Proj 2 announcements Digital certificates
CA venifies Google representative
That they are indeed from
own google. Come google. com has Pkg certificate for Google = Certificate authority Advartages: - r can obtain certificate from anyone, don't need. Disadvantages: -> if Sk is compromised, certificate still valid until experior

> mitigation: revocation lists: blacklisted

-Scalability: mitigated through hierarchy of curif

certificates; browser downloads periodically

Hierarchy PKI CA: Jerry VCB President Stanford president (2) PKCA hardcoded in browners President has PKp", egung) Ratuca has PKR", equity) Penguin has PKPE; expire Certificate

Certificate chain gets verified by browser - browser ducks all signatures verify again PK of pro in previous certificate and last PK is PK of CA Scalable because no one entity issues certificates for everyone Compromise of SK of entity N in hierarchy con head to impersonation of all entities underneda N, but no other entities be someone they are

Securing Internet Communication: TLS

CS 161: Computer Security
Prof. Raluca Ada Popa

Feb 22, 2018

Today's Lecture

- Applying crypto technology in practice
- Two simple abstractions cover many use cases for crypto:
 - "Sealed blob": Data that is encrypted and authenticated under a particular key
 - Secure channel: Communication channel that can't be eavesdropped on or tampered with
- Today: SSL a secure channel

Today's Lecture

- Goal #1: overview of SSL/TLS, the most prominent Internet security protocol
 - Secures the web via HTTPS
- Goal #2: cement understanding of crypto building blocks & how they're used together

Building Secure End-to-End Channels

- End-to-end = communication protections achieved all the way from originating client to intended server, between endpoints
 - With no need to trust intermediaries
- Dealing with threats:
 - Eavesdropping?
 - Encryption (including session keys)
 - Manipulation (injection, MITM)?
 - Integrity (use of a MAC); replay protection
 - Impersonation? (someone pretending as you)
 - Signatures

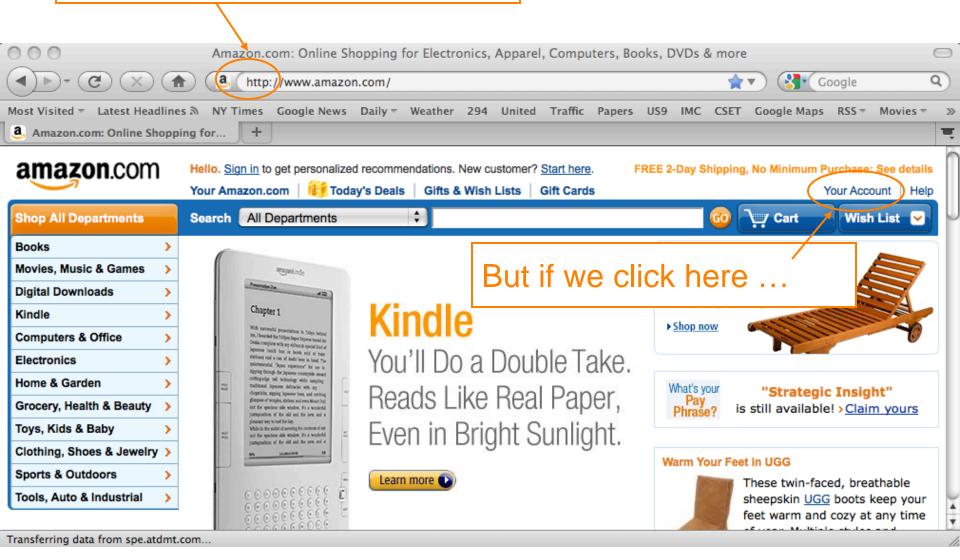
```
(Q: What's missing?)

A: Availability ...
```

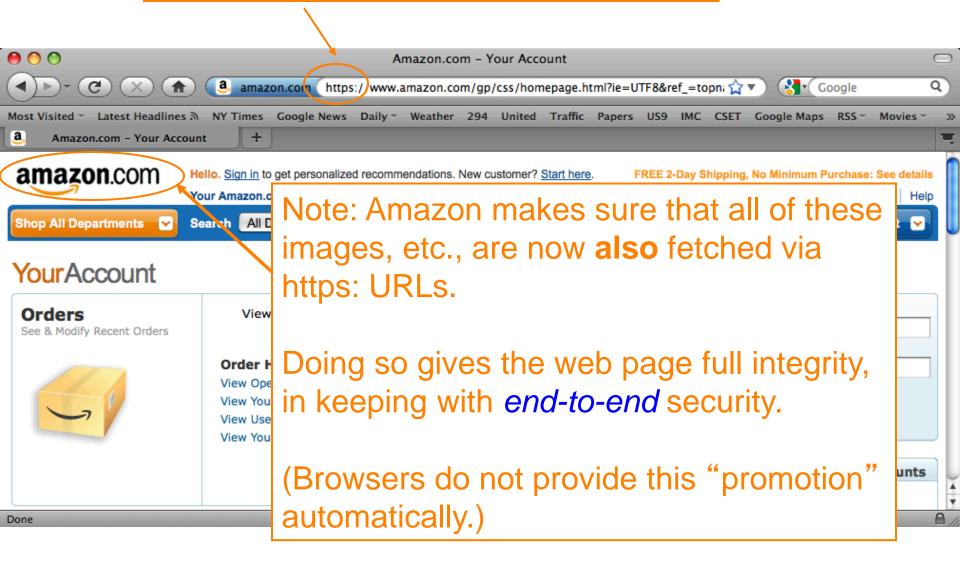
Building A Secure End-to-End Channel: SSL/TLS

- SSL = Secure Sockets Layer (predecessor)
- TLS = Transport Layer Security (standard)
 - Both terms used interchangeably
- Security for any application that uses TCP
 - Secure = encryption/confidentiality + integrity + authentication (of server, but *not* of client)
 - E.g., puts the 's' in "https"

