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## Lecture 07 Introduction to MIPS : Decisions II

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## "SO MANY GADGETS, SO MANY ACHES" NYT

Laptops "do not meet any of the ergonomic requirements for a computer system". Touch screens "should not be used heavily for typing" Texting is a problem because thumb bones have two bones instead of three ... "if you want to get injured, do a lot of texting". Advice? Take a break


## Review

- Memory is byte-addressable, but Iw and sw access one word at a time.
- A pointer (used by Iw and sw) is just a memory address, so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using conditional statements within if, while, do while, for.
- MIPS Decision making instructions are the conditional branches: beq and bne.
- New Instructions:

Iw, sw, beq, bne, j

## Last time: Loading, Storing bytes 1/2

- In addition to word data transfers
(lw, sw), MIPS has byte data transfers:
- load byte: Ib
- store byte: sb
- same format as Iw, sw
- E.g., lb \$s0, 3(\$s1)
- contents of memory location with address = sum of " 3 " + contents of register s1 is copied to the low byte position of register s0.


## Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
- lb: sign extends to fill upper 24 bits XXXX XxXX XxXX XxXX Xxxx Xxxx
...is copied to "sign-extend"
- Normally don't want to sign extend chars
- MIPS instruction that doesn't sign extend when loading bytes:
Cel a load byte unsigned:
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## Overflow in Arithmetic (1/2)

- Reminder: Overflow occurs when there is a "mistake" in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):

| 15 | 1111 |
| ---: | ---: |
| +3 | +0011 |
| 18 | 10010 |

- But we don't have room for 5 -bit solution, so the solution would be 0010, which is +2, and "wrong".


## Overflow in Arithmetic (2/2)

- Some languages detect overflow (Ada), some don't (most C implementations)
- MIPS solution is 2 kinds of arithmetic instructs:
- These cause overflow to be detected
- add (add)
- add immediate (addi)
- subtract (sub)
- These do not cause overflow detection
- add unsigned (addu)
- add immediate unsigned (addiu)
- subtract unsigned (subu)
- Compiler selects appropriate arithmetic - MIPS C compilers produce


## Two "Logic" Instructions

- Here are 2 more new instructions
- Shift Left: sll \$s1,\$s2,2 \#s1=s2<<2
- Store in \$s1 the value from $\$$ s2 shifted 2 bits to the left (they fall off end), inserting 0's on right; << in C.
- Before:0000 0002 ${ }_{\text {nex }}$ $00000000000000000000000000000010^{0_{\text {two }}}$
- After: $00000008_{\text {hex }}$ $00000000000000000000000000001000^{\text {two }}$
- What arithmetic effect does shift left have?
- Shift Right: sri is opposite shift; >>


## Loops in C/Assembly (1/3)

- Simple loop in C; A[] is an array of ints

$$
\begin{aligned}
& \text { do }\left\{\begin{array}{l}
\mathrm{g}=\mathrm{g}+\mathrm{A}[\mathrm{i} ;
\end{array} \text { } \mathrm{i}=\mathrm{i}+\mathrm{j} ;\right.
\end{aligned}
$$

\} while (i != h);

- Rewrite this as:

Loop: $\mathrm{g}=\mathrm{g}+\mathrm{A}[\mathrm{i}$;
$\mathrm{i}=\mathrm{i}+\mathrm{j}$;
if ( i ! $=\mathrm{h}$ ) goto Loop;

- Use this mapping:
g, h, i, j,
\$s1, \$s2, §s3, \$s4, \$s5

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## Loops in C/Assembly (2/3)

- Final compiled MIPS code:

Loop: sll \$t1,\$s3,2 \# \$t1= 4*/
addu \$t1,\$t1,\$s5 \# \$t1=addr A+4i
Iw \$t1,0(\$11) \#\$t1=A[i]
addu $\$ \mathbf{s} 1, \$ \mathrm{~s} 1, \$ \mathrm{t} 1 \quad$ \# g=g+A[i] addu $\$ \mathrm{ss} 3, \$ \mathrm{~s} 3, \$ \mathrm{ss} 4$ \# i=i+j bne \$s3,\$s2,Loop \# goto Loop
\# if il=h

- Original code:

$$
\begin{aligned}
& \text { Loop: g = g + [i]; } \\
& i=1+j ; \\
& \text { if ( ! = h) goto Loop; }
\end{aligned}
$$

## Loops in C/Assembly (3/3)

- There are three types of loops in C:
- while
- do ... while
- for
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch


## Administrivia

- HW2 is due Sunday at 23:59:59


## Inequalities in MIPS (1/4)

- Until now, we've only tested equalities ( $==$ and != in C). General programs need to test < and > as well.
- Introduce MIPS Inequality Instruction:
- "Set on Less Than"
- Syntax: slt reg1,reg2,reg3
- Meaning: reg1 = (reg2 < reg3) ;
 "reset" means "change to 0".


