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UC Berkeley CS61C : Machine Structures

Lecture 26 – Combinational Logic Blocks

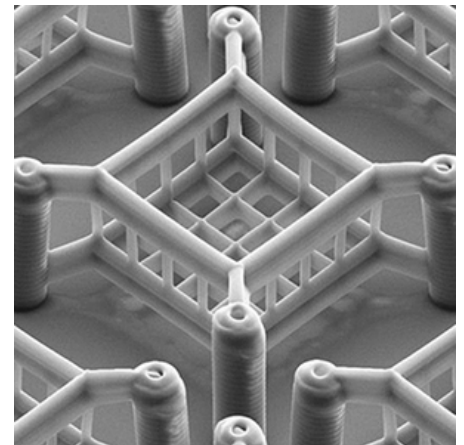


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Very fast 3D Micro Printer ⇒

A new company called Nanoscribe has developed a fabrication device that can create structures like the one at the right at the micro scale in minutes (instead of hours). The idea is that “tiny, ultrashort pulses from a near-infrared laser on a light-sensitive material solidifies on spot. Mirrors not motors



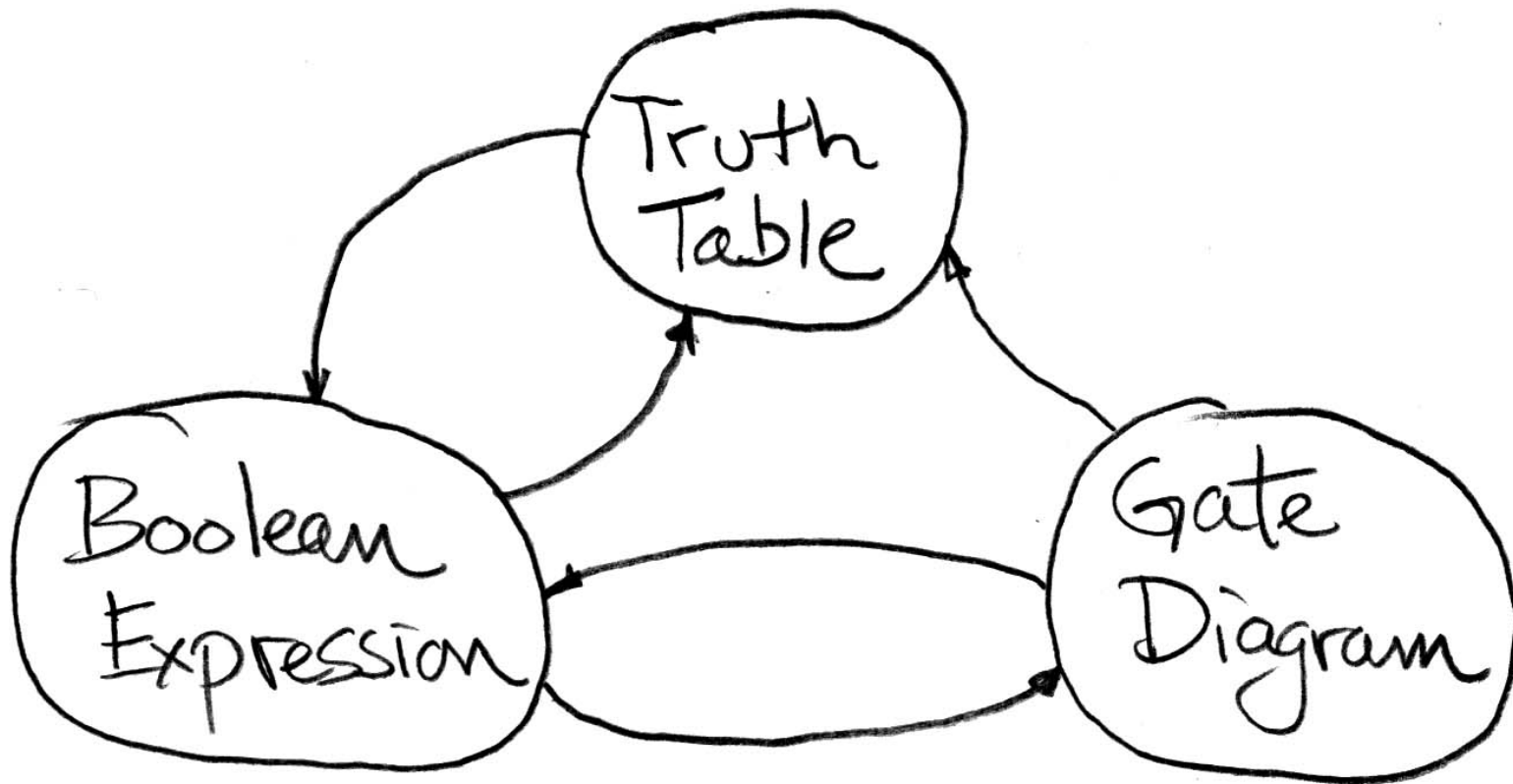
www.technologyreview.com/news/511856/micro-3-d-printer-creates-tiny-structures-in-seconds/

CS61C L26 Combinational Logic Blocks (1)

