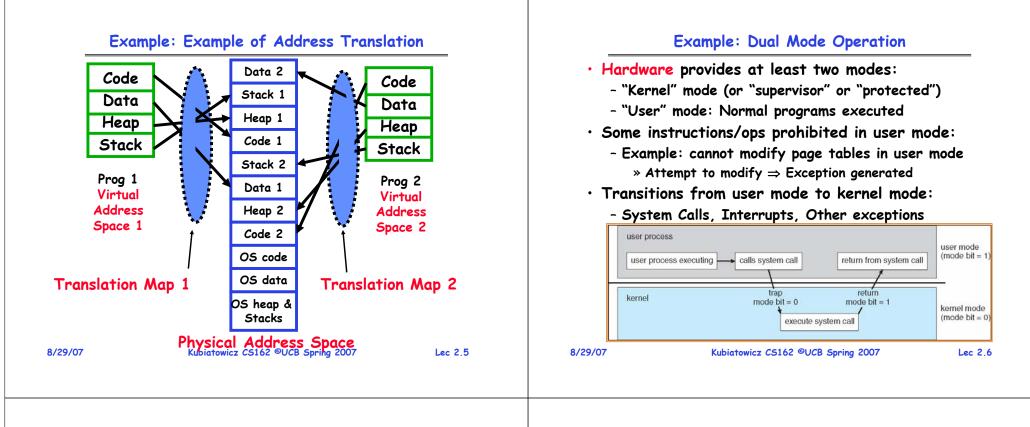
	Review: Virtual Machine Abstraction				
CS162 Openating Systems and	Application				
Operating Systems and Systems Programming Lecture 2	Operating System Physical Machine Interface				
Lecture 2 History of the World Parts 1—5 Operating Systems Structures August 29, 2007 Prof. John Kubiatowicz http://inst.eecs.berkeley.edu/~cs162	 Hardware Software Engineering Problem: Turn hardware/software quirks ⇒ what programmers want/need Optimize for convenience, utilization, security, reliability, etc For Any OS area (e.g. file systems, virtual memory, networking, scheduling): What's the hardware interface? (physical reality) What's the application interface? (nicer abstraction) 2/2/2012 Lec 2.2 				
 Example: Protecting Processes from Each Other Problem: Run multiple applications in such a way that they are protected from one another Goal: Keep User Programs from Crashing OS Keep User Programs from Crashing each other [Keep Parts of OS from crashing other parts?] (Some of the required) Mechanisms: Address Translation Dual Mode Operation Simple Policy: Programs are not allowed to read/write memory of other Programs or of Operating System 	 Example: Address Translation Address Space A group of memory addresses usable by something Each program (process) and kernel has potentially different address spaces. Address Translation: Translate from Virtual Addresses (emitted by CPU) into Physical Addresses (of memory). Apping often performed in Hardware by Memory Addresses (MMU) Virtual Addresses (MMU) 				

Lec 2.3

8/29/07



Goals for Today

- History of Operating Systems
 - Really a history of resource-driven choices
- Operating Systems Structures
- Operating Systems Organizations

Moore's Law Change Drives OS Change

	1981	2006	Factor
CPU MHz,	10	3200x4	1,280
Cycles/inst	3—10	0.25-0.5	6—40
DRAM capacity	128KB	4GB	32,768
Disk capacity	10MB	1TB	100,000
Net bandwidth	9600 b/s	1 Gb/s	110,000
# addr bits	16	32	2
#users/machine	10s	≤ 1	≤ 0.1
Price	\$25,000	\$4,000	0.2

Typical academic computer 1981 vs 2006

Note: Some slides and/or pictures in the following are adapted from slides ©2005 Silberschatz, Galvin, and Gagne. Many slides generated from my lecture notes by Kubiatowicz.

Lec 2.7

Moore's law effects

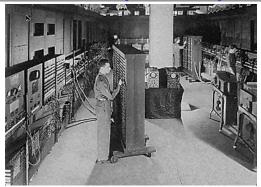
- Nothing like this in any other area of business
- Transportation in over 200 years:
 - 2 orders of magnitude from horseback @10mph to Concorde @1000mph
 - Computers do this every decade (at least until 2002)!
- What does this mean for us?
 - Techniques have to vary over time to adapt to changing tradeoffs
- I place a lot more emphasis on principles
 - The key concepts underlying computer systems
 - Less emphasis on facts that are likely to change over the next few years...
- Let's examine the way changes in \$/MIP has radically changed how OS's work

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Lec 2.9

Dawn of time ENIAC: (1945—1955)



• "The machine designed by Drs. Eckert and Mauchly was a monstrosity. When it was finished, the ENIAC filled an entire room, weighed thirty tons, and consumed two hundred kilowatts of power."

