## CS61c Spring 2015 Discussion 2 - C Memory Management \& MIPS

## 1 C Memory Management

1. In which memory sections (CODE, STATIC, HEAP, STACK) do the following reside?
```
#define C 2
const int val = 16;
char arr[] = "foo";
void foo(int arg){
    char *str = (char *) malloc (C*val);
    char *ptr = arr;
}
```

| $\arg [$ | $\operatorname{str}[$ | $]$ |  |
| :--- | :--- | :--- | :--- |
| $\operatorname{arr}[$ | $\ln [$ | $\operatorname{str}[$ | $]$ |
| $\operatorname{val}[$ | C | $[$ | $]$ |

2. What is wrong with the C code below?
```
int* ptr = malloc(4 * sizeof(int));
if(extra_large) ptr = malloc(10 * sizeof(int));
return ptr;
```

3. Write code to prepend (add to the start) to a linked list, and to free/empty the entire list.
struct ll_node \{ struct ll_node* next; int value; \}

| free_ll(struct ll_node** list) | prepend(struct ll_node** list, int value) |
| :--- | :--- |
|  |  |
|  |  |
|  |  |

Note: list points to the first element of the list, or to NULL if the list is empty.

## 2 MIPS Intro

1. Assume we have an array in memory that contains int* arr $=\{1,2,3,4,5,6,0\}$. Let the value of arr be a multiple of 4 and stored in register $\$ \mathbf{s} 0$. What do the following programs do?
a) $1 \mathrm{w} \$ \mathrm{t} 0,12(\$ \mathrm{~s} 0)$ add \$t1, \$t0, \$s0 sw \$t0, 4(\$t1)
b) addiu \$s1, \$s0, 27
lh \$t0, $-3(\$ \mathrm{~s} 1)$
c) addiu $\$$ s1, $\$$ s0, 24
lh \$t0\$, -3(\$s1)
```
d) addiu $t0, $0, 12
    sw $t0, 6($s0)
    e) addiu $t0, $0, 8
    sw $t0, -4($s0)
f) addiu $s1, $s0, 10
    addiu $t0, $0, 6
    sw $t0, 2($s1)
```

2. In 1), what other instructions could be used in place of each load/store without alignment errors?
3. What are the instructions to branch to label: on each of the following conditions?

| \$s0 < \$s1 | \$s0 <= \$s1 | \$s0 > 1 | \$s0 >= 1 |
| :--- | :--- | :--- | :--- |
|  |  |  |  |
|  |  |  |  |

## 3 Translating between C and MIPS

Translate between the C and MIPS code. You may want to use the MIPS Green Sheet as a reference. In all of the C examples, we show you how the different variables map to registers - you don't have to worry about the stack or any memory-related issues.

| C | MIPS |
| :---: | :---: |
| ```// $s0 -> a, $s1 -> b // $s2 -> c, $s3 -> z int a = 4, b = 5, c = 6, z; z = a + b + c + 10;``` |  |
| ```// $s0 -> int * p = intArr; // $s1 -> a; *p = 0; int a = 2; p[1] = p[a] = a;``` |  |
| $\begin{aligned} & \text { // \$s0 -> a, \$s1 -> b } \\ & \text { int } a=5, b=10 ; \\ & \text { if }(a+a==b)\{ \\ & \quad a=0 ; \\ & \} \text { else \{ } \\ & \quad b=a-1 ; \end{aligned}$ |  |
|  | ```addiu $s0, $0, 0 addiu $s1, $0, 1 addiu $t0, $0, 30 loop: beq $s0, $t0, exit addu $s1, $s1, $s1 addiu $s0, $s0, 1 j loop exit:``` |
| ```// $a0 -> n, $v0 -> sum int sum; for(sum=0;n>0;sum+=n--);``` |  |

