

Lecture 8: Practical Bison: Error Handling, etc.

HackJam:

- Hackers@Berkeley "HackJam"—a 12 hour hackathon hosted by Hackers@Berkeley and sponsored by Box.
- There will be food served throughout the event and prizes awarded at the end.
- Who should come: Anyone interested in hacking, regardless of experience. There will be helpful students and engineers from Box there to help anyone who wants to learn.
- Time: 11am-11pm Saturday, September 29th.
- Place: Wozniak Lounge, Soda Hall.

Programming Contest: Also this Saturday, from 10-3:30.

Parsing Errors

- One purpose of the parser is to filter out errors that show up in parsing
- Later stages should not have to deal with possibility of malformed constructs
- Parser must *identify* error so programmer knows what to correct
- Parser should *recover* so that processing can continue (and other errors found).
- Parser might even *correct* error (e.g., PL/C compiler could “correct” some Fortran programs into equivalent PL/1 programs!)

Identifying Errors

- All of the valid parsers we've seen identify syntax errors as soon as possible.
- *Valid prefix property*: all the input that is shifted or scanned is the beginning of some valid program...
- ... But the rest of the input might not be.
- So in principle, deleting the lookahead (and subsequent symbols) and inserting others will give a valid program.

Automating Recovery

- Unfortunately, best results require using semantic knowledge and hand tuning.
 - E.g., $a(i).y = 5$ might be turned to $a[i].y = 5$ if a is statically known to be a list, or $a(i).y = 5$ if a is a function.
- Some automatic methods can do an OK job that at least allows parser to catch more than one error.

Bison's Technique

- The special terminal symbol error is never actually returned by the lexer.
- Gets inserted by parser in place of erroneous tokens.
- Parsing then proceeds normally.

Example of Bison's Error Rules

Suppose we want to throw away bad statements and carry on

```
stmt : whileStmt
    | ifStmt
    | ...
    | error NEWLINE
    ;
```

Response to Error

- Consider erroneous text like

```
if x y: ...
```

- When parser gets to the *y*, will detect error.
- Then pops items off parsing stack until it finds a state that allows a shift or reduction on 'error' terminal
- Does reductions, then shifts 'error'.
- Finally, throws away input until it finds a symbol it can shift after 'error', according to the grammar.

Error Response, contd.

- So with our example:

```
stmt : whileStmt
      | ifStmt
      | ...
      | error NEWLINE
      ;
```

We see 'y', throw away the 'if x', so as to be back to where a stmt can start.

- Shift 'error' and throw away more symbols to NEWLINE. Then carry on.

Of Course, It's Not Perfect

- “Throw away and punt” is sometimes called “panic-mode error recovery”
- Results are often annoying.
- For example, in our example, there could be an INDENT after the NEWLINE, which doesn't fit the grammar and causes another error.
- Bison compensates in this case by not reporting errors that are too close together
- But in general, can get cascade of errors.
- Doing it right takes a lot of work.

Bison Examples

[See lecture15 directory.]