EECS 151/251A
Spring 2019
Digital Design and Integrated Circuits
Instructor:
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Lecture 3
Outline

- Hardware Description Language Overview
- Verilog Introduction
Hardware Description Languages
**Design Entry**

- Schematic entry/editing used to be the standard method in industry and universities.
- Used in EECS150 until 2002 and EE141 until recently

😊 Schematics are intuitive. They match our use of gate-level or block diagrams.

😊 Somewhat physical. They imply a physical implementation.

😄 Require a special tool (editor).

😢 Unless hierarchy is carefully designed, schematics can be confusing and difficult to follow on large designs.

- Hardware Description Languages (HDLs) are the new standard
  - except for PC board design, where schematics are still used.
**Hardware Description Languages**

- **Basic Idea:**
  - Language constructs describe circuits with two basic forms:
    - **Structural descriptions:** connections of components. Nearly one-to-one correspondence to schematic diagram.
    - **Behavioral descriptions:** use high-level constructs (similar to conventional programming) to describe the circuit function.

- Originally invented for simulation.
  - "logic synthesis" tools exist to automatically convert to gate level representation.
  - High-level constructs greatly improves designer productivity.
  - However, this may lead you to falsely believe that hardware design can be reduced to writing programs*

  ```
  "Structural" example:
  Decoder(output x0,x1,x2,x3; inputs a,b) {
    wire abar, bbar;
    inv(bbar, b);
    inv(abar, a);
    and(x0, abar, bbar);
    and(x1, abar, b);
    and(x2, a, bbar);
    and(x3, a, b);
  }
  ```

  ```
  "Behavioral" example:
  Decoder(output x0,x1,x2,x3; inputs a,b) {
    case [a b]
      00: [x0 x1 x2 x3] = 0x1;
      01: [x0 x1 x2 x3] = 0x2;
      10: [x0 x1 x2 x3] = 0x4;
      11: [x0 x1 x2 x3] = 0x8;
    endcase;
  }
  ```

  *Describing hardware with a language is similar, however, to writing a parallel program.

  **Warning:** this is a fake HDL!
Sample Design Methodology

Hierarchically defines structure and/or function of circuit.

HDL Specification

Simulation
Verification: Does the design behave as required with regards to function, timing, and power consumption?

Synthesis
Maps specification to resources of implementation platform (FPGA or ASIC).

Note: This is not the entire story. Other tools are useful for analyzing HDL specifications. More on this later.
Hardware Description Languages

Verilog:
- Simple C-like syntax for structural and behavior hardware constructs
- Mature set of commercial tools for synthesis and simulation
- Used in EECS 151 / 251A

VHDL:
- Semantically very close to Verilog
- More syntactic overhead
- Extensive type system for “synthesis time” checking

System Verilog:
- Enhances Verilog with strong typing along with other additions
- Somewhat less mature tool-flow

BlueSpec:
- Invented by Prof. Arvind at MIT
- Originally built within the Haskell programming language
- Now available commercially: bluespec.edu

Chisel:
- Developed at UC Berkeley
- Used in CS152, CS250
- Available at: chisel.eecs.berkeley.edu
Chisel: Constructing Hardware In a Scala Embedded Language

- Embeds hardware-description language in Scala, using Scala’s extension facilities: Hardware module is just data structure in Scala
- Different output routines generate different types of output (C, FPGA-Verilog, ASIC-Verilog) from same hardware representation
- Full power of Scala for writing hardware generators
  - Object-Oriented: Factory objects, traits, overloading, etc.
  - Functional: Higher-order functions, anonymous functions, currying
  - Compiles to JVM: Good performance, Java interoperability
Verilog: Brief History

- Invented as simulation language. Synthesis was an afterthought. Many of the basic techniques for synthesis were developed at Berkeley in the 80’s and applied commercially in the 90’s.
- Around the same time as the origin of Verilog, the US Department of Defense developed VHDL (A double acronym! VSIC (Very High-Speed Integrated Circuit) HDL). Because it was in the public domain it began to grow in popularity.
- Afraid of losing market share, Cadence opened Verilog to the public in 1990.
- An IEEE working group was established in 1993, and ratified IEEE Standard 1394 (Verilog) in 1995. We use IEEE Std 1364-2005.
- Verilog is the language of choice of Silicon Valley companies, initially because of high-quality tool support and its similarity to C-language syntax.
- VHDL is still popular within the government, in Europe and Japan, and some Universities.
- Most major CAD frameworks now support both.
Verilog Introduction
Verilog Introduction