Garcia, Spring 2014 © UCB

Review

- Use this table and techniques we learned to transform from 1 to another

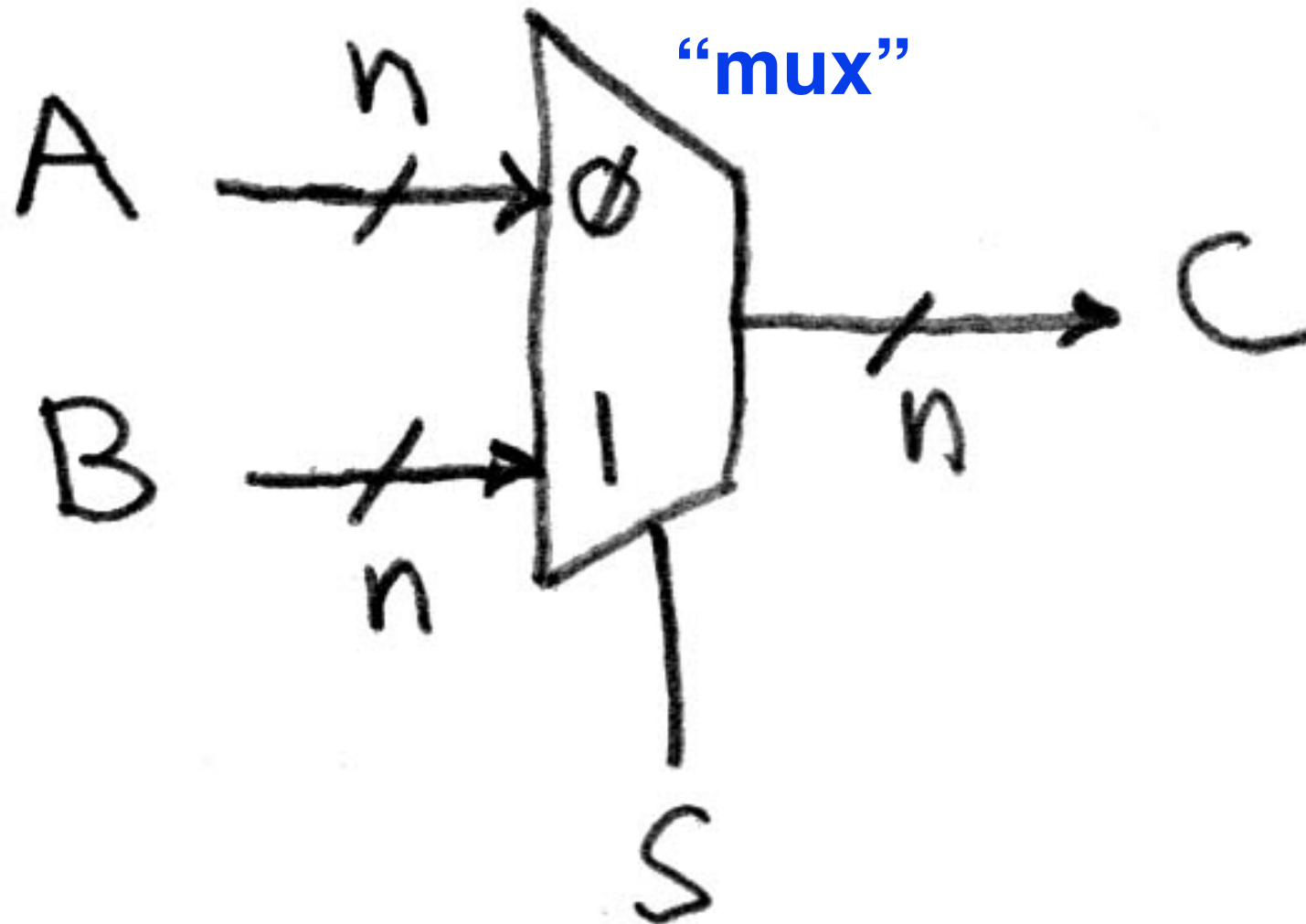


Today

- **Data Multiplexors**
- **Arithmetic and Logic Unit**
- **Adder/Subtractor**

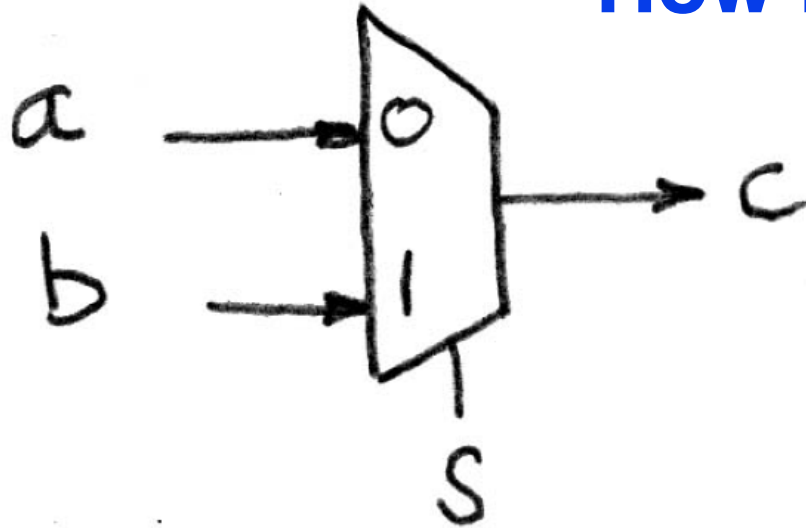


Data Multiplexor (here 2-to-1, n-bit-wide)



N instances of 1-bit-wide mux

How many rows in TT?