Regular web surfing - http: URL



Web surfing with TLS/SSL - https: URL



RSA Encryption

- A public-key encryption algorithm, not only digital signature
- The encrypt algorithm is similar to the verify algorithm, and the decrypt similar to the sign algorithm
- Small differences: encrypt the message with special padding, instead of signing a hash of the message

HTTPS Connection (SSL / TLS)

- Suppose a browser (client) wants to connect to a server who has a certificate from a trusted CA
- Client browser and server will exchange symmetric keys using SSL/TLS
- Then, they will send encrypted & authenticated traffic to each other

HTTPS Connection (SSL / TLS)

Browser

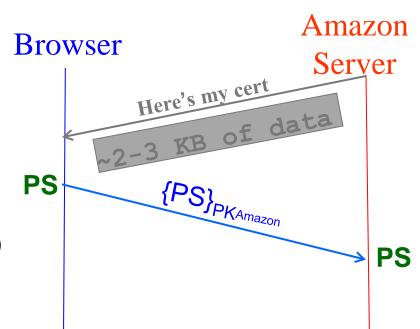
- Browser (client) connects to Amazon's HTTPS server
- Client picks 256-bit random number R_B, sends over list of crypto algorithms it supports
- Server picks 256-bit random number R_s, selects algorithms to use for this session
- Server sends over its certificate
- (all of this is in the clear)
- Client now validates cert

Hello. $My rnd #= R_B$. I support(TLS+RSA+AES128+SHA256) or (SSL+RSA+3DES+MD5) or My rnd $\#=R_S$. Let's use TLS+RSA+AES128+SHA256 Here's my cert ·2-3 KB of data

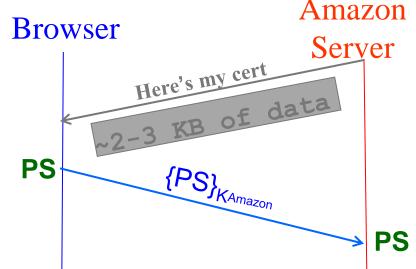
Amazon

Server

- For RSA, browser constructs "Premaster Secret" PS
- Browser sends PS encrypted using Amazon's public RSA key PK_{Amazon}
- Using PS, R_B, and R_S, browser & server derive symm. *cipher keys* (C_B, C_S) & MAC *integrity keys* (I_B, I_S)
 One pair to use in each direction



- For RSA, browser constructs "Premaster Secret" PS
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These <u>seed</u> a cryptographically strong pseudo-random number generator (PRNG). Q: why R_B and R_S, and not just a longer R_S?

A: just in sees one party's randomness is not

A: just in case one party's randomness is not good

- For RSA, browser constructs "Premaster Secret" PS
- Browser sends PS encrypted using Amazon's public RSA key K_{Amazon}
- Using PS, R_B, and R_S, browser & server derive symm. *cipher keys* (C_B, C_S) & MAC *integrity keys* (I_B, I_S)
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- Browser & server exchange MACs computed over entire dialog so far

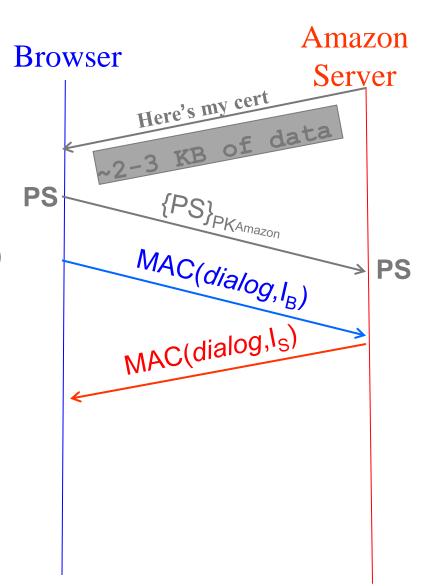
Amazon Browser Server Here's my cert PS MAC(dialog, I_B) MAC(dialog, ls)

Q: Why?

A: So they know they have same (C_B, C_S), (I_B, I_S)

- For RSA, browser constructs "Premaster Secret" PS
- Browser sends PS encrypted using Amazon's public RSA key K_{Amazon}
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- If good MAC, browser displays

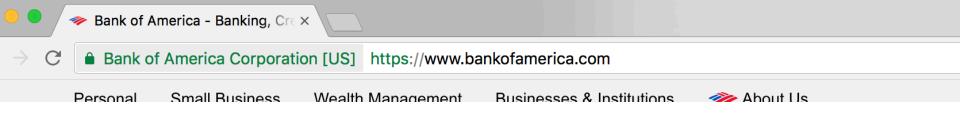




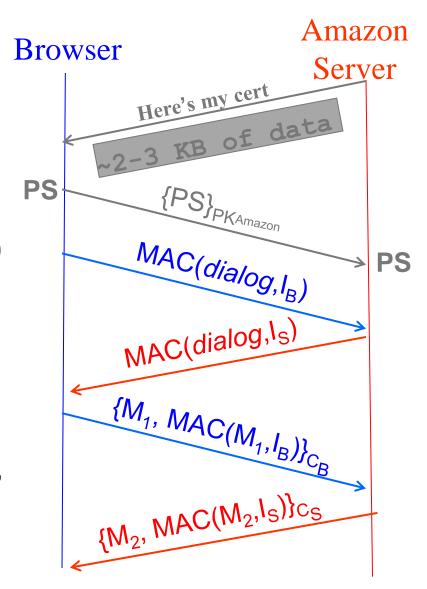
On Firefox:



On Chrome:



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- Browser & server exchange MACs computed over entire dialog so far
- If good MAC, Browser displays
- All subsequent communication encrypted w/ symmetric cipher (e.g., AES128) cipher keys in some chaining mode, MACs
 - Sequence #'s included to thwart replay attacks