- A **module** definition describes a component in a circuit
- Two ways to describe module contents:
  - **Structural Verilog**
    - List of sub-components and how they are connected
    - Just like schematics, but using text
    - Tedious to write, hard to decode
    - You get precise control over circuit details
    - May be necessary to map to special resources of the FPGA/ASIC
  - **Behavioral Verilog**
    - Describe what a component does, not how it does it
    - Synthesized into a circuit that has this behavior
    - Result is only as good as the tools
- Build up a hierarchy of modules. Top-level module is your entire design (or the environment to test your design).
Verilog Modules and Instantiation

- Modules define circuit components.
- Instantiation defines hierarchy of the design.

```verilog
module addr_cell (a, b, cin, s, cout);
  input     a, b, cin;
  output    s, cout;
endmodule
```

Note: A module is not a function in the C sense. There is no call and return mechanism. Think of it more like a hierarchical data structure.
module xor_gate (out, a, b);
input a, b;
output out;
wire aBar, bBar, t1, t2;
not invA (aBar, a);
not invB (bBar, b);
and and1 (t1, a, bBar);
and and2 (t2, b, aBar);
or or1 (out, t1, t2);
endmodule

Notes:
• The instantiated gates are not “executed”. They are active always.
• xor gate already exists as a built-in (so really no need to define it).
• Undeclared variables assumed to be wires. Don’t let this happen to you!
/* 2-input multiplexor in gates */
module mux2 (in0, in1, select, out);
    input in0,in1,select;
    output out;
    wire s0,w0,w1;
    not (s0, select);
    and (w0, s0, in0),
        (w1, select, in1);
    or  (out, w0, w1);
endmodule // mux2

Built-ins gates can have > 2 inputs. Ex:
and (w0, a, b, c, d);

Built-ins don’t need instance names

C++ style comments

Multiple instances can share the same “master” name.
module mux4 (in0, in1, in2, in3, select, out);
input in0, in1, in2, in3;
input [1:0] select;
output out;
wire w0, w1;

/* 2-input multiplexor in gates */
module mux2 (in0, in1, select, out);
    input in0, in1, select;
    output out;
    wire s0, w0, w1;
    not (s0, select);
    and (w0, s0, in0),
        (w1, select, in1);
    or  (out, w0, w1);
endmodule // mux2

module mux4 (in0, in1, in2, in3, select, out);
    input in0, in1, in2, in3;
    input [1:0] select;  // Signal array. Declares select[1], select[0]
    output out;
    wire w0, w1;

    mux2
    m0 (.select(select[0]), .in0(in0), .in1(in1), .out(w0)),
    m1 (.select(select[0]), .in0(in2), .in1(in3), .out(w1)),
    m3 (.select(select[1]), .in0(w0), .in1(w1), .out(out));
endmodule // mux4

Named ports. Highly recommended.
module foo (out, in1, in2);
  input       in1, in2;
  output      out;
assign out = in1 & in2;
endmodule

“continuous assignment”
Connects out to be the logical “and” of in1 and in2.

Shorthand for explicit instantiation of “and” gate (in this case).

The assignment continuously happens, therefore any change on the rhs is reflected in out immediately (except for the small delay associated with the implementation of the &).

Not like an assignment in C that takes place when the program counter gets to that place in the program.
module FullAdder(a, b, ci, r, co);
    input a, b, ci;
    output r, co;
    assign r = a ^ b ^ ci;
    assign co = a&ci | a&b | b&cin;
endmodule

module Adder(A, B, R);
    input [3:0] A;
    input [3:0] B;
    output [4:0] R;
    wire c1, c2, c3;
    FullAdder add0(.a(A[0]), .b(B[0]), .ci(1'b0), .co(c1), .r(R[0]),)
    add1(.a(A[1]), .b(B[1]), .ci(c1), .co(c2), .r(R[1]),)
    add2(.a(A[2]), .b(B[2]), .ci(c2), .co(c3), .r(R[2]),)
    add3(.a(A[3]), .b(B[3]), .ci(c3), .co(R[4]), .r(R[3]));
endmodule
## Verilog Operators

