Privilege Separation

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Part III OS Security

- Principles of OS Security
- State-of-the-art techniques & tools for OS Security

The Story Continues...

- Programs have vulnerabilities and may not find all of them ahead of time
- What can we do?
 - Build security mechanisms to minimize damage
- Examples
 - Priviledge separation to prevent priviledge escalation
 Isolation to protect other parts of the program and
 - other programs
 - Sandboxing to limit the damage it does to the system
 General concept: reference monitor

Operating System & Privilege

• OS's role

- Interface between hardware & applications
- Manages resources
- Provide protection to hardware & applications
- Privilege
 - Rights to perform certain operations
 - » E.g., writes to certain files & certain network operations

Principle of Least Privilege

• Give the user/program only the privilege it needs to get its task done

One of the most important principles in systems security

• Why?

- Limit the damage when program misbehaves or is compromised

- · What privileges should you give to your
 - -ssh server
 - Video game program

Mangement of Privileges

- Example: File privileges
- · Superuser/root mode - Access to everything
- · Windows privilege model
 - Previously, most programs require superuser mode to install/run
 - » Consequence: most users log on as administrator
 - Vista: User Account Control (UAC)
 - When user log on as a standard user, a token with basic privileges is assigned

 - When user log on as an administrator, two tokens are assigned
 One with basic privileges, the other with root privileges
 Normal programs will be started with basic privileges
 Programs require root privileges will be prompted for user permission

Privileged Programs

- Privilege management is coarse-grained in today's OS
 - Root can do anything
- Many programs run as root
 - Even though they only need to perform a small number of priviledged operations
- What's the problem?
 - Privileged programs are juicy targets for attackers
 - By finding a bug in parts of the program that do not need privilege, attacker can gain root

What Can We Do?

- Drop privilege as soon as possible
- Ex: a network daemon only needs privilege to bind to low port # (<1024) at the beginning

 Solution?
 - Drop privilege right after binding the port
- What benefit do we gain?
 - Even if attacker finds a bug in later part of the code, can't gain privilege any more
- · How to drop privilege?
 - Setuid programming in UNIX

Unix file permission

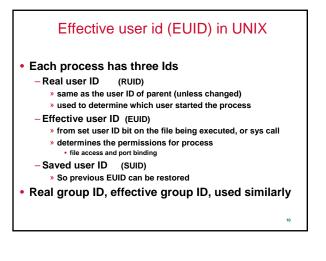
- Each file has owner and group
- Permissions set by owner
 - Read, write, execute
 - Owner, group, other
 - Represented by vector of four octal values



setid

- Only owner, root can change permissions

 This privilege cannot be delegated or shared
- Setid bits
 - Talk about this in a sec



Operations on UIDs

Root

- ID=0 for superuser root; can access any file

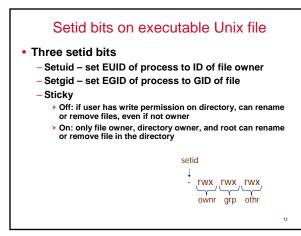
Fork and Exec

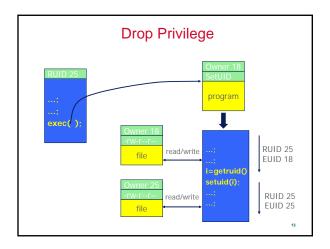
- Inherit three IDs, except exec of file with setuid bit

Setuid system calls

seteuid(newid) can set EUID to
 » Real ID or saved ID, regardless of current EUID
 » Any ID, if EUID=0

Details are actually more complicated
 – Several different calls: setuid, seteuid, setreuid







Administravia

- Photo submission due Oct 29
- Project 2

What Happens if you can't drop privilege?

- In what example scenarios does this happen?
 - A service loop
 - E.g., ssh
- Solution?
 - Privilege separation
 - Identifying operations that need privileges
 - Separate original code into master (priviledged) and slave (unprivileged)
- Example: ssh

Privilege Separation

• Process:

- Step 1: Identify which operations require privilege
- Step 2: rewrite programs into 2 or more parts
- Approach:
 - Manual
 - » Have been done on security-critical programs, e.g., ssh
 - » Labor-intensive and may miss privileged operations

- Automatic

- » Automatic inference of privileged operations using a few initial annotations
- » Automatic source-to-source rewriting Privileged code move into master
 - Unprivileged code move into slave
 - · Stubs for inter communication

Automatic Privilege Separation

- · Step 1: automatic inference of privileged data and operations
 - Given some initial annotations of privileged data and/or operations, infer what other data/operations are privileged
 - Idea: can be viewed as a form of static taint analysis - Approach:
 - - » Define qualifier _priv_ and _unpriv_ » Operations on _priv_ results in _priv_

priv b _priv_ c

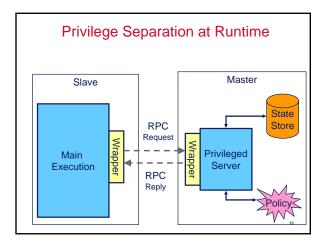
priv g

int _priv_ a; Int _priv_ f(); int b = f(a); c= c+b; g(c);

Automatic Privilege Separation

Step 2: automatic source-to-source transformation

- Move privileged data and code to Master
- For call to privileged functions, change the call site to a wrapper function which marshals the args on slave side and sends them to Master's stub
- Similar stubs on returns for unprivileged return values





Summary: Privilege Separation

- Only master is privileged, usually much smaller
- Slave is unprivileged
- Bug in slave cannot harm master, cannot gain privilege
- How to protect master from a compromised slave?
 - Fault isolation: e.g., running in different processes