

Section 9: Exam Review

Function Calls

CS164

Consider this program. Match each piece of data to its location relative to `*entry's*` starting `rsp`. * 6 points

```
(define (add3 x y z)
  (+ x (+ y z)))
```

```
(print (add3 1 2 3))
```

	rsp - 0	rsp - 8	rsp - 16	rsp - 24	rsp - 32	rsp - 40	rsp - 48
y	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
entry's return address	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
x	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
padding cell for stack alignment	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
add3's return address	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
z	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>

Drill Question

Stack Review:

Draw a Stack Diagram that Illustrates a
Function Call

(Feel free to consult lecture notes on Functions)

Name 2 differences between calling C functions vs. native functions.

Why do we need to align the stack pointer when calling a native function?

Practice Question

In this question you will complete a pseudocode description of how to compile function calls in our language (like in class).

Instructions: For each line in the pseudocode, circle exactly one of the alternatives separated by a slash in each set of { curly braces }.

How to compile (foo e1 e2) assuming that our stack index is currently set to stack_index:

1. Compile e1 with the stack index starting at stack_base { + / - } { 8 / 16 / 24 / 32 / 40 / 48 }. The resulting value will be stored in rax.
2. Store the resulting value on the { stack / heap }.
3. Compile e2 with the stack index starting at stack_base { + / - } { 8 / 16 / 24 / 32 / 40 / 48 }. The resulting value will be stored in rax.
4. Store the resulting value on the { stack / heap }.
5. { incremented / decremented } Rsp by { stack_base / (stack_base + 16) / (stack_base - 16) / stack_index / (stack_index + 16) / (stack_index - 16) }
6. { Save Rdi to the stack / Save Rdi to the heap / nothing }
7. Call the label for foo.
8. { incremented / decremented } Rsp by { stack_base / (stack_base + 16) / (stack_base - 16) / stack_index / (stack_index + 16) / (stack_index - 16) }
9. { Restore Rdi from the stack / Restore Rdi from the heap / nothing }

Practice Question

In this question you will complete a pseudocode description of how to compile function calls in our language (like in class).

Instructions: For each line in the pseudocode, circle exactly one of the alternatives separated by a slash in each set of **{ curly braces }**.

How to compile (foo e1 e2) assuming that our stack index is currently set to stack_index:

1. Compile e1 with the stack index starting at stack_base { + / - } { 8 / 16 / 24 / 32 / 40 / 38 }. The resulting value will be stored in rax.
2. Store the resulting value on the { stack / heap }.
3. Compile e2 with the stack index starting at stack_base { + / - } { 8 / 16 / 24 / 32 / 40 / 38 }. The resulting value will be stored in rax.
4. Store the resulting value on the { stack / heap }.
5. { incremented / decremented } Rsp by { stack_base / (stack_base + 16) / (stack_base - 16) / stack_index / (stack_index + 16) / (stack_index - 16) }
6. { Save Rdi to the stack / Save Rdi to the heap / nothing }
7. Call the label for foo.
8. { incremented / decremented } Rsp by { stack_base / (stack_base + 16) / (stack_base - 16) / stack_index / (stack_index + 16) / (stack_index - 16) }
9. { Restore Rdi from the stack / Restore Rdi from the heap / nothing }

Note: This question is somewhat ill-defined. Exam questions will have one well defined answer!

Problems from last week

For each expression below, will it result in a **compile-time** error, a **run-time** error, or a **valid** result when run in the **class compiler**?

- A.

```
(let ((v1 (vector 1 true)))  
    (let ((v2 (vector 1 true)))  
        (vector-get v1 1)))
```
- B.

```
(vector-length (pair 1 2))
```
- C.

```
(list? (pair 1 2))
```
- D.

```
(vector-get (if (not (char? #\a)) (vector 1 true) (vector 2 true)) 1)
```

Problems from Last Week

The OCaml code below implements a **left** projection operation in the compiler. Write **three** s-expressions of the form **(left ...)**. One that will result in a **run time** error, one that will result in a **compile time** error, and one that will result in the **number 22**.

```
| Lst [Sym "left"; e] ->  
  compile_exp tab stack_index e  
  @ ensure_pair (Reg Rax)  
  @ [Mov (Reg Rax, MemOffset (Reg Rax, Imm (-pair_tag)))]
```

Problems from Last Week

What does the **interpreter symbol table** look like at each point in the code?

```
(let ((z false) (x 10))
  do (A_EXP)
      (pair 1
            (if z
                (let (x (+ x 1)) (let (y 9) B_EXP))
                (let (y 8) (C_EXP)))
          (x + 1)
          (D_EXP)))
```

A_EXP: {z:false; x:10}

B_EXP: _____

C_EXP: _____

D_EXP: _____

Problems from Last Week

What does the **compiler symbol table** look like?
(var stack index in order of appearance)

```
(let ((z false) (x 10)
      do (A_EXP)
          (pair 1
                (if z
                    (let (x (+ x 1)) (let (y 9) B_EXP))
                    (let (y 8) (C_EXP))
                )
          (x + 1)
          (D_EXP)))
```

A_EXP: _{z:1; x:2}_
B_EXP: _____

C_EXP: _____
D_EXP: _____

Questions