

## Random Number Generation and Electronic Cash

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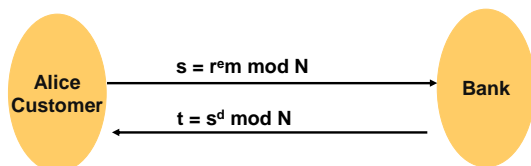
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## Building Block: Blind Signatures

- **Blind signature: achieve anonymity**
  - How can Alice get a signature from the Bank without the Bank knowing what message is being signed?
- **Protocol:**
  - generating blind signature on message  $m$  in RSA setting
  - Bank's private key  $(d, p, q)$ , public key  $(e, N)$



- Alice computes  $t/r \bmod N = m^d \bmod N$

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## Ecash Using Blind Signature

- How to use blind signature to build ecash?
- A valid \$1 bill is a pair  $(x, y)$ , where  $y = \text{hash}(x)^d \bmod N$ ,  $\text{hash}()$  is one-way function
- How does the ecash protocol work?
- Why do we need  $\text{hash}()$ ?
- How to prevent double spending?
- What to do for different denominations?
  - Nickles, dimes, dollars

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### Other Methods for Ecash

- **Use zero-knowledge proofs (out of scope)**
  - More building blocks of ZKP
  - Support many properties
    - » Identifying double spenders

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### Administrative Matters

- **Next class: talk about midterm scope**

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### Untrusted Storage

- **User's sensitive data often stored in untrusted storage**
  - Third party storage/out-sourced storage
  - Mis-configuration causes information leakage
  - Attacker hacks into system
  - Insider attack
- **Need to encrypt data to protect privacy**
- **Yet, need to perform certain operations on data for functionality**

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## Operations on Encrypted Data

- What kind of operations needed on encrypted data?
- What are your favorite applications that you would like to enable on encrypted data?
- Here we focus on searching on encrypted data

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## Motivating Example: Equality Search on Encrypted Data

- Searching encrypted e-mails on servers
- Searching encrypted files on servers
- Searching in encrypted databases



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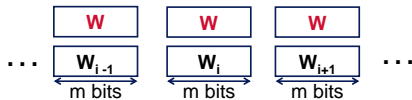
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## Sequential Scan

Search for  $W$



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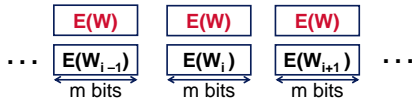
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## Sequential Scan – with encryption

Search for  $W$



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## Desired Properties

- **Word search is provably secure**
  - Provable encryption properties
  - Server cannot search for arbitrary words
  - Does not leak information about other words
  - Does not reveal query word
- **Efficiency**
  - Low computation overhead
  - Low space and communication overhead
  - Low management overhead

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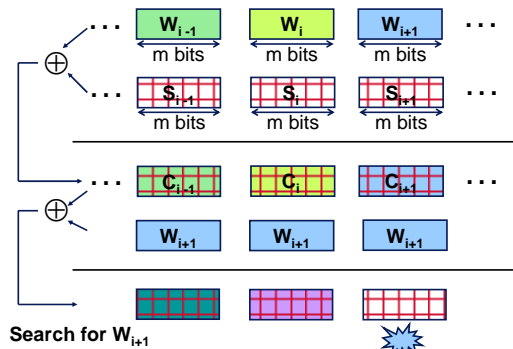
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## The Key Idea



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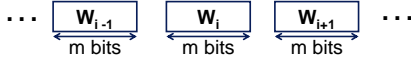
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### Setup and Notations

- Document: sequence of fixed length words



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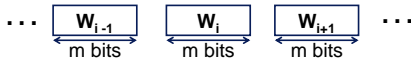
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### Setup and Notations

- Document: sequence of fixed length words



- $L_0, L_1, L_2, \dots$   
–sequence of pseudorandom  $n$ -bit blocks

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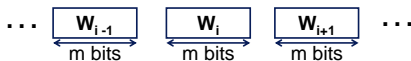
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### Setup and Notations

