Section 11: First-Class Functions

1. What happens when we call a function?

Complete the following diagram, which represents the stack immediately after the function call **(f 3 6 9)**. Assume that rsp is aligned to a multiple of 16 when we compile the call. By "rsp" in the left column, we mean the value of rsp before the function call.

Address	Value
rsp - 56	
rsp - 48	
rsp - 40	
rsp - 32	
rsp - 24	
rsp - 16	
rsp - 8	
rsp	<return address="" caller's="" for="" frame=""></return>

Where does the stack frame of the callee start?

How would the stack frame change if we dynamically communicate the number of arguments like we did in Homework 6?

2. What's the use of function pointers without lambdas?

Write a program in our language that only works once we support function pointers (do *not* include any lambda expressions at this point).

What can function pointers be useful for?

3. Handling lambdas <u>without</u> free variables

The C11 standard supports function pointers, but no lambdas. How could you change the following functions to be a valid C11 code while still maintaining its functionality?

```
typedef struct node {
   int value;
   struct node* next;
} node;
node* map(node* list, int (*f)(int)) {
   if(list == NULL) return NULL;
   node* next = map(list->next, f);
   node* new = malloc(sizeof(node));
   new->next = next;
   new->value = f(list->value);
   return new;
}
node* square(node* list) {
   return map(list, [](int n) { return n * n; });
                   ----- This is the problematic
   11
   11
                                               section b/c lambdas are
                                               not supported in C11.
   11
}
```

Interactive Version (runs in your browser!): <u>https://godbolt.org/z/jjhcafffT</u>

4. Handling lambdas <u>with</u> free variables

We will now explore what happens when we compile the following lambda in an environment where the variables \mathbf{x} and \mathbf{y} are defined:

(lambda (z) (+ z (+ y x)))

4.1 Compiling the creation of a lambda

4.1.1 Setting up the heap

On the left is the stack layout at the site where the lambda is created.

On the right, fill in the values that are transferred to the heap when we compile the lambda. *Assume that the label of the lambda implementation is called* **_1***ambda_1*.

Stack		Неар		
Address	Value	Address	Value	
rsp - 16	У	rdi		
rsp - 8	х	rdi + 8		
rsp	<return address=""></return>	rdi + 16		

Why can't we just inline the values of \mathbf{x} and \mathbf{y} in the body of the lambda?

4.1.2 Stack layout going into the body

We now want to explore the stack and heap layout that we would expect to see right after we jump to the label of the lambda function. Please show where **x**, **y** and **z** are located on the stack.

Stack				
Address	Value			
rsp - 24				
rsp - 16				
rsp - 8				
rsp	<return address=""></return>			
Who is responsible for placing \mathbf{x} and \mathbf{y} there? \Box caller \Box callee			callee	
Who is respo	nsible for placing z there?	caller	callee	
How do we obtain the values of \mathbf{x} and \mathbf{y} ?				

4.2 Compiling a call to a lambda

Finally, let's talk about how a lambda function is called. Feel free to look at the **compile.ml** file of the class compiler, starting at the line | **Call (f, args) when not is_tail ->.**

Please fill in the blanks in the following pseudocode that describes how to call a lambda function.

1.	Compile the arguments and place them on the	
2.	Compile the expression.	
3.	Ensure that rax is tagged with	
4.	Copy the content of rax to the	
	a. This is important for the function to be able to retrieve the and copy them onto the stack.	
5.	Load the address of the function into rax from the	by
	subtracting the type tag from	·
6.	Perform a call to the address now stored in	_, modifying the
	stack pointer before and after for alignment purposes.	

Here's some blank diagrams that you can use to help yourself visualize the algorithm!

Stack			Неар		
Address	Value		Address	Value	
rsp - 24			rdi		
rsp - 16		-	rdi + 8		
rsp - 8		-	rdi + 16		
rsp		_	rdi + 24		

5. The spectrum of first-class functions

Assume that variadic arguments and the *apply* function from Homework 6 are not supported.

		First-class functions				
	No first-class functions	Functions pointers w/o lambdas	Lambdas (w/o free variables)	Lambdas (w/ free variables)		
Is the function label known at the call site?						
Is the number of expected arguments known at the call site?						
Where is the code for the function generated?						
How is the function label created?						
How do we know which address to jump to at the call site?						
Do we need to use the heap? Why?						