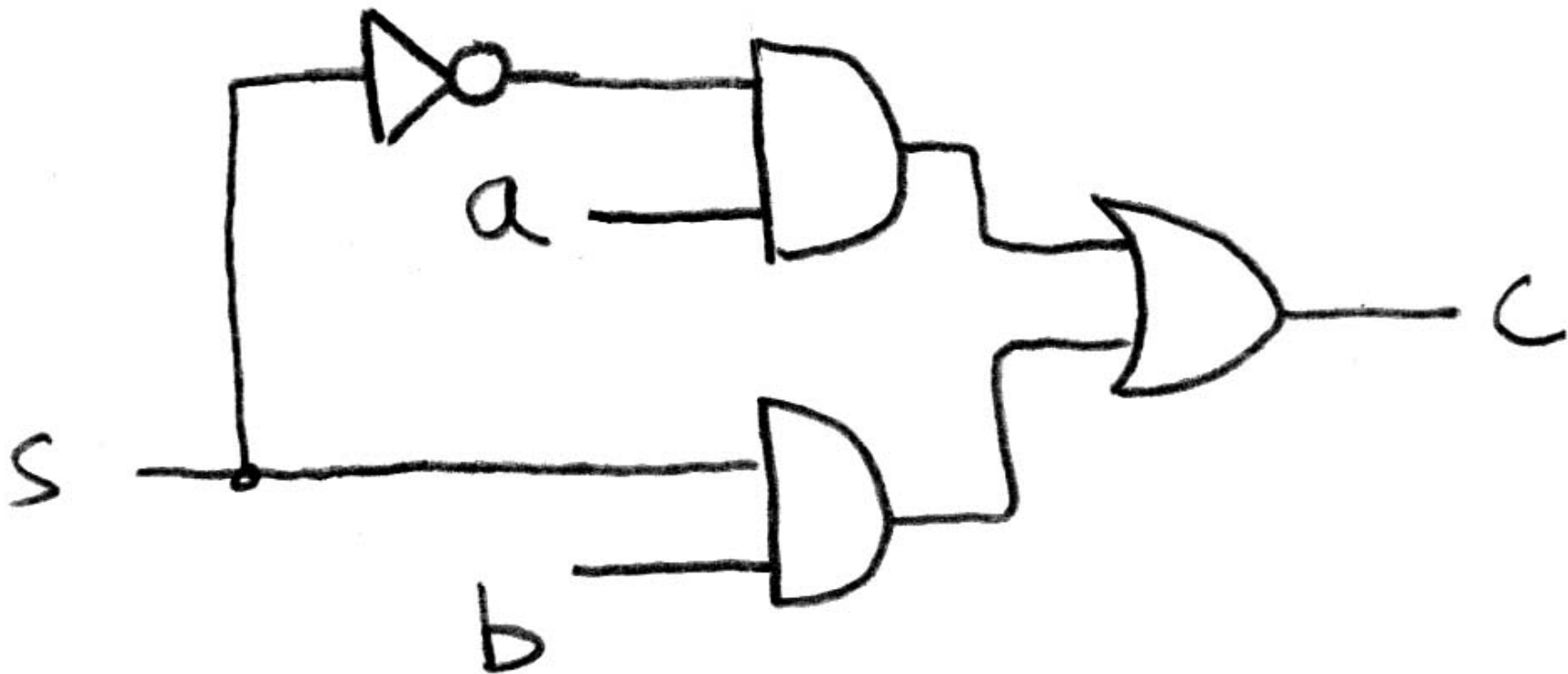


$$\begin{aligned}c &= \bar{s}a\bar{b} + \bar{s}ab + s\bar{a}b + sab \\ &= \bar{s}(a\bar{b} + ab) + s(\bar{a}b + ab) \\ &= \bar{s}(a(\bar{b} + b)) + s((\bar{a} + a)b) \\ &= \bar{s}(a(1) + s((1)b) \\ &= \bar{s}a + sb\end{aligned}$$