<table>
<thead>
<tr>
<th>Verilog Operator</th>
<th>Name</th>
<th>Functional Group</th>
</tr>
</thead>
<tbody>
<tr>
<td>()</td>
<td>bit-select or part-select</td>
<td></td>
</tr>
<tr>
<td>()</td>
<td>parenthesis</td>
<td></td>
</tr>
<tr>
<td>!</td>
<td>logical negation</td>
<td>Logical</td>
</tr>
<tr>
<td>~</td>
<td>negation</td>
<td>Bit-wise</td>
</tr>
<tr>
<td>&amp;</td>
<td>reduction AND</td>
<td>Reduction</td>
</tr>
<tr>
<td></td>
<td></td>
<td>reduction OR</td>
</tr>
<tr>
<td>~&amp;</td>
<td>reduction NAND</td>
<td>Reduction</td>
</tr>
<tr>
<td>^</td>
<td>reduction XOR</td>
<td>Reduction</td>
</tr>
<tr>
<td><del>^ or ^</del></td>
<td>reduction XNOR</td>
<td>Reduction</td>
</tr>
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<td>+</td>
<td>unary (sign) plus</td>
<td>Arithmetic</td>
</tr>
<tr>
<td>-</td>
<td>unary (sign) minus</td>
<td>Arithmetic</td>
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<td>multiply</td>
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<td>/</td>
<td>divide</td>
<td>Arithmetic</td>
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<td>+</td>
<td>binary plus</td>
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</tr>
<tr>
<td>-</td>
<td>binary minus</td>
<td>Arithmetic</td>
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<td>Shift</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>shift right</td>
<td>Shift</td>
</tr>
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<td>greater than</td>
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<td>greater than or equal to</td>
<td>Relational</td>
</tr>
<tr>
<td>&lt;</td>
<td>less than</td>
<td>Relational</td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal to</td>
<td>Relational</td>
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</tr>
<tr>
<td>!==</td>
<td>case inequality</td>
<td>Equality</td>
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<td>bit-wise AND</td>
<td>Bit-wise</td>
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<td>^=</td>
<td>bit-wise XOR</td>
<td>Bit-wise</td>
</tr>
<tr>
<td>~^ or ~^</td>
<td>bit-wise XNOR</td>
<td>Bit-wise</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bit-wise OR</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>logical AND</td>
<td>Logical</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>?;</td>
<td>conditional</td>
<td>Conditional</td>
</tr>
</tbody>
</table>
**Verilog Numbers**

**Constants:**

- 14  ordinary decimal number
- –14 2’s complement representation
- 12'b0000_0100_0110  binary number (“_” is ignored)
- 12'h046  hexadecimal number with 12 bits

**Signal Values:**

By default, Values are unsigned

\[
\text{e.g., } C[4:0] = A[3:0] + B[3:0];
\]

if \(A = 0110\) (6) and \(B = 1010\) (-6)

\[
C = 10000 \text{ not } 00000
\]

i.e., \(B\) is zero-padded, not sign-extended

```verilog
wire signed [31:0] x;
```

Declares a signed (2’s complement) signal array.
Verilog Assignment Types
Continuous Assignment Examples

```verilog
assign R = X | (Y & ~Z);
assign r = &X;
assign R = (a == 1'b0) ? X : Y;
assign P = 8'hff;
assign P = X * Y;
assign P[7:0] = {4{X[3]}, X[3:0]};
assign {cout, R} = X + Y + cin;
assign Y = A << 2;
assign Y = {A[1], A[0], 1'b0, 1'b0};
```

- **Example Reduction Operator**
- **Use of Bit-wise Boolean Operators**
- **Conditional Operator**
- **Example Constants**
- **Arithmetic Operators (use with care!)**
- **(ex: sign-extension)**
- **Bit Field Concatenation**
- **Bit Shift Operator**
- **Equivalent Bit Shift**
Non-continuous Assignments

A bit strange from a hardware specification point of view. Shows off Verilog roots as a simulation language.

“always” block example:

```verilog
module and_or_gate (out, in1, in2, in3);
    input    in1, in2, in3;
    output   out;
    reg      out;
always @(in1 or in2 or in3) begin
    out = (in1 & in2) | in3;
end
endmodule
```

“reg” type declaration. Not really a register in this case. Just a Verilog idiosyncrasy.

“sensitivity” list, triggers the action in the body.

 Isn’t this just: `assign out = (in1 & in2) | in3;`?

