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UCB CS61C : Machine Structures




Lecture 07
Introduction to MIPS : Decisions II

Guest Lecturer
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“SO MANY GADGETS, SO MANY ACHES” NYT


Laptops “do not meet any of the ergonomic requirements for a computer system”. Touch screens “should not be used heavily for typing” Texting is a problem because thumb bones have two bones instead of three ... “if you want to get injured, do a lot of texting”. Advice? Take a break



www.nytimes.com/2010/02/19/technology/19china.html

Review


- Memory is byte-addressable, but **lw** and **sw** access one word at a time.
- A pointer (used by **lw** and **sw**) is just a memory address, so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using conditional statements within **if**, **while**, **do while**, **for**.
- MIPS Decision making instructions are the conditional branches: **beq** and **bne**.
- New Instructions:
lw, sw, beq, bne, j



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Last time: Loading, Storing bytes 1/2

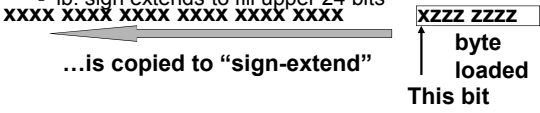
- In addition to word data transfers (**lw, sw**), MIPS has byte data transfers:
 - load byte: **lb**
 - store byte: **sb**
- same format as **lw, sw**
- E.g., **lb \$s0, 3(\$s1)**
 - contents of memory location with address = sum of “3” + contents of register **s1** is copied to the low byte position of register **s0**.




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Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
 - lb: sign extends to fill upper 24 bits



- Normally don't want to sign extend chars
- MIPS instruction that doesn't sign extend when loading bytes:
 - load byte unsigned: **lbu**




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Overflow in Arithmetic (1/2)

- Reminder: Overflow occurs when there is a “mistake” in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):

| | |
|-----|--------|
| 15 | 1111 |
| + 3 | + 0011 |
| 18 | 10010 |


 - But we don't have room for 5-bit solution, so the solution would be **0010**, which is **+2**, and “wrong”.



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Overflow in Arithmetic (2/2)

- Some languages detect overflow (Ada), some don't (most C implementations)
- MIPS solution is 2 kinds of arithmetic instructions:
 - These cause overflow to be detected
 - add (**add**)
 - add immediate (**addi**)
 - subtract (**sub**)
 - These do not cause overflow detection
 - add unsigned (**addu**)
 - add immediate unsigned (**addiu**)
 - subtract unsigned (**subu**)
- Compiler selects appropriate arithmetic
 - MIPS C compilers produce **addu, addiu, subu**



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Two "Logic" Instructions

- Here are 2 more new instructions
- Shift Left: **sll \$s1,\$s2,2 #s1=s2<<2**
 - Store in \$s1 the value from \$s2 shifted 2 bits to the left (they fall off end), inserting 0's on right; << in C.
 - Before: `0000 0002hex`
`0000 0000 0000 0000 0000 0000 0000 0010two`
 - After: `0000 0008hex`
`0000 0000 0000 0000 0000 0000 0000 1000two`
 - What arithmetic effect does shift left have?
- Shift Right: **srl** is opposite shift; >>

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Loops in C/Assembly (1/3)

- Simple loop in C; **A[]** is an array of ints

```
do { g = g + A[i];
    i = i + j;
} while (i != h);
```
- Rewrite this as:

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```
- Use this mapping:
`g, h, i, j, &A[0]`
`$s1, $s2, $s3, $s4, $s5`

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Loops in C/Assembly (2/3)

- Final compiled MIPS code:

```
Loop: sll $t1,$s3,2 # $t1=4*i
      addu $t1,$t1,$s5 # $t1=addr A+4i
      lw $t1,0($t1) # $t1=A[i]
      addu $s1,$s1,$t1 # g=g+A[i]
      addu $s3,$s3,$s4 # i=i+j
      bne $s3,$s2,Loop # goto Loop
      # if i!=h
```
- Original code:

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```

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Loops in C/Assembly (3/3)

- There are three types of loops in C:
 - while**
 - do ... while**
 - for**
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch

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Administrivia

- HW2 is due Sunday at 23:59:59

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Inequalities in MIPS (1/4)

- Until now, we've only tested equalities (`==` and `!=` in C). General programs need to test `<` and `>` as well.
 - Introduce MIPS Inequality Instruction:
 - "Set on Less Than"
 - Syntax: **slt reg1,reg2,reg3**
 - Meaning: `reg1 = (reg2 < reg3);`
- ```
if (reg2 < reg3)
 reg1 = 1;
else reg1 = 0;
```
- ← Same thing...
- 
- "set" means "change to 1",
- 
- "reset" means "change to 0".