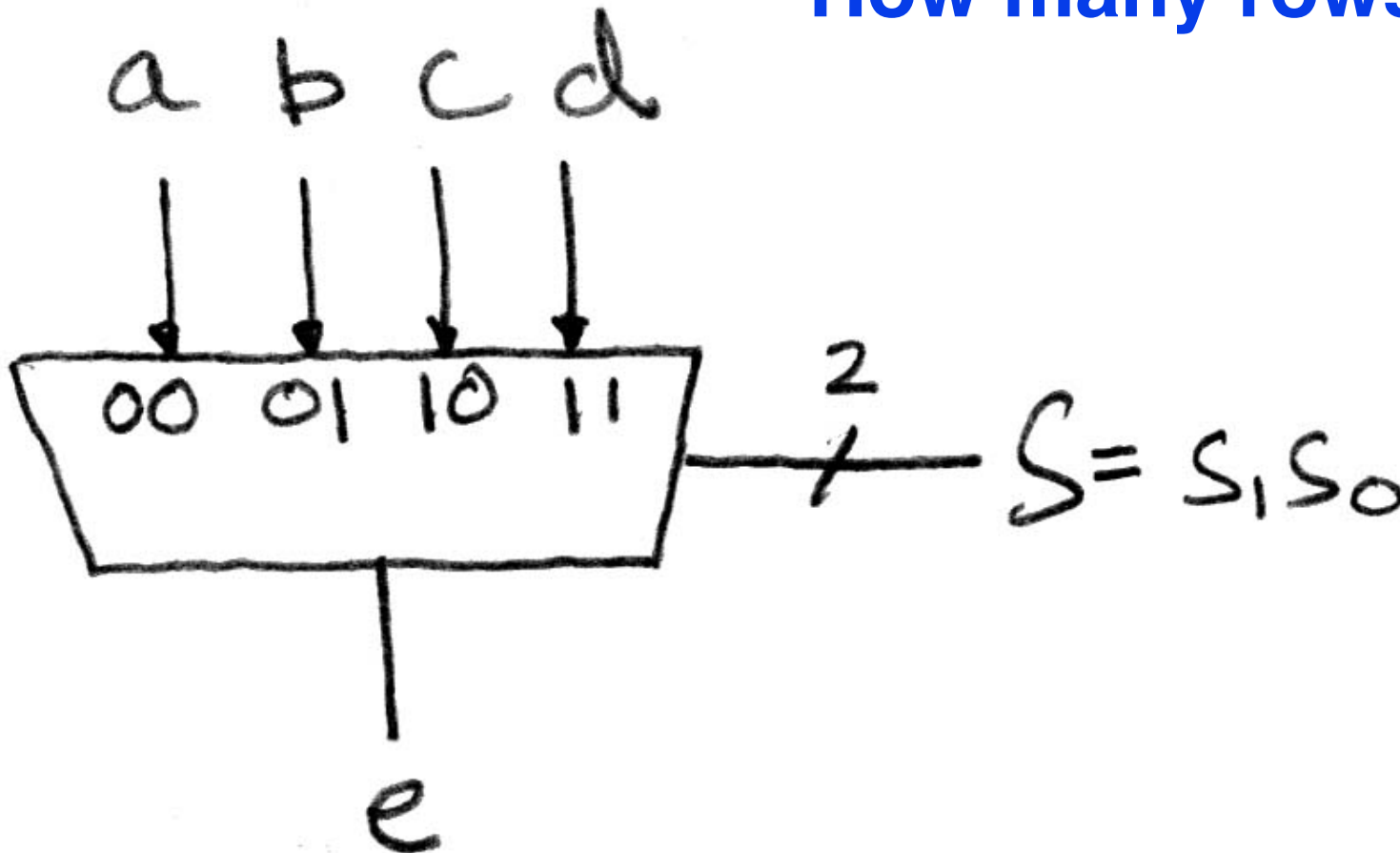
How do we build a 1-bit-wide mux?

$$\bar{s}a + sb$$



4-to-1 Multiplexor?

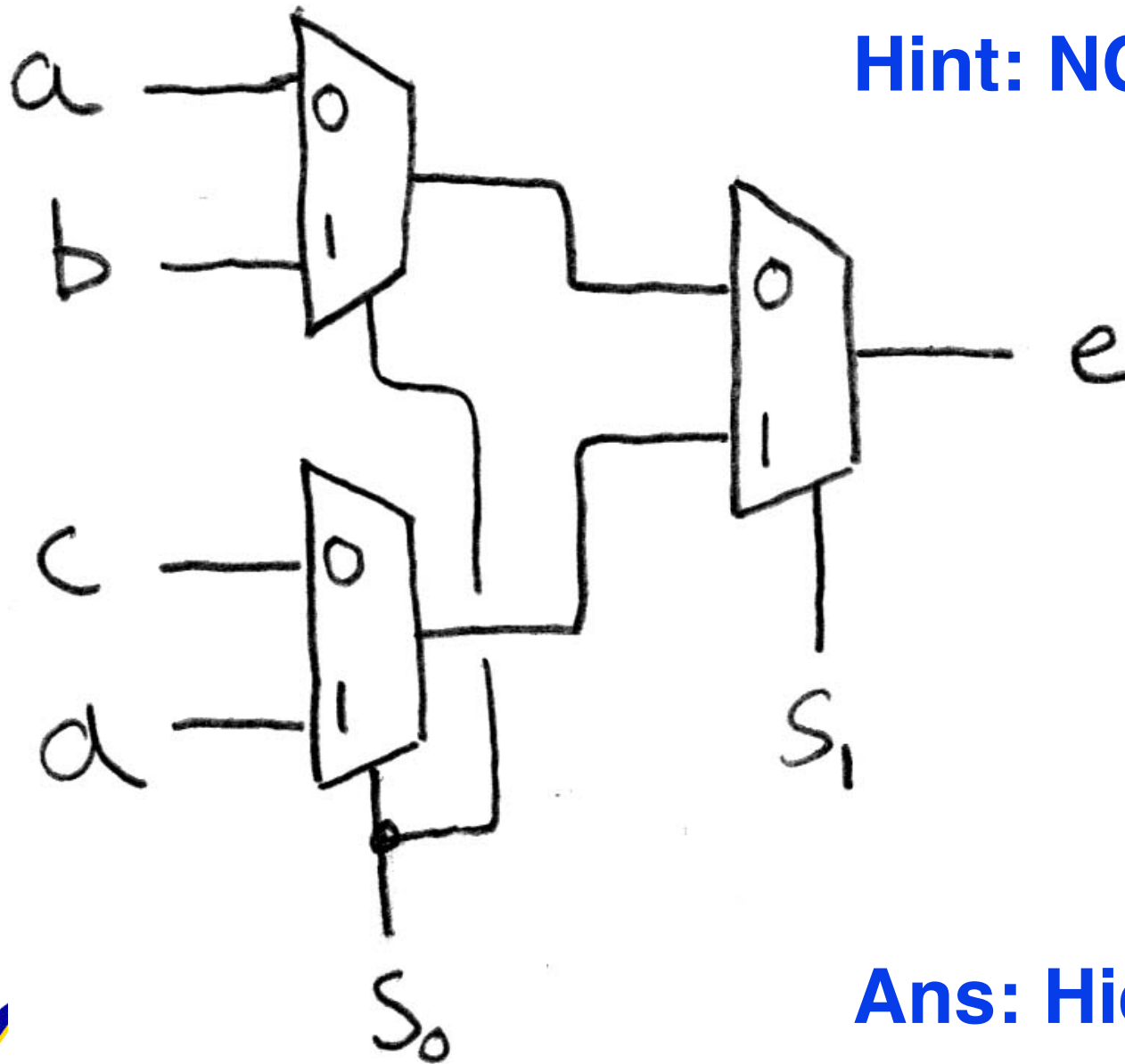
How many rows in TT?