## Inequalities in MIPS (2/4)

- How do we use this? Compile by hand: if ( $\mathrm{g}<\mathrm{h}$ ) goto Less; \#g:\$s0, h:\$s1
- Answer: compiled MIPS code...
slt \$t0,\$s0,\$s1 \# \$t0 = 1 if g<h bne \$t0,\$0,Less \# goto Less
\# if \$t0!=0
\# (if (g<h)) Less:
- Register \$0 always contains the value 0, so bne and beq often use it for comparison after an slt instruction.
- A slt $\Rightarrow$ bne pair means if(... < ...)goto...


## Inequalities in MIPS (3/4)

- Now we can implement <, but how do we implement $>, \leq$ and $\geq$ ?
- We could add 3 more instructions, but:
- MIPS goal: Simpler is Better
- Can we implement $\leq$ in one or more instructions using just slt and branches?
- What about >?
- What about $\geq$ ?


## Inequalities in MIPS (4/4)

\# a:\$s0, b:\$s1
slt \$t0,\$s0,\$s1 \# \$t0 = 1 if a<b beq \$t0,\$0,skip \# skip if a >= b <stuff>
\# do if a<b skip:
Two independent variations possible:
Use slt \$0,\$s1,\$so instead of slt \$t0,\$s0,\$s1
Use bne instead of beq

## Immediates in Inequalities

- There is also an immediate version of slt to test against constants: slti
- Helpful in for loops

$$
\text { if ( } g>=1 \text { ) goto Loop }
$$

## Loop:

M
stit \$10,\$s0,1 \#\$t0 = 1 if $\# \$ \mathrm{~s} 0<1(\mathrm{~g}<1)$ beq \$10,\$0,Loop \# goto Loop \# if $\$ 10=0$
\# (if (g>=1))

## What about unsigned numbers?

- Also unsigned inequality instructions:


## sltu, sltiu

... which sets result to 1 or $\mathbf{0}$ depending on unsigned comparisons

- What is value of \$t0, \$t1?
(\$s0 = FFFF FFFA hex , \$s1 = 0000 FFFA $_{\text {hex }}$ )

$$
\begin{array}{r}
\text { slt \$t0, \$s0, \$s1 } \\
\text { sltu \$t1, \$s0, \$s1 }
\end{array}
$$

## MIPS Signed vs. Unsigned - diff meanings!

- MIPS terms Signed/Unsigned "overloaded":
- Do/Don't sign extend
- (lb, lbu)
- Do/Don't overflow
- (add, addi, sub, mult, div)
- (addu, addiu, subu, multu, divu)
- Do signed/unsigned compare
- (slit, slti/sltu, sltiu)


## Peer Instruction

Loop:addi \$s0,\$s0,-1 \# i = i - 1 slti \$t0,\$s1,2 \# \$t0 = ( $j<2$ ) beq \$t0,\$0 ,Loop \# goto Loop if \$t0 == 0 slt \$t0,\$s1,\$s0 \# \$t0 = ( $j<i$ ) bne \$t0,\$0,Loop \# goto Loop if \$t0 != 0
(\$s0=i, \$s1=j)
What C code properly fills in the blank in loop below? do \{i--;\} while(__);

## "And in conclusion..."

- To help the conditional branches make decisions concerning inequalities, we introduce: "Set on Less Than" called slt, slti, sltu, sltiu
- One can store and load (signed and unsigned) bytes as well as words with lb, Ibu
- Unsigned add/sub don't cause overflow
- New MIPS Instructions:
sII, srl, Ib, lbu
slt, slti, sltu, sltiu
addu, addiu, subu


## Bonus Slides

## Example: The C Switch Statement (1/3)

- Choose among four alternatives depending on whether $k$ has the value 0, 1, 2 or 3. Compile this C code:

```
switch (k) {
case 0: f=i+j; break; /* k=0 */
case 1: f=g+h; break; /* k=1 */
case 2: f=g-h; break; /* k=2 */
case 3: f=i-j; break; /* k=3 */
}
```


## Example: The C Switch Statement (2/3)

- This is complicated, so simplify.
- Rewrite it as a chain of if-else statements, which we already know how to compile: if(k==0) f=i+j;
else if(k==1) f=g+h; else if(k==2) f=g-h; else if( $k==3$ ) f=i-j;
- Use this mapping:
f:\$s0, g:\$s1, h:\$s2,
i:\$s3, j:\$s4, k:\$s5


## Example: The C Switch Statement (3/3)

- Final compiled MIPS code:
bne \$s5,\$0,L1 \# branch k!=0 add \$s0,\$s3,\$s4 \#k==0 so f=i+j j Exit \#end of case so Exit
L1: addi \$t0,\$s5,-1 \#\$t0=k-1
bne \$t0,\$0,L2 \# branch k!=1
add $\$ \mathbf{s} 0, \$ s 1, \$ s 2$ \#k==1 so $f=g+h$ j Exit \#end of case so Exit
L2: addi \$t0,\$s5,-2 \# \$t0=k-2
bne \$t0,\$0,L3 \# branch k!=2
sub \$s0,\$s1,\$s2 \#k==2 so f=g-h
j Exit \#end of case so Exit
addi \$t0,\$s5,-3 \# \$t0=k-3
bne \$t0,\$0,
sub $\$$ s0,\$s3,\$s4