• http://ei.cs.vt.edu/~history/ENIAC.Richey.HTML

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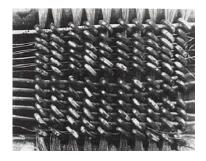
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Lec 2.10

History Phase 1 (1948—1970) Hardware Expensive, Humans Cheap

- When computers cost millions of \$'s, optimize for more efficient use of the hardware!
 Lack of interaction between user and computer
- User at console: one user at a time
- Batch monitor: load program, run, print
- Optimize to better use hardware
 - When user thinking at console, computer idle \Rightarrow BAD!
 - Feed computer batches and make users wait
 - Autograder for this course is similar
- No protection: what if batch program has bug?

Core Memories (1950s & 60s)



The first magnetic core memory, from the IBM 405 Alphabetical Accounting Machine.

- Core Memory stored data as magnetization in iron rings
 - Iron "cores" woven into a 2-dimensional mesh of wires
 - Origin of the term "Dump Core"
 - Rumor that IBM consulted Life Saver company
- See: http://www.columbia.edu/acis/history/core.html

Lec 2.11

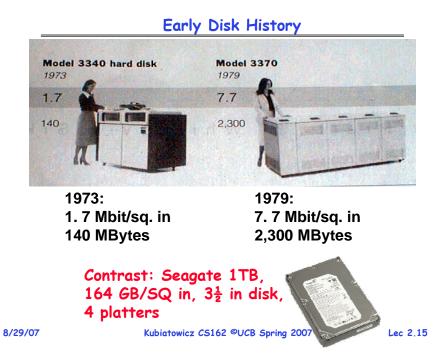
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History Phase $1\frac{1}{2}$ (late 60s/early 70s)

- Data channels, Interrupts: overlap I/O and compute
 - DMA Direct Memory Access for I/O devices
 - I/O can be completed asynchronously
- Multiprogramming: several programs run simultaneously
 - Small jobs not delayed by large jobs
 - More overlap between I/O and CPU
 - Need memory protection between programs and/or OS
- Complexity gets out of hand:
 - Multics: announced in 1963, ran in 1969
 - » 1777 people "contributed to Multics" (30-40 core dev)
 - » Turing award lecture from Fernando Corbató (key researcher): "On building systems that will fail"
 - OS 360: released with 1000 known bugs (APARs) » "Anomalous Program Activity Report"
- OS finally becomes an important science:
 - How to deal with complexity???
 - UNIX based on Multics, but vastly simplified
- 8/29/07

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Lec 2.13



A Multics System (Circa 1976)



• The 6180 at MIT IPC, skin doors open, circa 1976:

- "We usually ran the machine with doors open so the operators could see the AQ register display, which gave you an idea of the machine load, and for convenient access to the EXECUTE button, which the operator would push to enter BOS if the machine crashed."
- http://www.multicians.org/multics-stories.html 8/29/07 Kubiatowicz C5162 ©UCB Spring 2007 Lec 2.14



- Computers available for tens of thousands of dollars instead of millions
- · OS Technology maturing/stabilizing
- Interactive timesharing:
 - Use cheap terminals (~\$1000) to let multiple users interact with the system at the same time
 - Sacrifice CPU time to get better response time
 - Users do debugging, editing, and email online
- Problem: Thrashing
 - Performance very non-linear response with load
 - Thrashing caused by many factors including
 - » Swapping, queueing



Administrivia: What is this CS16x???

- Why change CS162? Only minor changes since 1990's...
 - Slides!
 - Java version of Nachos
 - Content: More crypto/security, less databases and distributed filesystems
 - Time to update again!!
- Most CS students take CS 162 and 186
 - But, not all take EE 122, CS 169/161
 - We'd like all students to have a basic understanding of key concepts from these classes
- Each class introduces the same topics with classspecific biases
 - Concurrency in an Operating System versus in a Database management system
- Introduce concepts with a common framework 8/29/07 Kubiatowicz C5162 ©UCB Spring 2007 Lec 2.17