$$e = \bar{s}_1\bar{s}_0a + \bar{s}_1s_0b + s_1\bar{s}_0c + s_1s_0d$$

Is there any other way to do it?

Hint: NCAA tourney!

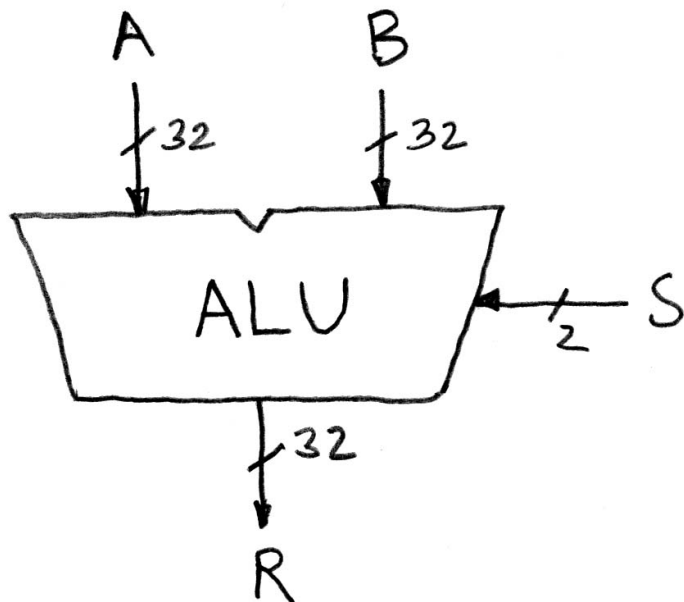


Ans: Hierarchically!



Arithmetic and Logic Unit

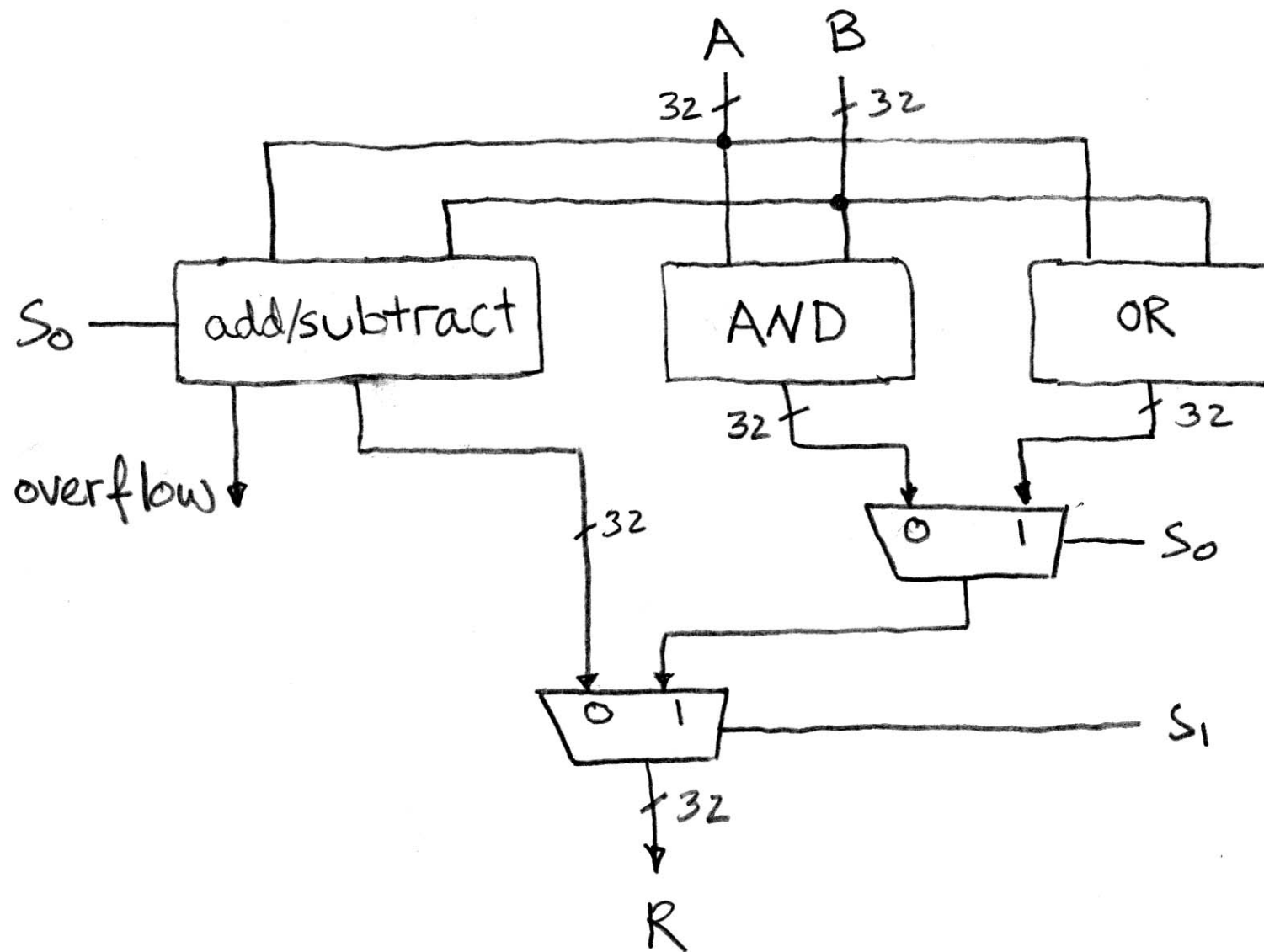
- Most processors contain a special logic block called “Arithmetic and Logic Unit” (ALU)
- We’ll show you an easy one that does ADD, SUB, bitwise AND, bitwise OR



