_3	s_2	s_1	s_0

$$s_i =$$

$$c_{i+1} =$$



Adder/Subtractor – One-bit adder (2/2)...

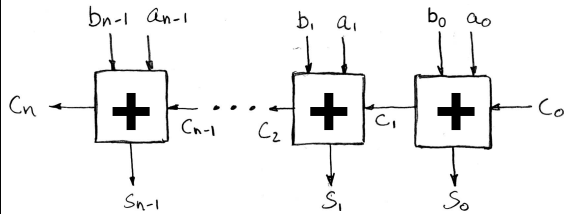


$$s_i = \text{XOR}(a_i, b_i, c_i)$$

$$c_{i+1} = \text{MAJ}(a_i, b_i, c_i) = a_i b_i + a_i c_i + b_i c_i$$



N 1-bit adders \Rightarrow 1 N-bit adder



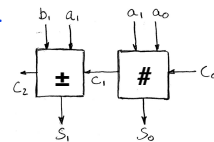
What about overflow?
Overflow = c_n ?



What about overflow?

Consider a 2-bit signed # & overflow:

- 10 = -2 + -2 or -1
- 11 = -1 + -2 only
- 00 = 0 NOTHING!
- 01 = 1 + 1 only



Highest adder

- $C_1 = \text{Carry-in} = C_{in}$, $C_2 = \text{Carry-out} = C_{out}$
- No C_{out} or $C_{in} \Rightarrow$ NO overflow!

What • C_{in} , and $C_{out} \Rightarrow$ NO overflow!

op?

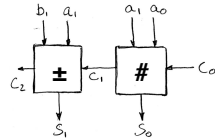
- C_{in} , but no $C_{out} \Rightarrow$ A,B both > 0, overflow!
- C_{out} , but no $C_{in} \Rightarrow$ A,B both < 0, overflow!



What about overflow?

Consider a 2-bit signed # & overflow:

- 10 = -2
- 11 = -1
- 00 = 0
- 01 = 1



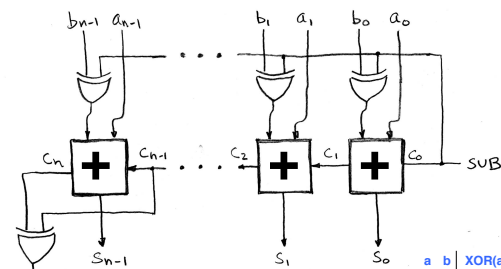
Overflows when...

- C_{in} , but no $C_{out} \Rightarrow$ A,B both > 0, overflow!
- C_{out} , but no $C_{in} \Rightarrow$ A,B both < 0, overflow!

$$\text{overflow} = c_n \text{ XOR } c_{n-1}$$



Extremely Clever Subtractor



XOR serves as
conditional inverter!

a	b	XOR(a,b)
0	0	0
0	1	1
1	0	1
1	1	0



Peer Instruction

- 1) Truth table for mux with 4-bits of signals has 2^4 rows
- 2) We could cascade N 1-bit shifters to make 1 N-bit shifter for sll, srl

	12
a)	FF
b)	FT
c)	TF
d)	TT



“And In conclusion...”

- Use muxes to select among input
 - S input bits selects 2^S inputs
 - Each input can be n-bits wide, indep of S
- Can implement muxes hierarchically
- ALU can be implemented using a mux
 - Coupled with basic block elements
- N-bit adder-subtractor done using N 1-bit adders with XOR gates on input
 - XOR serves as conditional inverter