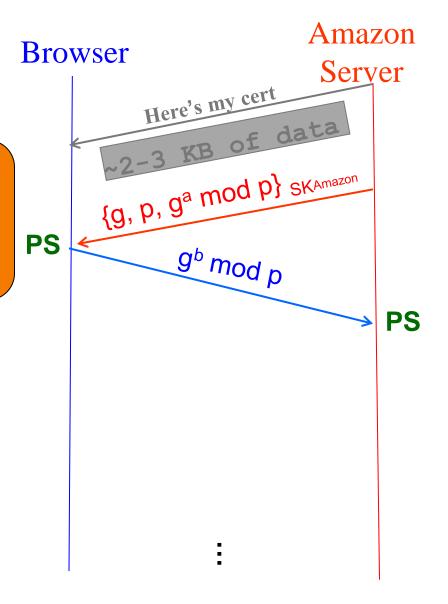


Alternative: Key Exchange via Diffie-Hellman

 For Diffie-Hellman, server generates random a, sends public params and g^a mod p

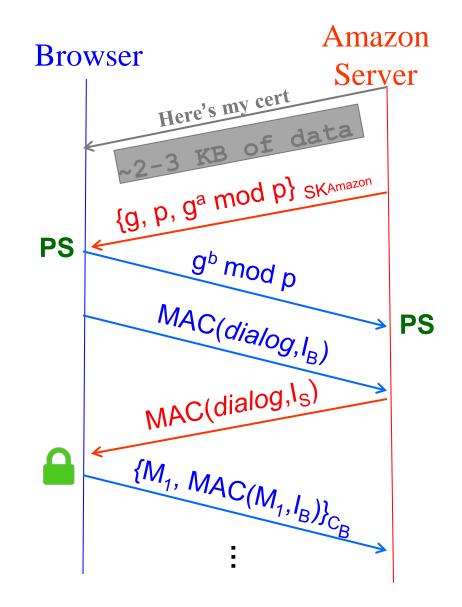
Q: How can we prevent MITM?

A: Server signs g^a mod p using SK_{Amazon}, verity using PK_{Amazon} from server certificate



Alternative: Key Exchange via Diffie-Hellman

- For Diffie-Hellman, server generates random a, sends public params and g^a mod p
 - Signed with server's private key
- Browser verifies signature using PK from certificate
- Browser generates random b, computes PS = g^{ab} mod p, sends to server
- Server also computes
 PS = g^{ab} mod p
- Remainder is as before: from PS, R_B, and R_S, browser & server derive symm. *cipher keys* (C_B, C_S) and MAC *integrity keys* (I_B, I_S), etc...



RSA versus Diffie-Hellman

 Forward secrecy: If attacker steals long term secret key of server, SK_{Amazon}, should not be able to read past conversations (cannot compromise past session keys (C_B, C_S) & (I_B, I_S))

- Why matters?
 - Attackers log traffic now. Compromise key in future and try to decrypt the traffic.

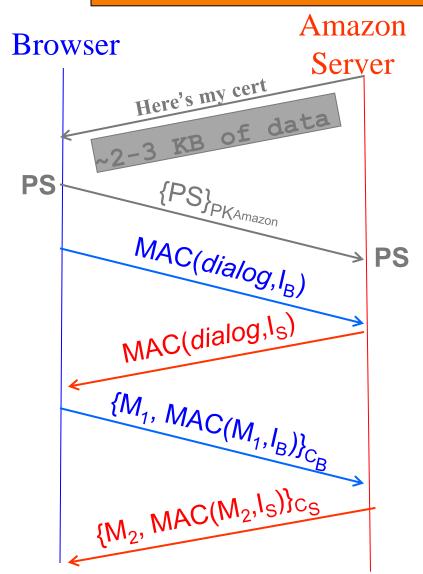
Exchange with RSA

Q: Forward secrecy?
A: No forward secrecy because attacker can decrypt PS and knows R_B, and R_S and computes secrets

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- If good MAC, Browser displays



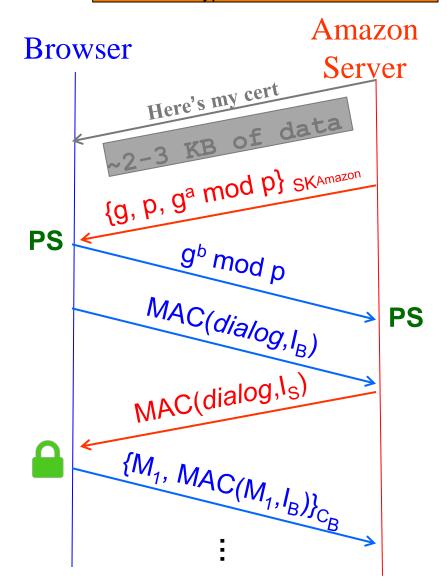
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 - Sequence #'s thwart replay attacks



Exchange via Diffie-Hellman

Q: Forward secrecy?
A: Has forward secrecy because shared secret never sent over the network! If attacker as SK_{Amazon} , cannot decrypt a.

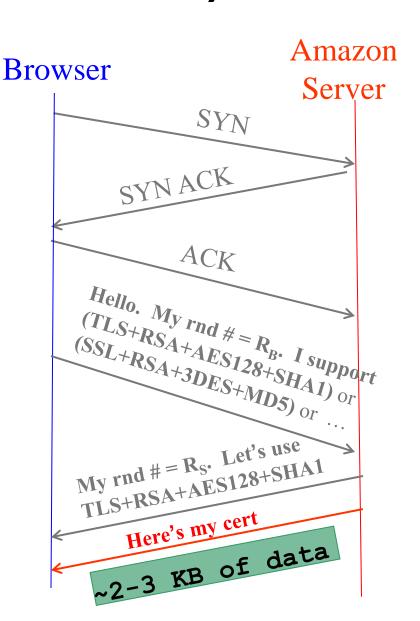
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- Server sends over its certificate
- (all of this is in the clear)
- Client now validates cert