when $S=00$, $R=A+B$
when $S=01$, $R=A-B$
when $S=10$, $R=A \text{ AND } B$
when $S=11$, $R=A \text{ OR } B$



Our simple ALU



Adder/Subtractor Design -- how?

- Truth-table, then determine canonical form, then minimize and implement as we've seen before
- Look at breaking the problem down into smaller pieces that we can cascade or hierarchically layer



Adder/Subtractor – One-bit adder LSB...

$$\begin{array}{cccc} & a_3 & a_2 & a_1 & a_0 \\ + & b_3 & b_2 & b_1 & b_0 \\ \hline s_3 & s_2 & s_1 & s_0 \end{array}$$

a_0	b_0	s_0	c_1
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

$$s_0 =$$

$$c_1 =$$



Adder/Subtractor – One-bit adder (1/2)...

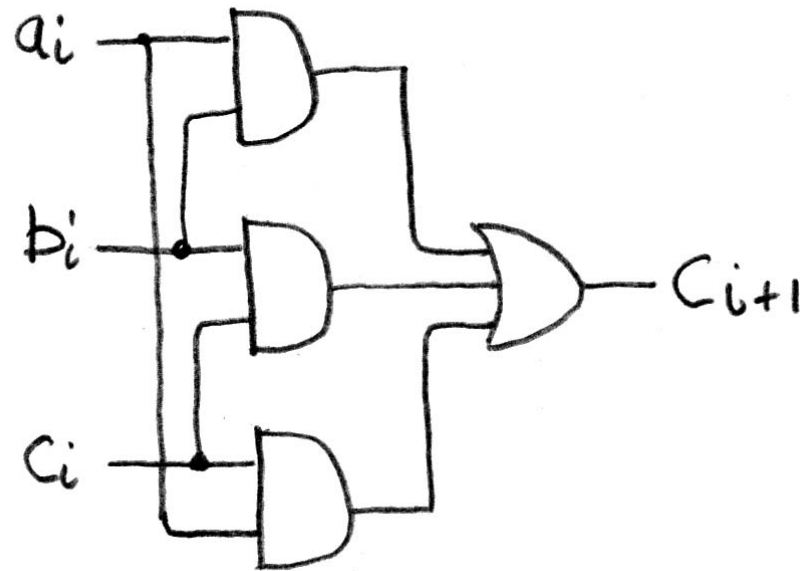
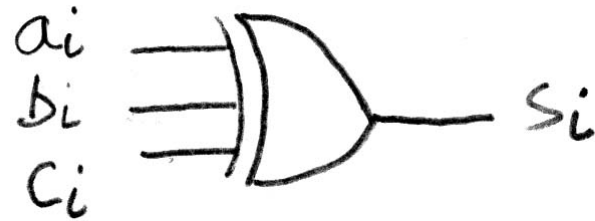
				a_i	b_i	c_i	S_i	C_{i+1}
				0	0	0	0	0
				0	0	1	1	0
				0	1	0	1	0
+	a_3	a_2	a_1	a_0	0	1	1	0
	b_3	b_2	b_1	b_0	0	1	0	1
	s_3	s_2	s_1	s_0	1	0	0	1
				1	0	1	0	1
				1	1	0	0	1
				1	1	1	1	1