Why bother?
Always Blocks

Always blocks give us some constructs that are impossible or awkward in continuous assignments.

case statement example:

module mux4 (in0, in1, in2, in3, select, out);
    input in0, in1, in2, in3;
    input [1:0] select;
    output      out;
    reg         out;

always @ (in0 in1 in2 in3 select)
    case (select)
        2'b00: out=in0;
        2'b01: out=in1;
        2'b10: out=in2;
        2'b11: out=in3;
    endcase
endmodule // mux4

Could'nt we just do this with nested "if"s?

Well yes and no!
Always Blocks

Nested if-else example:

```verilog
module mux4 (in0, in1, in2, in3, select, out);
    input in0, in1, in2, in3;
    input [1:0] select;
    output      out;
    reg         out;

    always @ (in0 in1 in2 in3 select)
        if (select == 2'b00) out=in0;
        else if (select == 2'b01) out=in1;
        else if (select == 2'b10) out=in2;
        else out=in3;
endmodule // mux4
```

Nested if structure leads to “priority logic” structure, with different delays for different inputs (in3 to out delay > than in0 to out delay). Case version treats all inputs the same.
module FullAdder(a, b, ci, r, co);
    input a, b, ci;
    output r, co;

    assign r = a ^ b ^ ci;
    assign co = a&ci + a&b + b&cin;
endmodule

module Adder(A, B, R);
    input [3:0] A;
    input [3:0] B;
    output [4:0] R;

    wire c1, c2, c3;
    FullAdder
    add0 (.a(A[0]), .b(B[0]), .ci(1'b0), .co(c1), .r(R[0]) ),
    add1 (.a(A[1]), .b(B[1]), .ci(c1), .co(c2), .r(R[1]) ),
    add2 (.a(A[2]), .b(B[2]), .ci(c2), .co(c3), .r(R[2]) ),
    add3 (.a(A[3]), .b(B[3]), .ci(c3), .co(R[4]), .r(R[3]) );
endmodule
Example - Ripple Adder Generator

Parameters give us a way to generalize our designs. A module becomes a “generator” for different variations. Enables design/module reuse. Can simplify testing.

```
module Adder(A, B, R);
    parameter N = 4;
    input [N-1:0] A;
    input [N-1:0] B;
    output [N:0] R;
    wire [N:0] C;

    genvar i;
    generate
        for (i=0; i<N; i=i+1) begin:bit
            FullAdder add(.a(A[i], .b(B[i]), .ci(C[i]), .co(C[i+1]), .r(R[i]));
        end
    endgenerate

    assign C[0] = 1'b0;
    assign R[N] = C[N];
endmodule
```

Declare a parameter with default value.

Note: this is not a port. Acts like a “synthesis-time” constant.

Replace all occurrences of “4” with “N”.

variable exists only in the specification - not in the final circuit.

Keyword that denotes synthesis-time operations

For-loop creates instances (with unique names)

Adder adder4 ( ... );
Overwrite parameter N at instantiation.

Adder #(N(64))
adder64 ( ... );
More on Generate Loop

Permits variable declarations, modules, user defined primitives, gate primitives, continuous assignments, initial blocks and always blocks to be instantiated multiple times using a for-loop.

// Gray-code to binary-code converter
module gray2bin1 (bin, gray);
    parameter SIZE = 8;
    output [SIZE-1:0] bin;
    input  [SIZE-1:0] gray;

genvar i;

generate
    for (i=0; i<SIZE; i=i+1)
    begin:
        assign bin[i] = ^gray[SIZE-1:i];
    end
endgenerate
endmodule

variable exists only in the specification - not in the final circuit.

Keywords that denotes synthesis-time operations

For-loop creates instances of assignments
Loop must have constant bounds

generate if-else-if based on an expression that is deterministic at the time the design is synthesized.

generate case: selecting case expression must be deterministic at the time the design is synthesized.
Modularity is essential to the success of large designs.
High-level primitives enable direct synthesis of behavioral descriptions (functions such as additions, subtractions, shifts (<< and >>), etc.)