Administrivia: CS 194-3/16x

- Mondays and Wednesdays 9-10:30 in 306 Soda
 - Taught by Anthony Josephy: high teaching ratings!
- Primary content is similar to CS 162
 - With CS 186, 161, and 169, and EE 122 topics
- 4 units with CS Upper Division credit
- 3-4 Projects (tentative)
 - Nachos Phase 1
 - Multi-core programming
 - Secure iTunes-like e-commerce site with a Peer-to-Peer content distribution network
- We need some bold students to try the course
 - Might need to be cancelled otherwise
 - Great way to get 186 & 122 material as well
- Targeted at Sophomores/First term Juniors 8/29/07 Kubiatowicz C5162 ©UCB Spring 2007

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Administrivia: Back to CS162

- Cs162-xx accounts:
 - Make sure you got an account form
 - » We have more forms for those of you who didn't get one
 - If you haven't logged in yet, you need to do so
- Nachos readers:
 - TBA: Will be down at Copy Central on Hearst
 - Will include lectures and printouts of all of the code
- Video archives available off lectures page
 - Just click on the title of a lecture for webcast
 - Only works for lectures that I have already given!
 - Still working on Webcast
- No slip days on first design document for each phase
 - Need to get design reviews in on time
- Don't know Java well?
 - Talk CS 9G self-paced Java course

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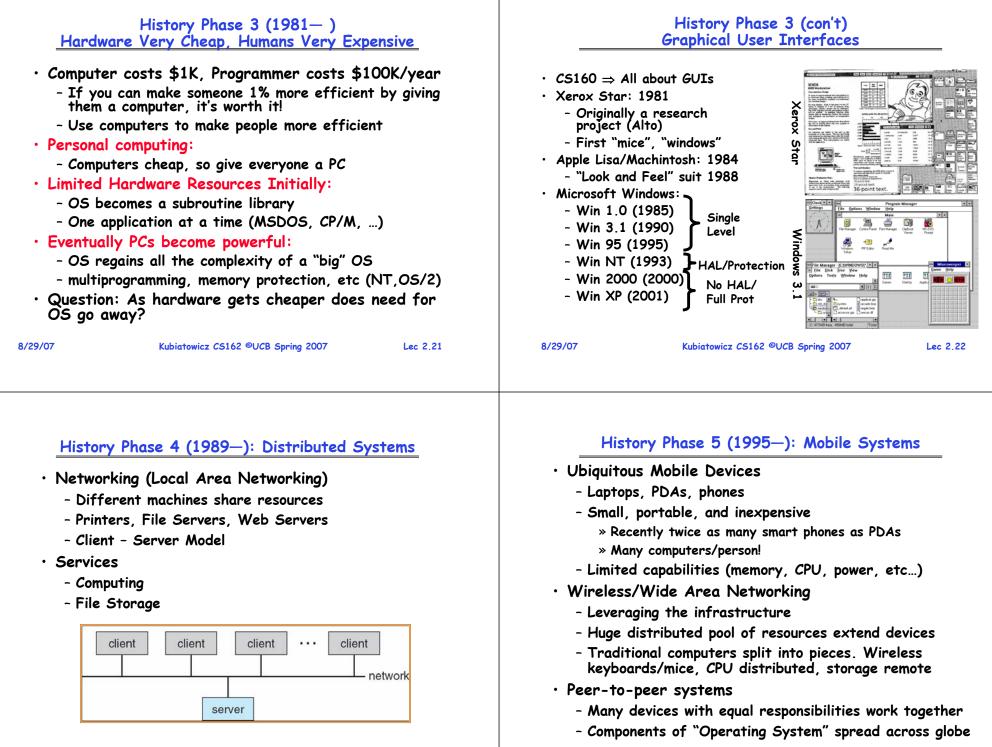
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Administriva: Almost Time for Project Signup

- Project Signup: Watch "Group/Section Assignment Link"
 - 4-5 members to a group
 - » Everyone in group must be able to *actually* attend same section
 - » The sections assigned to you by Telebears are temporary!
 - Only submit once per group!
 - » Everyone in group must have logged into their cs162-xx accounts once before you register the group
 - » Make sure that you select at least 2 potential sections
 - » Due date: Thursday 9/6 by 11:59pm
- Sections:
 - No sections tomorrow
 - Go to Telebears-assigned Section next week

