### CS 61C Midterm

7/23/2007

<table>
<thead>
<tr>
<th>Last Name</th>
<th>Key</th>
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<tbody>
<tr>
<td>First Name</td>
<td>Answer</td>
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<tr>
<td>SID</td>
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<tr>
<td>Login</td>
<td>cs61c-</td>
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<td>First Letter of Login</td>
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<td>Second Letter of Login</td>
<td>a b c d e f g h i j k l m n o p q r s t u v w x y z</td>
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<tr>
<td>Lab Time</td>
<td>101</td>
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</tbody>
</table>

| Name of the person to your left |
| Name of the person to your right |

All the work is my own. I had no prior knowledge of the exam contents nor will I share the contents with others in CS61C who have not taken it yet. (please sign)

### Instruction

- This booklet contains 7 numbered pages including the cover page. Put all answers on these pages (feel free to use the back of any page for scratch work); don’t hand in any stray pieces of paper.
- Please turn off all pagers, cell phones & beepers. Remove all hats & headphones. Place your backpacks, laptops and jackets at the front. Sit in every other seat. Nothing may be placed in the “no fly zone” spare seat/desk between students.
- Fill in the front of this page and put your name & login on every sheet of paper for Question 0.
- You have 180 minutes to complete this exam. The exam is closed book, no computers, PDAs or calculators. You may use one pages (US Letter, front and back) of notes, plus the green reference sheet from COD 3/e.
- There may be partial credit for incomplete answers; write as much of the solution as you can. We will deduct points if your solution is far more complicated than necessary. When we provide a blank, please fit your answer within the space provided. You have 3 hours...relax.

<table>
<thead>
<tr>
<th>Question</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>Total</th>
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<tbody>
<tr>
<td>Points</td>
<td>1</td>
<td>13</td>
<td>16</td>
<td>16</td>
<td>14</td>
<td>15</td>
<td>75</td>
</tr>
<tr>
<td>Score</td>
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<td>13</td>
<td>16</td>
<td>16</td>
<td>14</td>
<td>15</td>
<td>75</td>
</tr>
</tbody>
</table>

1/7
**Question 1: The CS61C Grab Bag (13pts, 20min)**

a) How many bits does it take to address N things?  
(hint: you may use the ceiling or floor function)

\[\text{ceil} (\log (N))\]

b) Let: \(x = 0xfde7\) \(z = 0x80d5\)

i. What is \(x\) in binary?

\[1111 \ 1101 \ 1110 \ 0111\]

ii. What is the value of \(x\) if it was interpreted as a 2’s Complement?  
(answer in decimal)

\[-537\]

iii. What is the encoding of the result of part ii in 1’s Complement?  
(answer in hex)

\[fde6\]

iv. Let \(x + y = z\). Interpreting \(z\) as a sign-magnitude number, what 1’s Complement encoding of \(y\) will make this true?  
(answer in hex)

\[0144\]

c) Match all the choices on the right with the most appropriate items on the left. Some choices will not be used, and others may be used more than once. If more than one choice on the right matches an item on the left, list all of those choices.

<table>
<thead>
<tr>
<th></th>
<th>Next-fit</th>
<th>A. (2^{38})</th>
</tr>
</thead>
<tbody>
<tr>
<td>K</td>
<td>Copying garbage collection</td>
<td>B. Makes request for memory</td>
</tr>
<tr>
<td>I</td>
<td>Linker</td>
<td>C. Uses static predefined chunks</td>
</tr>
<tr>
<td>D, L</td>
<td>Assembler</td>
<td>D. Resolves jump statements</td>
</tr>
<tr>
<td>F, J</td>
<td>Buddy Memory Allocation</td>
<td>E. Converts C to MAL</td>
</tr>
<tr>
<td>H</td>
<td>Loader</td>
<td>F. Resolves branch statements</td>
</tr>
<tr>
<td>B</td>
<td>Compiler</td>
<td>G. (2^{27})</td>
</tr>
<tr>
<td>E</td>
<td>256 Gibi</td>
<td>H. Can vary the size of chunks</td>
</tr>
<tr>
<td>A</td>
<td></td>
<td>I. Wastes half the space</td>
</tr>
<tr>
<td></td>
<td></td>
<td>J. Makes the relocation table</td>
</tr>
<tr>
<td></td>
<td></td>
<td>K. Attempt to improve First-fit</td>
</tr>
<tr>
<td></td>
<td></td>
<td>L. Makes executable</td>
</tr>
</tbody>
</table>

d) Suppose you want to add a new pseudoinstruction to your assembler to operate on memory directly. The syntax of the instruction is:

\text{maddi register immediate (ie maddi $t2, 14)}

It will add the given immediate to the value in memory that the register holds the address for. You may assume that the absolute value of the immediate is less than \(2^{15}\). Write out what the instruction to the right translates to in TAL.

\[\text{maddi}\ $s0, \ 0x\text{beef}\]

\[\text{lw}\ $at, \ 0($s0)\]

\[\text{addi}\ $at, \ $at, \ 0x\text{beef}\]

\[\text{sw}\ $at, \ 0($s0)\]
**Question 2: The Matrix (16pts, 25min)**

In this problem you will help write some functions to deal with matrices in the mathematical sense. Our ultimate goal is to generate an add function for two matrices. You can use `strlen` and `strcpy` from `string.h`, but no other library functions (reference info on next page). For all of the following parts you may not need to use of all the variables declared, but you may **not** declare any additional variables.

```c
struct matrixStruct {
    int *data;
    char *name;
    int numRows;
    int numCols;
};

typedef struct matrixStruct *matrix;
```

a) In one line of C complete the `getElement` function. You are guaranteed valid input, and you need to return the value in the matrix `A` at row `r` and column `c`. The matrix numbering starts with 0, so if `r=0` and `c=0` it is the upper left entry.

```c
int getElement(matrix A, int r, int c) {
    return (A->data)[c + r*(A->numCols)];
}(many possible ways)
```

b) We want to generate a string that represents the name of the sum matrix. The string should be the name of `A`, concatenated with " + " (note the spaces before and after the ‘+’), and concatenated with the name of `B`. You are guaranteed valid input. Complete `getSumName`.

Example: `A->name = "PB"   B->name="J"` result of `getSumName(A,B) = "PB + J"

```c
char* getSumName(matrix A, matrix B) {
    int newLength, i;
    char *newName, *temp;
    newLength = strlen(A->name) + strlen(B->name) + 4;
    newName = (char*) malloc(sizeof(char) * newLength);
    temp = newName;
    temp = strcpy(temp, A->name);
    temp += strlen(A->name);
    temp = strcpy(temp, " + ");
    temp+=3;
    temp = strcpy(temp, B->name);

    return newName;
}
```
Question 2 (continued): The Matrix (16pts, 25min)
c) Now that we have helping functions, let’s write out the add function. It will be given two matrices \( A \) and \( B \), and should return their sum. If the input matrices’ dimensions do not agree, they can’t be added so return \( \text{NULL} \). You are guaranteed non-NULL input matrices. The result matrix should have its name set as described in Part B. You may not declare any new variables or functions.

```c
matrix add(matrix A, matrix B) {
    int i, j;
    matrix result;
    if((A->numRows != B->numRows)||(A->numCols != B->numCols))
        return NULL;
    result = (matrix) malloc(sizeof(struct matrixStruct));
    result->numRows = A->numRows;
    result->numCols = A->numCols;
    result->name = getSumName(A, B);
    result->data = (int* ) malloc(sizeof(int) * result->numRows * result->numCols);
    for(i = 0; i < result->numRows; i++)
        for(j = 0; j < result->numCols; j++)
            (result->data)[j + i*(A->numCols)] =
                getElement(A, i, j) + getElement(B,i, j);
    return result;
}
```

d) Fill in the function below to free a matrix. You have no guarantee of the integrity of \( A \).

```c
void freeMatrix(matrix A) {
    if(A==NULL)
        return;
    free(A->data);
    free(A->name);
    free(A);
}
```

String Reference (copied from K & R)

- `char *strcpy(s, ct)`:
  copy string \( ct \) to string \( s \), including ‘\0’ return \( s \)

- `size_t strlen(cs)`:
  return length of \( cs \)
**Question 3: Floating Away (16pts, 35min)**

A new hybrid floating point standard has been developed that uses 24 bits. It uses the same tricks and style as the IEEE 754 standard, but does it with different amounts of bits. This standard uses 24 bits, with the first bit being the sign, the next 6 bits being the exponent, and the last 17 bits being the significand (mantissa). When expressing values that aren't part of an encoded float, you may give your answer as a linear combination of powers of 2.