Certificates

- Browser compares domain name in cert w/ URL
 - Note: this provides an end-to-end property
 (as opposed to say a cert associated with an IP address)
- Browser accesses <u>separate</u> cert belonging to **issuer**
 - These are hardwired into the browser and trusted!
 - -There could be a *chain* of these ...
- Browser applies issuer's public key to verify signature S, obtaining hash of what issuer signed
 - Compares with its own SHA-256 hash of Amazon's cert
- Assuming hashes match, now have high confidence it's indeed Amazon ...
 - assuming signatory is trustworthy

= assuming didn't lose private key; assuming didn't sign thoughtlessly

End-to-End ⇒ **Powerful Protections**

- Attacker runs a sniffer to capture our WiFi session?
 - (maybe by breaking crummy WEP security)
 - But: encrypted communication is unreadable
 - No problem!
- DNS cache poisoning gives client wrong IP address
 - Client goes to wrong server
 - But: detects impersonation
 - No problem!
- Attacker hijacks our connection, injects new traffic
 - But: data receiver rejects it due to failed integrity check
 - No problem!

Powerful Protections, cont.

- Attacker manipulates routing to run us by an eavesdropper or take us to the wrong server?
 - But: they can't read; we detect impersonation
 - No problem!
- Attacker slips in as a Man In The Middle?
 - But: they can't read, they can't inject
 - They can't even replay previous encrypted traffic
 - No problem!

Validating Amazon's Identity, cont.

- Browser retrieves cert belonging to the issuer
 - -These are hardwired into the browser and trusted!
- What if browser can't find a cert for the issuer?



This Connection is Untrusted

You have asked Firefox to connect securely to www.mikestoolbox.org, but we can't confirm that your connection is secure.

Normally, when you try to connect securely, sites will present trusted identification to prove that you are going to the right place. However, this site's identity can't be verified.

What Should I Do?

If you usually connect to this site without problems, this error could mean that someone is trying to impersonate the site, and you shouldn't continue.

Get me out of here!

Technical Details

www.mikestoolbox.org uses an invalid security certificate.

The certificate is not trusted because the issuer certificate is not trusted.

(Error code: sec_error_untrusted_issuer)

I Understand the Risks



Verify Certificate

Safari can't verify the identity of the website "www.mikestoolbox.org".

The certificate for this website was signed by an unknown certifying authority. You might be connecting to a website that is pretending to be "www.mikestoolbox.org", which could put your confidential information at risk. Would you like to connect to the website anyway?

?

Show Certificate

Cancel

Continue

Validating Amazon's Identity, cont.

- Browser retrieves cert belonging to the issuer
 - -These are hardwired into the browser and trusted!
- What if browser can't find a cert for the issuer?
- If it can't find the cert, then warns the user that site has not been verified
 - Can still proceed, just without authentication
- Q: Which end-to-end security properties do we lose if we incorrectly trust that the site is whom we think?
- A: All of them!
 - Goodbye confidentiality, integrity, authentication
 - Man in the middle attacker can read everything, modify, impersonate

SSL / TLS Limitations

- Properly used, SSL / TLS provides powerful endto-end protections
- So why not use it for everything??
- Issues:
 - Cost of public-key crypto (fairly minor)
 - o Takes non-trivial CPU processing (but today a minor issue)
 - o Note: symmetric key crypto on modern hardware is non-issue
 - Hassle of buying/maintaining certs (fairly minor)

SSL / TLS Limitations

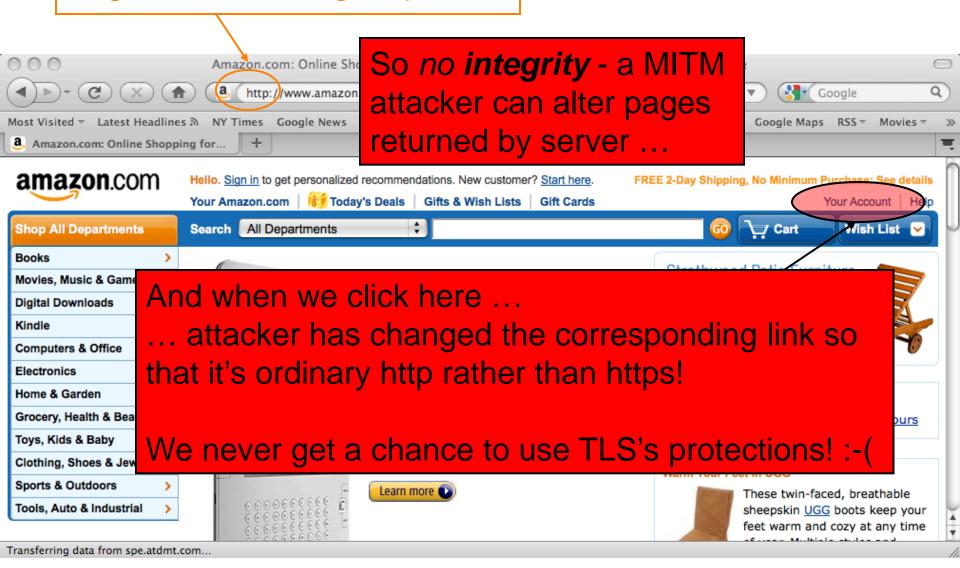
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 - o Note: symmetric key crypto on modern hardware is non-issue
 - Hassle of buying/maintaining certs (fairly minor)
 - Integrating with other sites that don't use HTTPS
 - Latency: extra round trips ⇒ 1st page slower to load

SSL / TLS Limitations, cont.

Problems that SSL / TLS does not take care of ?