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## Inequalities in MIPS (2/4)

- How do we use this? Compile by hand:  
**if (g < h) goto Less; #g:\$s0, h:\$s1**
- Answer: compiled MIPS code...  
**slt \$t0,\$s0,\$s1 # \$t0 = 1 if g<h**  
**bne \$t0,\$0,Less # goto Less**  
**# if \$t0!=0**  
**# (if (g<h)) Less:**
- Register **\$0** always contains the value 0, so **bne** and **beq** often use it for comparison after an **slt** instruction.
- A **slt** → **bne** pair means **if(... < ...)goto...**



## Inequalities in MIPS (3/4)

- Now we can implement **<**,  
but how do we implement **>**, **≤** and **≥** ?
- We could add 3 more instructions, but:
  - MIPS goal: Simpler is Better
- Can we implement **≤** in one or more instructions using just **slt** and branches?
  - What about **>**?
  - What about **≥**?



## Inequalities in MIPS (4/4)

```
a:$s0, b:$s1
slt $t0,$s0,$s1 # $t0 = 1 if a<b
beq $t0,$0,skip # skip if a >= b
<stuff> # do if a<b
skip:
```

Two independent variations possible:

Use **slt \$t0,\$s1,\$s0** instead of  
**slt \$t0,\$s0,\$s1**

Use **bne** instead of **beq**



## Immediates in Inequalities

- There is also an immediate version of **slt** to test against constants: **slti**
  - Helpful in **for** loops

**C** if (g >= 1) goto Loop

```
Loop: ...
M
I slti $t0,$s0,1 # $t0 = 1 if
P # $s0<1 (g<1)
S beq $t0,$0,Loop # goto Loop
if $t0==0
(if (g>=1))
```

An **slt** → **beq** pair means **if(... ≥ ...)goto...**



## What about unsigned numbers?

- Also unsigned inequality instructions:  
**sltu, sltiu**  
...which sets result to **1** or **0** depending on unsigned comparisons
- What is value of **\$t0**, **\$t1**?  
(**\$s0 = FFFF FFFA<sub>hex</sub>**, **\$s1 = 0000 FFFA<sub>hex</sub>**)  
**slt \$t0, \$s0, \$s1**  
**sltu \$t1, \$s0, \$s1**



## MIPS Signed vs. Unsigned – diff meanings!

- MIPS terms Signed/Unsigned “overloaded”:
  - Do/Don't sign extend
    - (**lb, lbu**)
  - Do/Don't overflow
    - (**add, addi, sub, mult, div**)
    - (**addu, addiu, subu, multu, divu**)
  - Do signed/unsigned compare
    - (**slt, slti/sltu, sltiu**)



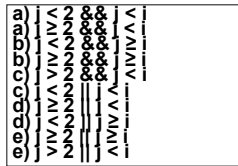
## Peer Instruction

```
Loop: addi $s0,$s0,-1 # i = i - 1
 slti $t0,$s1,2 # $t0 = (j < 2)
 beq $t0,$0,Loop # goto Loop if $t0 == 0
 slt $t0,$s1,$s0 # $t0 = (j < i)
 bne $t0,$0,Loop # goto Loop if $t0 != 0
```

(\$s0=i, \$s1=j)

What C code properly fills in the blank in loop below?

do {i--;} while(\_\_\_\_);



## “And in conclusion...”

- To help the conditional branches make decisions concerning inequalities, we introduce: “Set on Less Than” called **slt, slti, sltu, sltiu**
- One can store and load (signed and unsigned) bytes as well as words with **lb, lbu**
- Unsigned add/sub don’t cause overflow
- New MIPS Instructions:
  - sll, srl, lb, lbu**
  - slt, slti, sltu, sltiu**
  - addu, addiu, subu**



## Bonus Slides



## Example: The C Switch Statement (1/3)

- Choose among four alternatives depending on whether k has the value 0, 1, 2 or 3. Compile this C code:

```
switch (k) {
 case 0: f=i+j; break; /* k=0 */
 case 1: f=g+h; break; /* k=1 */
 case 2: f=g-h; break; /* k=2 */
 case 3: f=i-j; break; /* k=3 */
}
```



## Example: The C Switch Statement (2/3)

- This is complicated, so simplify.
- Rewrite it as a chain of if-else statements, which we already know how to compile:
 

```
if(k==0) f=i+j;
 else if(k==1) f=g+h;
 else if(k==2) f=g-h;
 else if(k==3) f=i-j;
```
- Use this mapping:
 

```
f:$s0, g:$s1, h:$s2,
i:$s3, j:$s4, k:$s5
```



## Example: The C Switch Statement (3/3)

- Final compiled MIPS code:
 

```
bne $s5,$0,L1 # branch k!=0
add $s0,$s3,$s4 #k==0 so f=i+j
j Exit # end of case so Exit
L1: addi $t0,$s5,-1 # $t0=k-1
bne $t0,$0,L2 # branch k!=1
add $s0,$s1,$s2 #k==1 so f=g+h
j Exit # end of case so Exit
L2: addi $t0,$s5,-2 # $t0=k-2
bne $t0,$0,L3 # branch k!=2
sub $s0,$s1,$s2 #k==2 so f=g-h
j Exit # end of case so Exit
L3: addi $t0,$s5,-3 # $t0=k-3
bne $t0,$0,Exit # branch k!=3
sub $s0,$s3,$s4 # k==3 so f=i-j
Exit:
```

