**An Abstract View of the Critical Path**

Critical Path (Load Operation) =
- Delay clock through PC (FFs) +
- Instruction Memory’s Access Time +
- Register File’s Access Time +
- ALU to Perform a 32-bit Add +
- Data Memory Access Time +
- Stable Time for Register File Write

**Improve Critical Path ➔ Improve Clock**

- “Critical path” (longest path through logic) determines length of clock period
- To reduce clock period ➔ decrease path through CL by inserting State

**Review: Single cycle datapath**

- 5 steps to design a processor
  1. Analyze instruction set ➔ datapath requirements
  2. Select set of datapath components & establish clock methodology
  3. Assemble datapath meeting the requirements
  4. Analyze implementation of each instruction to determine setting of control points that effects the register transfer
  5. Assemble the control logic

- *MIPS is the hard part*
- *MIPS makes that easier*
  - Instructions same size
  - Source registers always in same place
  - Immediates same size, location
  - Operations always on registers/immediates

**Review Datapath (1/3)**

- Datapath is the hardware that performs operations necessary to execute programs.
- Control instructs datapath on what to do next.
- Datapath needs:
  - access to storage (general purpose registers and memory)
  - computational ability (ALU)
  - helper hardware (local registers and PC)
Review Datapath (3/3)

1. Instruction Fetch
2. Decode/Register Read
3. Execute
4. Memory
5. Write Back

Gotta Do Laundry

- Ann, Brian, Cathy, Dave each have one load of clothes to wash, dry, fold, and put away
- Washer takes 30 minutes
- Dryer takes 30 minutes
- "Folder" takes 30 minutes
- "Stasher" takes 30 minutes to put clothes into drawers

Sequential Laundry

- Sequential laundry takes 8 hours for 4 loads

Pipelined Laundry

- Pipelined laundry takes 3.5 hours for 4 loads!

General Definitions

- Latency: time to completely execute a certain task
  - for example, time to read a sector from disk is disk access time or disk latency
  - Instruction latency is time from when instruction starts to time when it finishes.
- Throughput: amount of work that can be done over a period of time

Pipelining Lessons (0/2)

- Terminology:
  - Issue: When instruction goes into first stage of pipe.
  - Commit: when instruction finishes last stage
Pipelining Lessons (1/2)

1. Pipelining doesn't help latency of single task, it helps throughput of entire workload.
2. Multiple tasks operating simultaneously using different resources.
3. Potential speedup = Number pipe stages
4. Time to “fill” pipeline and time to “drain” it reduces speedup: 2.3X v. 4X in this example.

Steps in Executing MIPS

1. **IFetch**: Fetch Instruction, Increment PC
2. **Decode**: Instruction, Read Registers
3. **Execute**: Mem-ref: Calculate Address, Arith-log: Perform Operation
4. **Memory**: Load: Read Data from Memory, Store: Write Data to Memory
5. **Write Back**: Write Data to Register

Review: Datapath for MIPS

- Use datapath figure to represent pipeline

Pipelining Lessons (2/2)

1. Suppose new Washer takes 20 minutes, new Stasher takes 20 minutes. How much faster is pipeline?
2. Pipeline rate limited by slowest pipeline stage.
3. Unbalanced lengths of pipe stages also reduces speedup.

Graphical Pipeline Representation

(In Reg, right half highlight read, left half write)

- Every instruction must take same number of steps, also called pipeline “stages”, so some will go idle sometimes.

Pipelined Execution Representation

- Every instruction must take same number of steps, also called pipeline “stages”, so some will go idle sometimes.
Example

• Suppose 2 ns for memory access, 2 ns for ALU operation, and 1 ns for register file read or write; compute instruction throughput

• Nonpipelined Execution:
  • lw: IF + Read Reg + ALU + Memory + Write Reg = 2 + 1 + 2 + 2 + 1 = 8 ns
  • add: IF + Read Reg + ALU + Write Reg = 2 + 1 + 2 + 1 = 6 ns

• Pipelined Execution:
  • Max(IF, Read Reg, ALU, Memory, Write Reg) = 2 ns

Example

• Suppose 2 ns for memory access, 2 ns for ALU operation, and 1 ns for register file read or write; compute instruction latency

• Nonpipelined Execution:
  • lw: IF + Read Reg + ALU + Memory + Write Reg = 2 + 1 + 2 + 2 + 1 = 8 ns
  • add: IF + Read Reg + ALU + Write Reg = 2 + 1 + 2 + 1 = 6 ns

• Pipelined Execution:
  • SUM(IF, Read Reg, ALU, Memory, Write Reg) = 10 ns

Things to Remember

• Optimal Pipeline
  • Each stage is executing part of an instruction each clock cycle.
  • One instruction finishes during each clock cycle.
  • On average, executes far more quickly.