- SQL injection / XSS / server-side coding/logic flaws
- Vulnerabilities introduced by server inconsistencies

Regular web surfing: http: URL



"sslstrip" attack

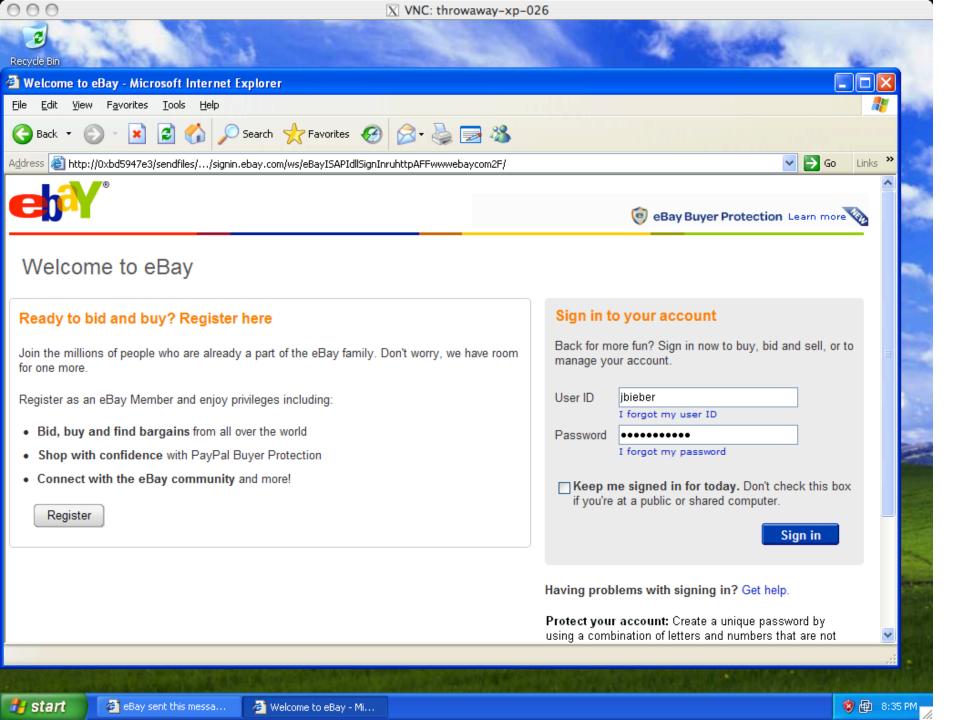
SSL / TLS Limitations, cont.

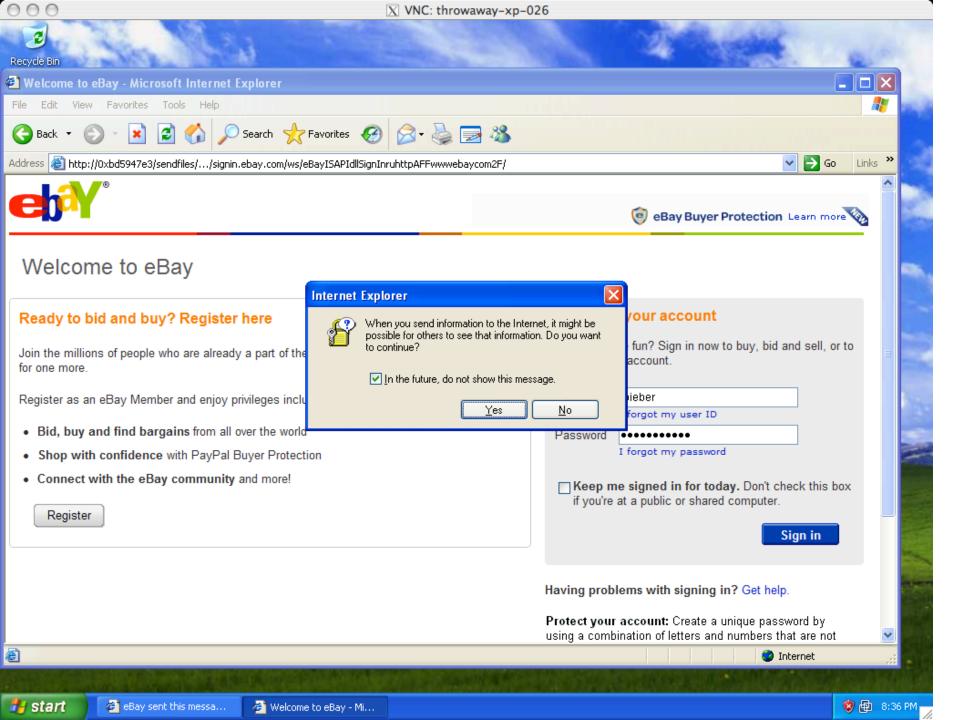
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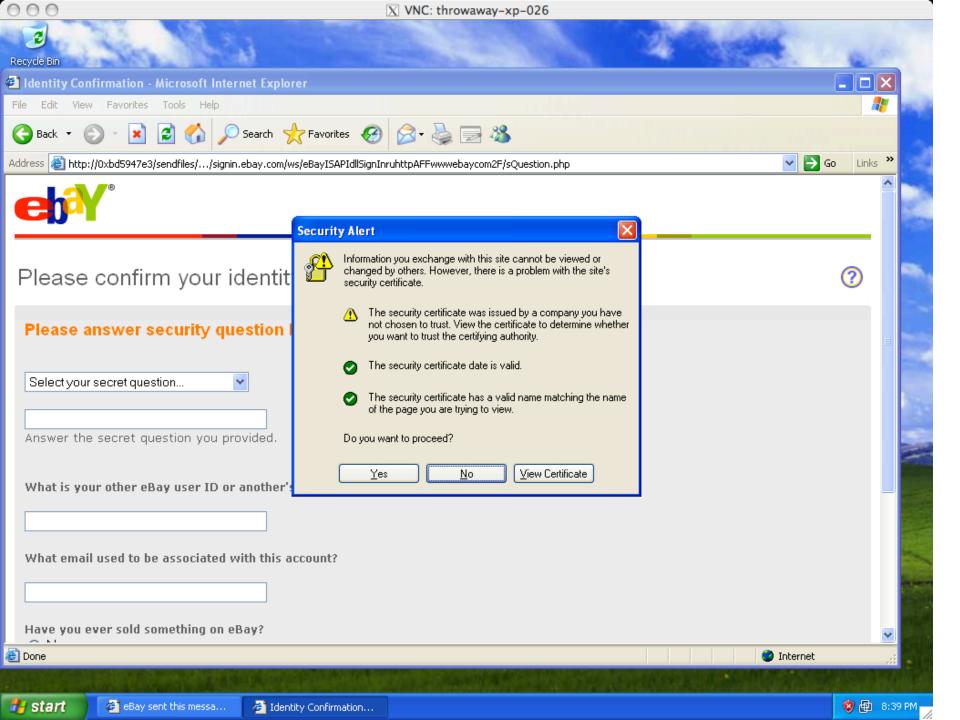
- SQL injection / XSS / server-side coding/logic flaws
- Vulnerabilities introduced by server inconsistencies
- Browser coding/logic flaws
- User flaws
 - Weak passwords
 - Phishing
- Issues of trust ...