- Document: sequence of fixed length words



- $L_0, L_1, L_2, \dots$   
–sequence of pseudorandom  $n$ -bit blocks
- Pseudorandom Function  $F_K$   
– maps  $n$  bits to  $(m-n)$  bits

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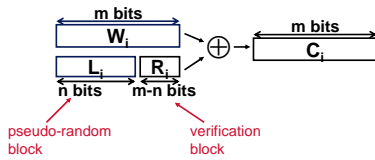
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### Basic Scheme (Encryption)



We xor the word  $W_i$  with

$$L_i \leftarrow \text{pseudorandom bits}$$

$$R_i \leftarrow F_K(L_i)$$

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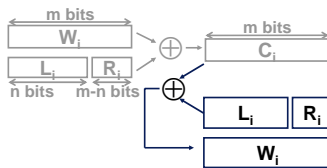
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### Basic Scheme (Decryption)



$$L_i \leftarrow \text{pseudorandom bits}$$

$$R_i \leftarrow F_K(L_i)$$

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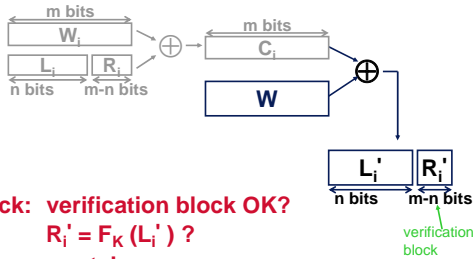
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### Basic Scheme (Searches)

Search for word  $W$ , give server  $W$  and  $K$



Check: verification block OK?

$$R'_i = F_K(L'_i) ?$$

Yes  $\Rightarrow$  match,

$$\text{false positive rate} = 1 / 2^{m-n}$$

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## Controlled Searches and Query Isolation

- To keep server from searching for arbitrary words & To avoid leaking information about other words
- In encryption:
  - Replace
  - $R_i \leftarrow F_K(L_i)$
  - with
  - $R_i \leftarrow F_{K_i}(L_i)$ , where  $K_i = F'_K(W_i)$
- To search for word  $W$ :
  - Reveal
  - $K_w = F'_K(W)$
- Enhancements:
  - Check only for "word occurs at least once" in document
  - Check only for "word occurs at least N times" in document

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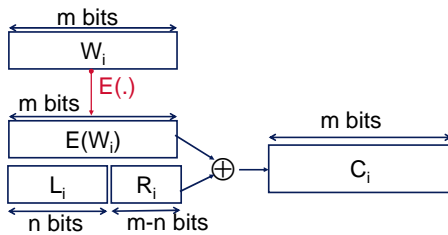
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## Hidden Queries



$L_i \leftarrow$  pseudorandom bits

$R_i \leftarrow F_{K_i}(L_i)$

where  $K_i = F'_K(E(W_i))$

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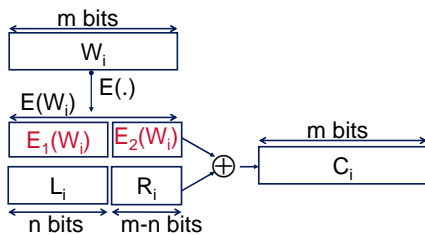
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## Final Scheme (Encryption)



$L_i \leftarrow$  pseudorandom bits

$R_i \leftarrow F_{K_i}(L_i)$

where  $K_i = F'_K(E_1(W_i))$

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Summary for Keyword Search on Encrypted Data  
(Symmetric Key Case)

- **Provable security**
  - Provable secrecy
  - Controlled search
  - Query isolation
  - Hidden queries
- **Simple and efficient**
  - $O(\text{length of document})$  stream cipher, block cipher and MAC operations for encryption/decryption
  - $O(\text{length of document})$  MAC operations for search
  - Almost no space and communication overhead
  - Easy to add documents
  - Convenient key management :  
user needs only one master key

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Conclusion

- Ecash
- Search/computation on encrypted data

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