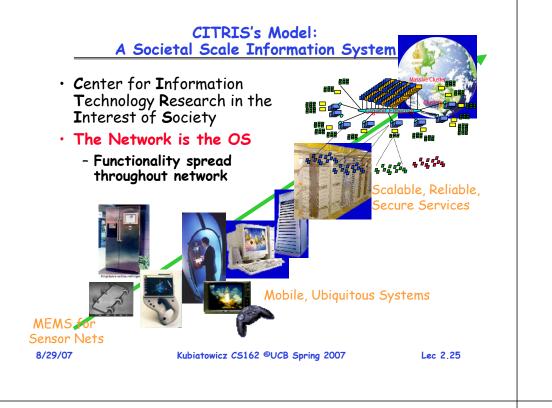
	Section	Time Location		ТА	
Γ	101	Th 10:00-11:00A	81 Evans	Kelvin Lwin	
	102	Th 12:00-1:00P 155 Barrows		Kelvin Lwin	
	103	Th 2:00-3:00P	75 Evans	Todd Kosloff	
	104	Th 4:00-5:00P	B51 Hildebrand	Todd Kosloff	
8	105	F 10:00-11:00A	4 Evans	Thomas Kho	



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Moore's Law Reprise: Modern Laptop

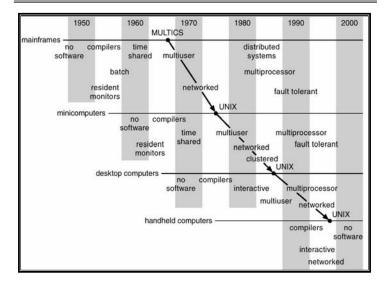
	1981	2005	2006 Ultralight Laptop
CPU MHz,	10	3200×4	1830
Cycles/inst	3—10	0.25-0.5	0.25-0.5
DRAM capacity	128KB	4GB	2GB
Disk capacity	10MB	1TB	100 <i>G</i> B
Net bandwidth	9600 b/s	1 Gb/s	1 Gb/s (wired) 54 Mb/s (wireless) 2 Mb/s (wide-area)
# addr bits	16	32	32
#users/machine	10s	≤ 1	$\leq \frac{1}{4}$
Price	\$25,000	\$4,000	\$2500

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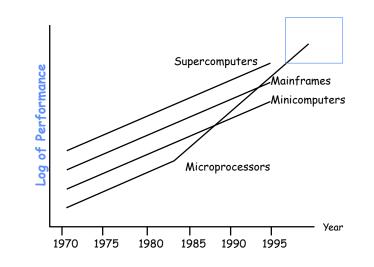
Lec 2.26