$$S_i =$$

$$C_{i+1} =$$

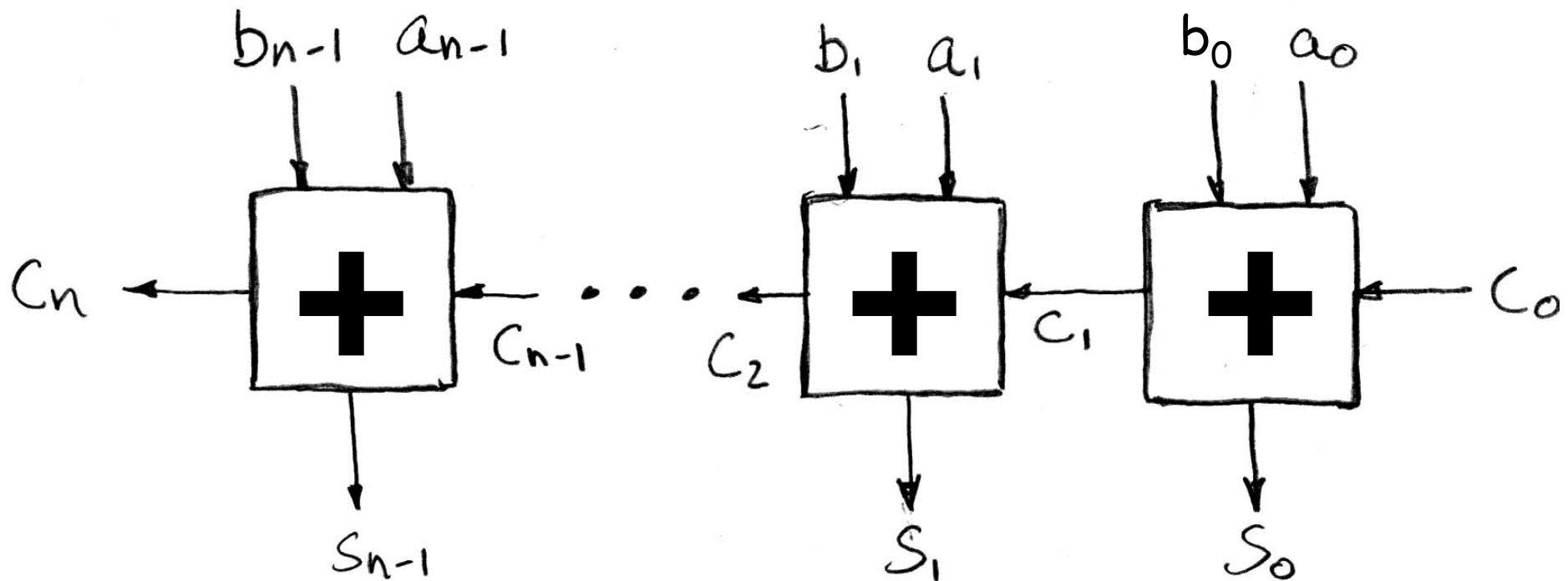


Adder/Subtractor – One-bit adder (2/2)...



$$s_i = \text{XOR}(a_i, b_i, c_i)$$
$$c_{i+1} = \text{MAJ}(a_i, b_i, c_i) = a_i b_i + a_i c_i + b_i c_i$$

N 1-bit adders \Rightarrow 1 N-bit adder

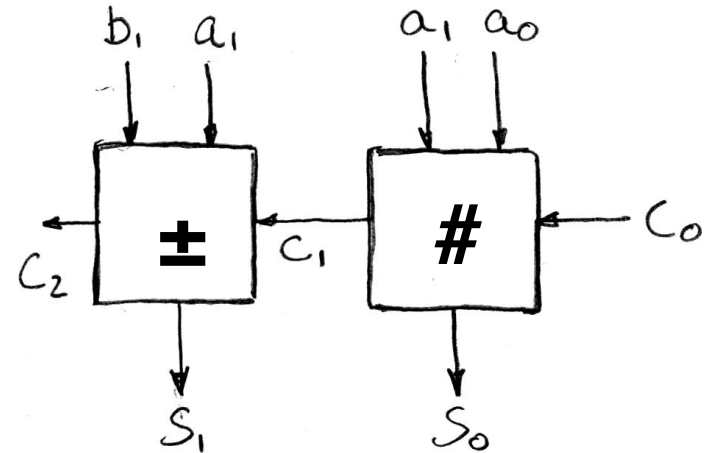


**What about overflow?
Overflow = c_n ?**

What about overflow?

Consider a 2-bit signed # & overflow:

- 10 = -2 + -2 or -1
- 11 = -1 + -2 only
- 00 = 0 NOTHING!
- 01 = 1 + 1 only



Highest adder

- $C_1 = \text{Carry-in} = C_{in}$, $C_2 = \text{Carry-out} = C_{out}$
- No C_{out} or $C_{in} \Rightarrow \text{NO overflow!}$

What op? • C_{in} , and $C_{out} \Rightarrow \text{NO overflow!}$

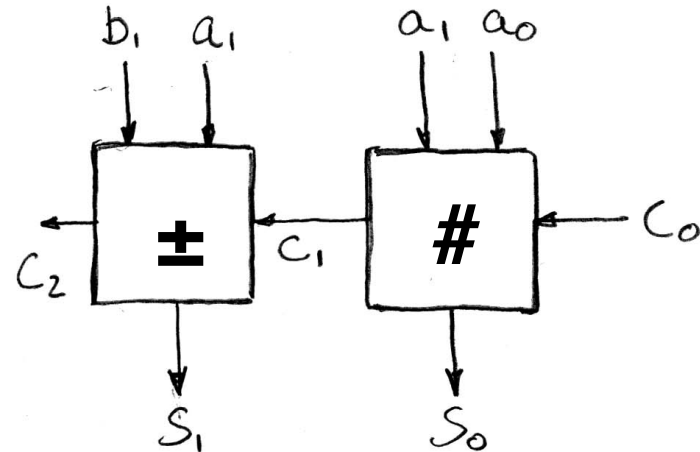
- C_{in} , but no $C_{out} \Rightarrow A, B \text{ both } > 0, \text{ overflow!}$
- C_{out} , but no $C_{in} \Rightarrow A, B \text{ both } < 0, \text{ overflow!}$



What about overflow?