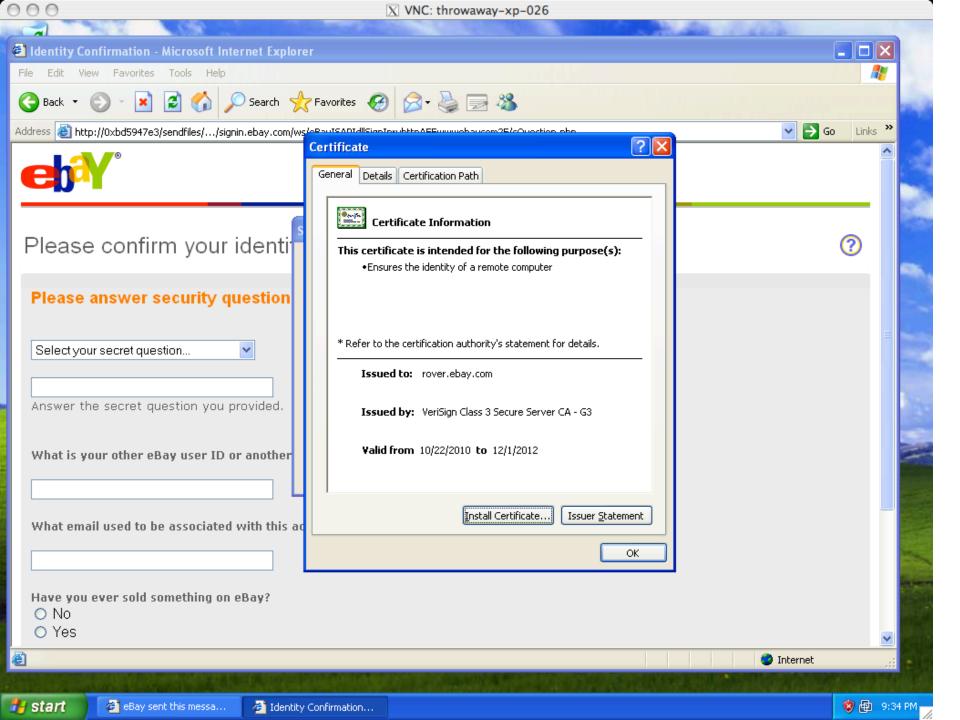
TLS/SSL Trust Issues

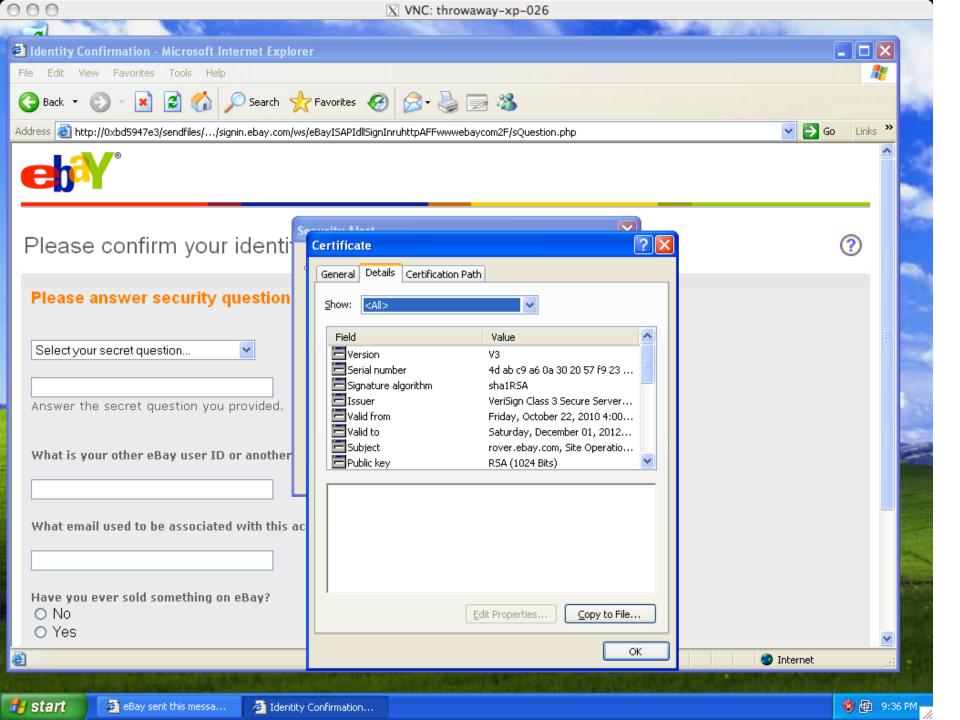
• User has to make correct trust decisions ...