Migration of Operating-System Concepts and Features

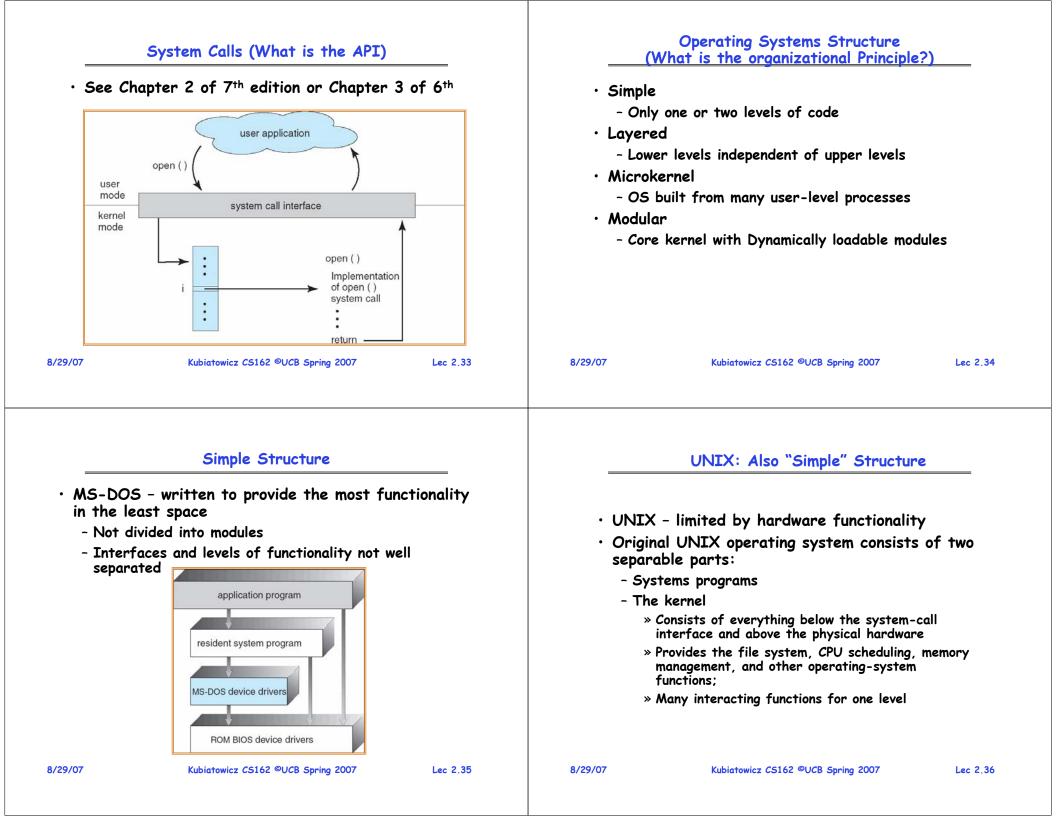


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Compare: Performance Trends (from CS152)



Lec 2.27



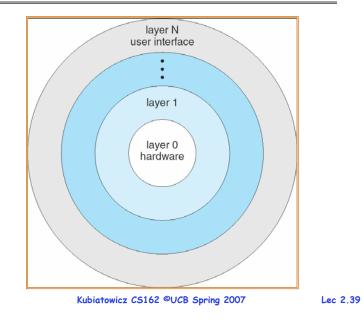
UNIX System Structure

User Mode	20	Applications	(the users)			
		Standard Libs shells and commands compilers and interpreters system libraries				
		syster	n-call interface to the ke	ernel		
Kernel Mode	Kernel	signals terminal handling character I/O system terminal drivers file system swapping block I/O system disk and tape drivers		CPU scheduling page replacement demand paging virtual memory		
		kerne	kernel interface to the hardware			
Hardware		terminal controllers terminals device controllers disks and tapes		memory controllers physical memory		
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Layered Structure

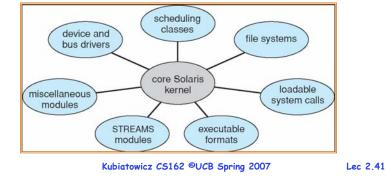
 Operating system is divided many layers (levels) Each built on top of lower layers Bottom layer (layer 0) is hardware Highest layer (layer N) is the user interface Each layer uses functions (operations) and services of only lower-level layers Advantage: modularity ⇒ Easier debugging/Maintenance Not always possible: Does process scheduler lie above or below virtual memory layer? Need to reschedule processor while waiting for paging May need to page in information about tasks Important: Machine-dependent vs independent layers Easier migration between platforms Easier evolution of hardware platform Good idea for you as well!
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Microkernel Structure Moves as much from the kernel into "<i>user</i>" space Small core OS running at kernel level
- OS Services built from many independent user-level processes
 Communication between modules with message passing Benefits: Easier to extend a microkernel
 Easier to port OS to new architectures More reliable (less code is running in kernel mode) Fault Isolation (parts of kernel protected from other parts) More secure
 Detriments: Performance overhead severe for naïve implementation
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Layered Operating System



Modules-based Structure

- Most modern operating systems implement modules
 - Uses object-oriented approach
 - Each core component is separate
 - Each talks to the others over known interfaces
 - Each is loadable as needed within the kernel
- Overall, similar to layers but with more flexible



Implementation Issues (How is the OS implemented?)

- Policy vs. Mechanism
 - Policy: What do you want to do?
 - Mechanism: How are you going to do it?
 - Should be separated, since both change
- Algorithms used
 - Linear, Tree-based, Log Structured, etc...
- Event models used
 - threads vs event loops
- Backward compatability issues
 - Very important for Windows 2000/XP
- System generation/configuration
 - How to make generic OS fit on specific hardware

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Conclusion

- Rapid Change in Hardware Leads to changing OS
 - Batch \Rightarrow Multiprogramming \Rightarrow Timeshare \Rightarrow Graphical UI \Rightarrow Ubiquitous Devices \Rightarrow Cyberspace/Metaverse/??
- \cdot OS features migrated from mainframes \Rightarrow PCs
- Standard Components and Services
 - Process Control
 - Main Memory
 - I/O
 - File System
 - UI
- Policy vs Mechanism
 - Crucial division: not always properly separated!
- Complexity is always out of control
 - However, "Resistance is NOT Useless!"

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