- Consider a 2-bit signed # & overflow:

10 = -2
11 = -1
00 = 0
01 = 1



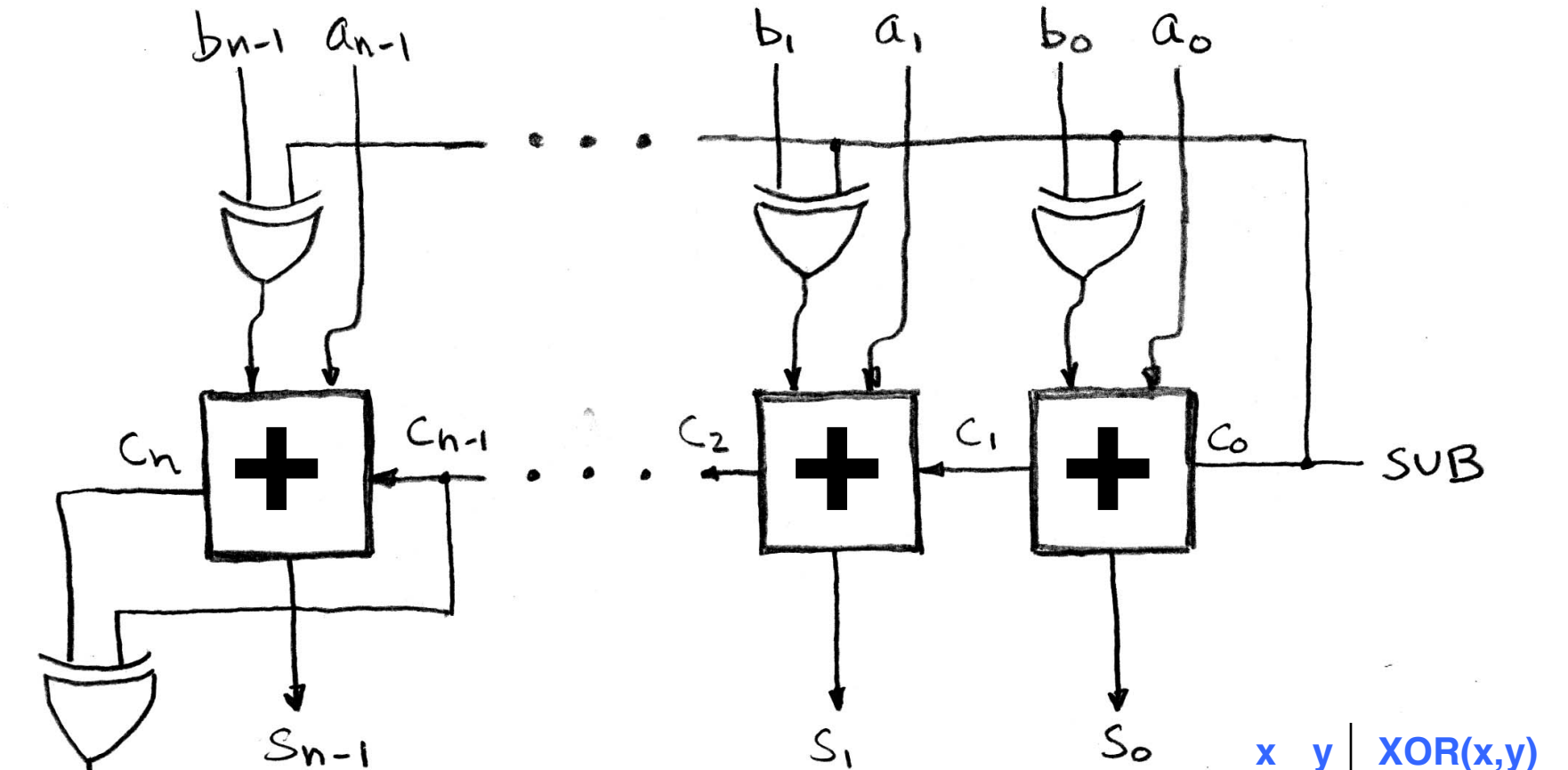
- Overflows when...

- C_{in} , but no $C_{out} \Rightarrow A, B$ both > 0 , overflow!
- C_{out} , but no $C_{in} \Rightarrow A, B$ both < 0 , overflow!

$$\text{overflow} = c_n \text{ XOR } c_{n-1}$$



Extremely Clever Subtractor



**XOR serves as
conditional inverter!**

x	y	XOR(x,y)
0	0	0
0	1	1
1	0	1
1	1	0



Peer Instruction

- 1) Truth table for mux with 4-bits of signals has 2^4 rows
- 2) We could cascade N 1-bit shifters to make 1 N-bit shifter for sll, srl

	12
a)	FF
b)	FT
c)	TF
d)	TT



Peer Instruction Answer

- 1) Truth table for mux with 4-bits of signals controls 16 inputs, for a total of 20 inputs, so truth table is 2^{20} rows... **FALSE**
- 2) We could cascade N 1-bit shifters to make 1 N-bit shifter for sll, srl ... **TRUE**

1) Truth table for mux with 4-bits of signals is 2^4 rows long

2) We could cascade N 1-bit shifters to make 1 N-bit shifter for sll, srl

	12
a)	FF
b)	FT
c)	TF
d)	TT



“And In conclusion...”

- **Use muxes to select among input**
 - **S input bits selects 2^S inputs**
 - **Each input can be n-bits wide, indep of S**
- **Can implement muxes hierarchically**
- **ALU can be implemented using a mux**
 - **Coupled with basic block elements**
- **N-bit adder-subtractor done using N 1-bit adders with XOR gates on input**
 - **XOR serves as conditional inverter**