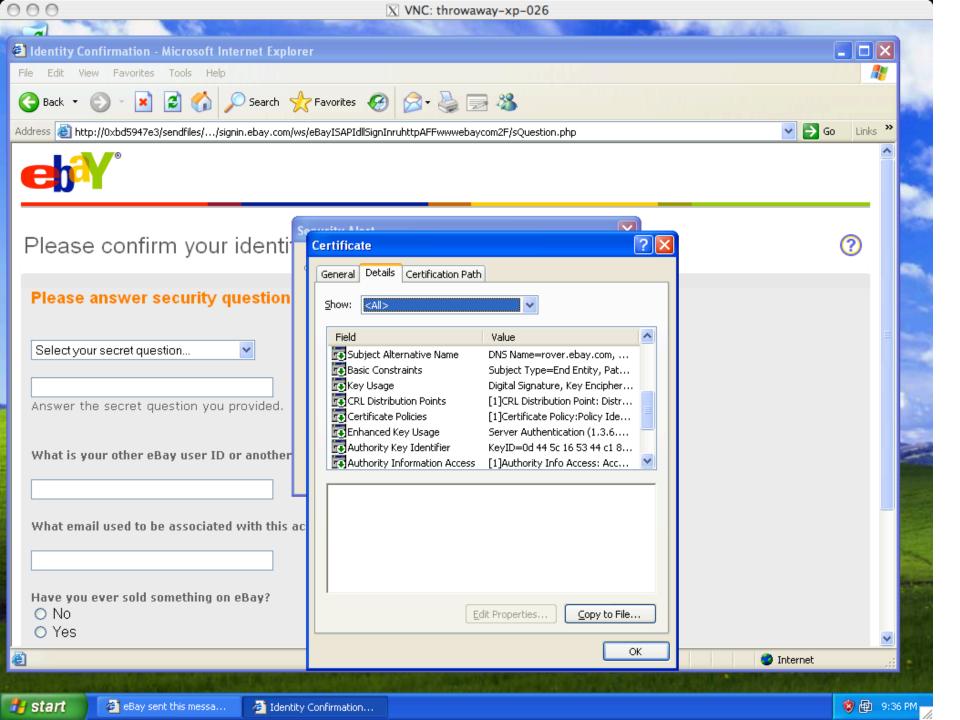


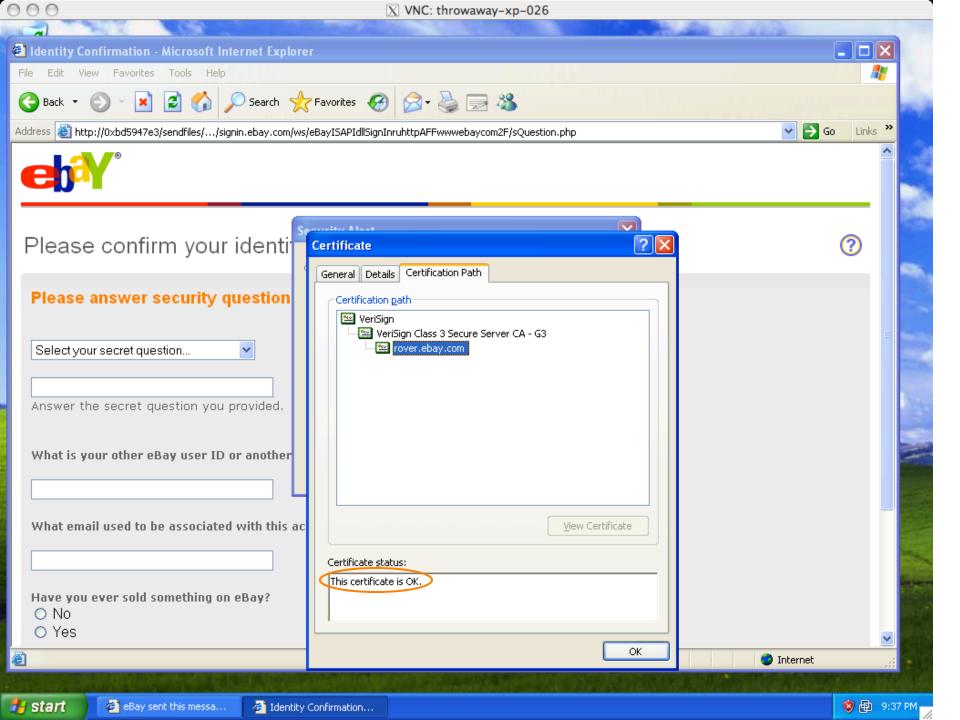




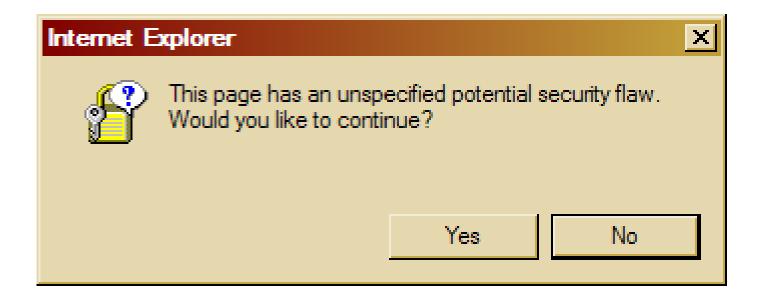








The equivalent as seen by most Internet users:



TLS/SSL Trust Issues, cont.

- "Commercial certificate authorities protect you from anyone from whom they are unwilling to take money."
 - Matt Blaze, circa 2001
- So how many CAs do we have to worry about, anyway?

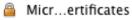
Q



Click to lock the System Roots keychain.