<table>
<thead>
<tr>
<th>sign</th>
<th>exponent</th>
<th>significand (also called mantissa or fraction)</th>
</tr>
</thead>
<tbody>
<tr>
<td>1 bit</td>
<td>6 bits</td>
<td>17 bits</td>
</tr>
</tbody>
</table>

a) Like the IEEE 754 standard, it uses a bias. What should it be such that it has the same effect as the bias for standard floats?

What is the largest non-infinite number possible?

What is the smallest positive integer that is not precisely representable with the new hybrid float?


b) This new hybrid standard uses unbiased rounding with one guard bit. Let: \( x = 0x360000 \) \( y = 0x320004 \) \( z = 0x368001 \). With the hybrid standard, \( x + y = z \). For what range of values of \( y \) (give hex of encoded floats) will produce the same result?

\[ 0x320004 \leq y \leq 0x350005 \]

0x320003 also correct to interpretation

How many valid float encodings are in that range?


c) Now you are designing a new floating point standard to meet range and precision requirements. We will do it in the the same spirit as the IEEE 754 standard, but with a different bit allocation. Bits are expensive so you want to use the minimum:

- Range from +/- \( 10^{150} \) to \( 10^{-150} \)
- Precise enough to distinguish \((1 + 1/(512 \text{ Mebi}))\) from \(1\)

What is the minimum number of bits for the exponent field?

What is the minimum number of bits for the significand field?

What is the minimum number of total bits?
**Question 4: I’m Bringing Hexy Back (14pts, 20min)**

We want to write a MIPS procedure called `xtoi`. It takes in a C string of a number written in hexadecimal in ASCII characters and it returns the number as an integer. You are guaranteed valid input that is only characters ‘0’ - ‘9’ and ‘a’ - ‘f’. The input string will be less than or equal to 8 characters. The header for an equivalent C function would look like:

```c
int xtoi(char* inpStr)
```

If `*inpStr = “0a4“`, running `xtoi(inpStr)` would return `164`. Fill in the following MIPS function to complete it. Be sure to obey MIPS register conventions. You are allowed to use pseudoinstructions, and you may need less lines than provided, but you can **not** add any extra lines in. (Hint: Your green sheet has an ASCII chart on the back)

```assembly
# a0 is pointer to start of string

xtoi:  addi $sp, $sp, -12
        sw $s0, 0($sp)
        sw $s1, 4($sp)
        sw $ra, 8($sp)
        add $s0, $a0, $0
        add $s1, $0, $0

loop:  lb $a0, 0($s0)
        beq $a0, $0, allDone
        addi $s0, $s0, 1
        jal digit
        sll $s1, $s1, 4
        add $s1, $s1, $v0
        j loop

allDone: add $v0, $s1, $0
         lw $s0, 0($sp)
         lw $s1, 4($sp)
         lw $ra, 8($sp)
         addi $sp, $sp, 12
         jr $ra

# subroutine that takes a hex digit in $a0 and returns value as a nibble

digit:  addi $v0, $a0, -48
        slti $t0, $a0, 10
        bne $t0, $0, dDone
        addi $v0, $a0, -87

digitDone: jr $ra
```
**Question 5: Play that funky music (15pts, 20min)**

Use the following MIPS code to answer the questions after it. You may assume the first instruction is at address 0x40000000.

```mips
j   funk

tasty:
  addiu $t1, $t1, 1
  sw   $t1, 0($t0)

funk:
  and   $s0, $s0, $s1
  la    $t0, funk
  lw    $t1, 0($t0)
  addi  $t2, $0, 0x3f
  and   $t2, $t2, $t1
  addi  $t3, $0, 0x27
  bne   $t2, $t3, tasty
```

a) Convert the 4th instruction, *(funk: and $s0, $s0, $s1)* into machine code. Express your answer in binary, as clearly as possible.

```
000000 10000 10001 10000 00000 100100
```

b) Given that `$s0 ← x` and `$s1 ← y`, with regard to the contents of `$s0` and `$s1`, what four lines of C have the same result as the above MIPS?

```
x = x&y;
x = x | y;
x = x ^ y;
x = ~(x | y);
```

c) Come up with a single MIPS instruction which has the same result as the provided code.

```
nor $s0, $s1, $0
```

d) Translate the final instruction *(bne $t2, $t3, tasty)* into machine code. Express your answer in hexadecimal.

```
0x154BFFF7
```