**Example: A 32-bit ALU**

**Function Table**

<table>
<thead>
<tr>
<th>F2</th>
<th>F1</th>
<th>F0</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0</td>
<td>0</td>
<td>A + B</td>
</tr>
<tr>
<td>0</td>
<td>0</td>
<td>1</td>
<td>A + 1</td>
</tr>
<tr>
<td>0</td>
<td>1</td>
<td>0</td>
<td>A - B</td>
</tr>
<tr>
<td>0</td>
<td>1</td>
<td>1</td>
<td>A - 1</td>
</tr>
<tr>
<td>1</td>
<td>0</td>
<td>X</td>
<td>A * B</td>
</tr>
</tbody>
</table>
Module Definitions

2-to-1 MUX
module mux32two(i0, i1, sel, out);
  input [31:0] i0, i1;
  input sel;
  output [31:0] out;

  assign out = sel ? i1 : i0;
endmodule

3-to-1 MUX
module mux32three(i0, i1, i2, sel, out);
  input [31:0] i0, i1, i2;
  input [1:0] sel;
  output [31:0] out;
  reg [31:0] out;

  always @ (i0 or i1 or i2 or sel)
  begin
    case (sel)
      2'b00: out = i0;
      2'b01: out = i1;
      2'b10: out = i2;
      default: out = 32'bx;
    endcase
  end
endmodule

32-bit Adder
module add32(i0, i1, sum);
  input [31:0] i0, i1;
  output [31:0] sum;

  assign sum = i0 + i1;
endmodule

32-bit Subtracter
module sub32(i0, i1, diff);
  input [31:0] i0, i1;
  output [31:0] diff;

  assign diff = i0 - i1;
endmodule

16-bit Multiplier
module mul16(i0, i1, prod);
  input [15:0] i0, i1;
  output [31:0] prod;

  // this is a magnitude multiplier
  // signed arithmetic later
  assign prod = i0 * i1;
endmodule
Top-Level ALU Declaration

Given submodules:

- module mux32two(i0, i1, sel, out);
- module mux32three(i0, i1, i2, sel, out);
- module add32(i0, i1, sum);
- module sub32(i0, i1, diff);
- module mul16(i0, i1, prod);

Declaration of the ALU Module:

module alu(a, b, f, r);
  input [31:0] a, b;
  input [2:0] f;
  output [31:0] r;
  wire [31:0] addmux_out, submux_out;
  wire [31:0] add_out, sub_out, mul_out;
  mux32two adder_mux(.io(b), .i1(32'd1), .sel(f[0]), .out(addmux_out));
  mux32two sub_mux(.io(b), .i1(32'd1), .sel(f[0]), .out(submux_out));
  add32 our_adder(.i0(a), .i1(addmux_out), .sum(add_out));
  sub32 our_subtractor(.i0(a), .i1(submux_out), .diff(sub_out));
  mul16 our_multiplier(.i0(a[15:0]), .i1(b[15:0]), .prod(mul_out));
  mux32three output_mux(.i0(add_out), .i1(sub_out), .i2(mul_out), .sel(f[2:1]), .out(r));
endmodule
Top-Level ALU Declaration, take 2

- No Hierarchy:
- Declaration of the ALU Module:

```verilog
module alu(a, b, f, r);
    input [31:0] a, b;
    input [2:0] f;
    output [31:0] r;
    always @ (a or b or f)
        case (f)
            3'b000:  r = a + b;
            3'b001:  r = a + 1'b1;
            3'b010:  r = a - b;
            3'b011:  r = a - 1'b1;
            3'b100:  r = a * b;
            default: r = 32'bx;
        endcase
endmodule
```

Will this synthesize into 2 adders and 2 subtractors or 1 of each?
We use behavioral modeling at the bottom of the hierarchy
Use instantiation to 1) build hierarchy and, 2) map to FPGA and ASIC resources not supported by synthesis.
Favor continuous assign and avoid always blocks unless:
- no other alternative: ex: state elements, case
- helps readability and clarity of code: ex: large nested if else
Use named ports.
Verilog is a big language. This is only an introduction.
- Complete IEEE Verilog-Standard document (1364-2005) linked to class website.
- Harris & Harris book chapter 4 is a good source.
- Be careful of what you read on the web. Many bad examples out there.
- We will be introducing more useful constructs throughout the semester. Stay tuned!
Final thoughts on Verilog Examples

Verilog looks like C, but it describes hardware:
Entirely different semantics: multiple physical elements with parallel activities and temporal relationships.

A large part of digital design is knowing how to write Verilog that gets you the desired circuit. First understand the circuit you want then figure out how to code it in Verilog. If you try to write Verilog without a clear idea of the desired circuit, you will struggle.

As you get more practice, you will know how to best write Verilog for a desired result.

Be suspicious of the synthesis tools! Check the output of the tools to make sure you get what you want.