RUNTIME ORGANIZATION (Solutions)

In the Objective Caml language, we can define the following functions.

```
(* Compute n choose k *)
let binom n k = ...
let test x y =
  let a = binom x y in
  let b = binom x (y+1) in
  a + b (* Return the sum *)
```

If we compile this code and then disassemble the result we get the following (using Intel syntax):

```
08049950 <camlTest__test_61>:
 8049950:
                sub
                                                      ; esp -= 12
                                                      ; *(esp+4) = eax
                       DWORD PTR [esp+0x4],eax
 8049953:
                mov
 8049957:
                       DWORD PTR [esp],ebx
                                                      ; *esp = ebx
                mov
                       8049930 <camlTest__binom_58>
 804995a:
                call
                                                      ; call binom
 804995f:
                       DWORD PTR [esp+0x8],eax
                                                      ; *(esp+8) = eax
                mov
                       ebx,DWORD PTR [esp]
                                                      ; ebx = *esp
 8049963:
                mov
                                                      ; ebx += 2
                       ebx,0x2
 8049966:
                add
                       eax, DWORD PTR [esp+0x4]
 8049969:
                mov
 804996d:
                call
                       8049930 <camlTest_binom_58>
 8049972:
                mov
                       ebx,DWORD PTR [esp+0x8]
 8049976:
                       eax, [ebx+eax*1-0x1]
                                                      ; eax = ebx + eax - 1
                lea
 804997a:
                add
                       esp,0xc
 804997d:
                ret
```

Describe the calling conventions used for binom. Where are the arguments n and k stored? Where is the result located on return?

The argument n is passed in the eax register, and k in ebx. We can tell this because eax has the same value both times binom is called, but the second time ebx's value is incremented (by 2 instead of 1 because OCaml uses the least significant bit as a tag bit). The return value is passed in eax, which we can tell because we can trace the two values added for a + b back to the values of eax right after the two calls to binom.

Draw a diagram showing the layout of the stack and the register file right before the second call to binom. Your diagram should show where each argument and local variable is stored.

| Ste | tack (growing downward): | | Registers: | |
|-----|--------------------------|--|------------|-------|
| | RA | | | x |
| | a | | ebx: | y + 1 |
| | Х | | | |
| | у | | | |

Now consider the following assembly code.

```
080483b4 <test>:
 80483b4:
                push
                        ebp
                                                     ; esp -= 4; *esp = ebp
 80483b5:
                        ebp,esp
                 mov
 80483b7:
                 sub
                        esp,0x18
 80483ba:
                        DWORD PTR [ebp-0x14],ecx
                 mov
 80483bd:
                        DWORD PTR [ebp-0x18],edx
                 mov
 80483c0:
                        eax, DWORD PTR [ebp-0x14]
                 mov
 80483c3:
                 mov
                        edx, DWORD PTR [eax]
 80483c5:
                        eax, DWORD PTR [ebp-0x18]
                 mov
 80483c8:
                        eax, DWORD PTR [eax]
                 mov
 80483ca:
                        edx,eax
                                                     ; edx = edx * eax
                 imul
 80483cd:
                        eax, DWORD PTR [ebp-0x14]
                 mov
 80483d0:
                        ecx, DWORD PTR [eax+0x4]
                 mov
                        eax, DWORD PTR [ebp-0x18]
 80483d3:
                 mov
                        eax, DWORD PTR [eax+0x4]
 80483d6:
                 mov
 80483d9:
                 imul
                        eax,ecx
 80483dc:
                 lea
                        eax, [edx+eax*1]
                                                   ; eax = edx + eax
 80483df:
                        DWORD PTR [ebp-0x4], eax
                 mov
                        eax, DWORD PTR [ebp-0x4]
 80483e2:
                 mov
 80483e5:
                 leave
                                                    ; esp = ebp; pop ebp
 80483e6:
                 ret
```

This function takes two parameters, passed in registers ecx and edx respectively. Its result is returned in register eax. Decompile (translate) this assembly into equivalent C code. Hint: This code implements a well-known mathematical operation.

The function implements a two-dimensional dot product.

```
struct vector_t {
  int x;
  int y;
};

int dot(struct vector_t* u, struct vector_t* v) {
  int p;
  p = u->x * v->x + u->y * v->y;
  return p;
}
```

Reasoning: The function uses three different offsets from ebp, so we have three local variables. Two of these (ebp-0x14 and ebp-0x18) correspond to the arguments; we know this because the argument registers are stored to them and they are never overwritten. The third variable (ebp-0x4) is an int because it stores a sum of products of 32-bit integers.

We can tell the arguments are pointers because we dereference them. We use two offsets with each of them and treat the dereferenced values as ints, so we know they point to int arrays or structs with int members. We interpret them as structs here because we recognize the form of the dot product, but it would be equally correct to give the arguments type int *.