• What makes this work?
  • Similarities between instructions allow us to use same stages for all instructions (generally).
  • Each stage takes about the same amount of time as all others: little wasted time.

Pipeline Summary

• Pipelining is a BIG IDEA
  • widely used concept

  • What makes it less than perfect? …

Administrivia

• Project 2 – Sunday
• HW6 out now
  • Due next Tuesday

Pipeline Hazard: Matching socks in later load

A depends on D; stall since folder tied up
Problems for Computers

- Limits to pipelining: **Hazards** prevent next instruction from executing during its designated clock cycle
  - **Structural hazards**: HW cannot support this combination of instructions (single person to fold and put clothes away)
  - **Control hazards**: Pipelining of branches & other instructions stall the pipeline until the hazard; “bubbles” in the pipeline
  - **Data hazards**: Instruction depends on result of prior instruction still in the pipeline (missing sock)

**Structural Hazard #1: Single Memory (1/2)**

- Solution:
  - Infeasible and inefficient to create second memory
  - (We’ll learn about this more next week)
  - So simulate this by having two Level 1 Caches (a temporary smaller [of usually most recently used] copy of memory)
  - Have both an L1 Instruction Cache and an L1 Data Cache
  - Requires complex hardware to control when both caches miss!

**Structural Hazard #1: Single Memory (2/2)**

**Structural Hazard #2: Registers (1/2)**

- Fact: Register access is VERY fast: takes less than half the time of ALU stage
- Solution: introduce convention
  - Always write to registers during first half of each clock cycle
  - Always read from registers during second half of each clock cycle (easy when async)
  - Result: can perform read and write during same clock cycle

**Structural Hazard #2: Registers (2/2)**

**Control Hazard: Branching (1/7)**

- Where do we do the compare for the branch?
Control Hazard: Branching (2/7)

• We put branch decision-making hardware in ALU stage
  • therefore two more instructions after the branch will always be fetched, whether or not the branch is taken

• Desired functionality of a branch
  • if we do not take the branch, don’t waste any time and continue executing normally
  • if we take the branch, don’t execute any instructions after the branch, just go to the desired label

Control Hazard: Branching (3/7)

• Initial Solution: Stall until decision is made
  • insert “no-op” instructions: those that accomplish nothing, just take time
  • Drawback: branches take 3 clock cycles each (assuming comparator is put in ALU stage)
  • Drawback: Will still fetch inst at branch+4. Must either decode branch in IF or squash fetched branch+4.

Control Hazard: Branching (4/7)

• Optimization #1:
  • move asynchronous comparator up to Stage 2
  • as soon as instruction is decoded (Opcode identifies is as a branch), immediately make a decision and set the value of the PC (if necessary)
  • Benefit: since branch is complete in Stage 2, only one unnecessary instruction is fetched, so only one no-op is needed
  • Side Note: This means that branches are idle in Stages 3, 4 and 5.

Control Hazard: Branching (5/7)

• Insert a single no-op (bubble)

<table>
<thead>
<tr>
<th>Instruction</th>
<th>Time (clock cycles)</th>
</tr>
</thead>
<tbody>
<tr>
<td>add</td>
<td></td>
</tr>
<tr>
<td>beq</td>
<td></td>
</tr>
<tr>
<td>lw</td>
<td></td>
</tr>
</tbody>
</table>

• Impact: 2 clock cycles per branch instruction ⇒ slow

Control Hazard: Branching (6/7)

• Optimization #2: Redefine branches
  • Old definition: if we take the branch, none of the instructions after the branch get executed by accident
  • New definition: whether or not we take the branch, the single instruction immediately following the branch gets executed (called the branch-delay slot)