Keychains

login



System

System Roots



[i]

Copy

A-Trust-Qual-02

Root certificate authority

Expires: Tuesday, December 2, 2014 3:00:00 PM PT

This certificate is valid

| | Name | | Kind | Expires | Keychain | |
|--------------------------------|--------------------|----------------------------------|-------------|--------------------------|--------------|---|
| | 1 | A-CERT ADVANCED | certificate | Oct 23, 2011 7:14:14 AM | System Roots | |
| | 100 | A-Trust-nQual-01 | certificate | Nov 30, 2014 3:00:00 PM | System Roots | |
| | 100 | A-Trust-nQual-03 | certificate | Aug 17, 2015 3:00:00 PM | System Roots | |
| | No. | A-Trust-Qual-01 | certificate | Nov 30, 2014 3:00:00 PM | System Roots | • |
| | 100 | A-Trust-Qual-02 | certificate | Dec 2, 2014 3:00:00 PM | System Roots | |
| Category All Items Passwords | - | AAA Certificate Services | certificate | Dec 31, 2028 3:59:59 PM | System Roots | |
| | 27 | AC Raíz Certicámara S.A. | certificate | Apr 2, 2030 2:42:02 PM | System Roots | |
| | | AddTrust Class 1 CA Root | certificate | May 30, 2020 3:38:31 AM | System Roots | |
| | | AddTrust External CA Root | certificate | May 30, 2020 3:48:38 AM | System Roots | |
| Secure Notes | | AddTrust Public CA Root | certificate | May 30, 2020 3:41:50 AM | System Roots | |
| My Certificates | 1 | AddTrust Qualified CA Root | certificate | May 30, 2020 3:44:50 AM | System Roots | |
| | No. | Admin-Root-CA | certificate | Nov 9, 2021 11:51:07 PM | System Roots | |
| Certificates | 100 | AdminCA-CD-T01 | certificate | Jan 25, 2016 4:36:19 AM | System Roots | |
| | No. | AffirmTrust Commercial | certificate | Dec 31, 2030 6:06:06 AM | System Roots | |
| | 1 | AffirmTrust Networking | certificate | Dec 31, 2030 6:08:24 AM | System Roots | |
| | | AffirmTrust Premium | certificate | Dec 31, 2040 6:10:36 AM | System Roots | |
| | 1 | AffirmTrust Premium ECC | certificate | Dec 31, 2040 6:20:24 AM | System Roots | |
| | No. | America Onliation Authority 1 | certificate | Nov 19, 2037 12:43:00 PM | System Roots | |
| | Service Control | America Onliation Authority 2 | certificate | Sep 29, 2037 7:08:00 AM | System Roots | |
| | No. | AOL Time Wcation Authority 1 | certificate | Nov 20, 2037 7:03:00 AM | System Roots | |
| | Service Control | AOL Time Wcation Authority 2 | certificate | Sep 28, 2037 4:43:00 PM | System Roots | |
| | No. | Apple Root CA | certificate | Feb 9, 2035 1:40:36 PM | System Roots | |
| | 1 | Apple Root Certificate Authority | certificate | Feb 9, 2025 4:18:14 PM | System Roots | |
| | No. | Application CA G2 | certificate | Mar 31, 2016 7:59:59 AM | System Roots | Ā |
| | - | ApplicationCA | certificate | Dec 12, 2017 7:00:00 AM | System Roots | ▼ |

167 items

TLS/SSL Trust Issues

- "Commercial certificate authorities protect you from anyone from whom they are unwilling to take money."
 - Matt Blaze, circa 2001
- So how many CAs do we have to worry about, anyway?
- Of course, it's not just their greed that matters ...

News

Solo Iranian hacker takes credit for Comodo certificate attack

Security researchers split on whether 'ComodoHacker' is the real deal

By Gregg Keizer

March 27, 2011 08:39 PM ET



Comments (5) Pecommended (37)





Computerworld - A solo Iranian hacker on Saturday claimed responsibility for stealing multiple SSL certificates belonging to some of the Web's biggest sites, including Google, Microsoft, Skype and Yahoo.

Early reaction from security experts was mixed, with some believing the hacker's claim, while others were dubious.

Last week, conjecture had focused on a state-sponsored attack, perhaps funded or conducted by the Iranian government, that hacked a certificate reseller affiliated with U.S.-based Comodo.

On March 23, Comodo acknowledged the attack, saying that eight days earlier, hackers had obtained nine bogus certificates for the log-on sites of Microsoft's Hotmail, Google's Gmail, the Internet phone and chat service Skype and Yahoo Mail. A certificate for Mozilla's Firefox add-on site was also acquired.

Fraudulent Google certificate points to Internet attack

Is Iran behind a fraudulent Google.com digital certificate? The situation is similar to one that happened in March in which spoofed certificates were traced back to Iran.



by Elinor Mills | August 29, 2011 1:22 PM PDT



A Dutch company appears to have issued a digital certificate for Google.com to someone other than Google, who may be using it to try to re-direct traffic of users based in Iran.

Yesterday, someone reported on a Google support site that when attempting to log in to Gmail the browser issued a warning for the digital certificate used as proof that the site is legitimate, according to this thread on a Google support forum site.

Final Report on DigiNotar Hack Shows Total Compromise of CA Servers

The attacker who penetrated the Dutch CA DigiNotar last year had complete control of all eight of the company's certificate-issuing servers during the operation and he may also have issued some rogue certificates that have not yet been identified. The final report from a

Evidence Suggests DigiNotar, Who Issued Fraudulent Google Certificate, Was Hacked *Years* Ago

from the diginot dept

The big news in the security world, obviously, is the fact that a **fraudulent Google certificate made its way out into the wild**, apparently targeting internet users in Iran. The Dutch company DigiNotar has put out a statement saying that **it discovered a breach** back on July 19th during a security audit, and that fraudulent certificates were generated for "several dozen" websites. The only one known to have gotten out into the wild is the Google one.

TLS/SSL Trust Issues

- "Commercial certificate authorities protect you from anyone from whom they are unwilling to take money."
 - Matt Blaze, circa 2001
- So how many CAs do we have to worry about, anyway?
- Of course, it's not just their greed that matters ...
- ... and it's not just their diligence & security that matters ...
 - "A decade ago, I observed that commercial certificate authorities protect you from anyone from whom they are unwilling to take money. That turns out to be wrong; they don't even do that much." Matt Blaze, circa 2010

Conclusion

- Use SSL/TLS to secure communications end-to-end
- Relies on trustworthiness of certificates