Control Hazard: Branching (7/7)

• Notes on Branch-Delay Slot
  • Worst-Case Scenario: can always put a no-op in the branch-delay slot
  • Better Case: can find an instruction preceding the branch which can be placed in the branch-delay slot without affecting flow of the program
    - re-ordering instructions is a common method of speeding up programs
    - compiler must be very smart in order to find instructions to do this
    - usually can find such an instruction at least 50% of the time
  • Jumps also have a delay slot…
Example: Nondelayed vs. Delayed Branch

Nondelayed Branch

<table>
<thead>
<tr>
<th>Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>or $8, $9, $10</td>
</tr>
<tr>
<td>add $1, $2, $3</td>
</tr>
<tr>
<td>sub $4, $5, $6</td>
</tr>
<tr>
<td>beh $1, $4, Exit</td>
</tr>
<tr>
<td>xor $10, $1, $11</td>
</tr>
</tbody>
</table>

Delayed Branch

<table>
<thead>
<tr>
<th>Instruction</th>
</tr>
</thead>
<tbody>
<tr>
<td>add $1, $2, $3</td>
</tr>
<tr>
<td>sub $4, $5, $6</td>
</tr>
<tr>
<td>beh $1, $4, Exit</td>
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<td>or $8, $9, $10</td>
</tr>
<tr>
<td>xor $10, $1, $11</td>
</tr>
</tbody>
</table>

Data Hazards (1/2)

- Consider the following sequence of instructions

  ```
  add $t0, $t1, $t2
  sub $t4, $t0, $t3
  and $t5, $t0, $t6
  or $t7, $t0, $t8
  xor $t9, $t0, $t10
  ```

Data Hazard Solution: Forwarding

Fix by forwarding result as soon as we have it to where we need it:

```
add $t0, $t1, $t2
sub $t4, $t0, $t3
and $t5, $t0, $t6
or $t7, $t0, $t8
xor $t9, $t0, $t10
```

“or” hazard solved by register hardware

Data Hazard: Loads (1/4)

- Forwarding works if value is available (but not written back) before it is needed. But consider...

```
lw $t0, 0($t1)
sub $t3, $t0, $t2
and $t5, $t0, $t4
or $t7, $t0, $t8
xor $t9, $t0, $t10
```

Need result before it is calculated!
Must stall use (sub) 1 cycle and then forward...

Data Hazard: Loads (2/4)

- Hardware must stall pipeline
- Called “interlock”

```
lw $t0, 0($t1)
sub $t3, $t0, $t2
and $t5, $t0, $t4
or $t7, $t0, $t6
```
Data Hazard: Loads (3/4)

- Instruction slot after a load is called "load delay slot"
- If that instruction uses the result of the load, then the hardware interlock will stall it for one cycle.
- If the compiler puts an unrelated instruction in that slot, then no stall
- Letting the hardware stall the instruction in the delay slot is equivalent to putting a nop in the slot (except the latter uses more code space)

Data Hazard: Loads (4/4)

- Stall is equivalent to nop

\[ \text{lw} \; \$t0, \; 0(\$t1) \]
\[ \text{sub} \; \$t3, \$t0, \$t2 \]
\[ \text{and} \; \$t5, \$t0, \$t4 \]
\[ \text{or} \; \$t7, \$t0, \$t6 \]

C.f. Branch Delay vs. Load Delay

- Load Delay occurs only if necessary (dependent instructions).
- Branch Delay always happens (part of the ISA).
- Why not have Branch Delay interlocked?
  - Answer: Interlocks only work if you can detect hazard ahead of time. By the time we detect a branch, we already need its value ... hence no interlock is possible!

Historical Trivia

- First MIPS design did not interlock and stall on load-use data hazard
- Real reason for name behind MIPS: Microprocessor without Interlocked Pipeline Stages
  - Word Play on acronym for Millions of Instructions Per Second, also called MIPS
  - Load/Use ➔ Wrong Answer!

"And in Conclusion...

- Pipeline challenge is hazards
- Forwarding helps w/many data hazards
- Delayed branch helps with control hazard in 5 stage pipeline