

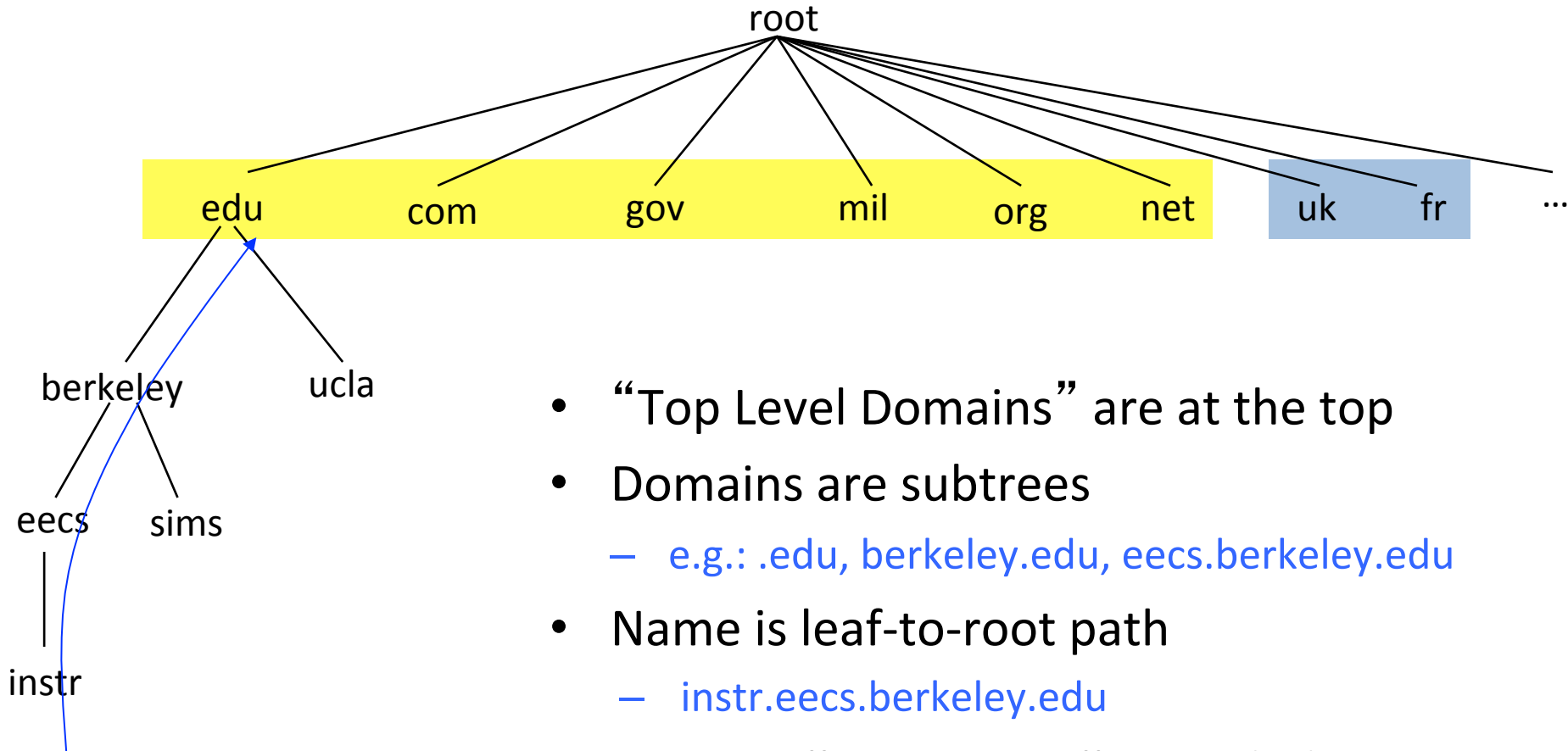
DNS and HTTP

CS 168

Domain Name Service

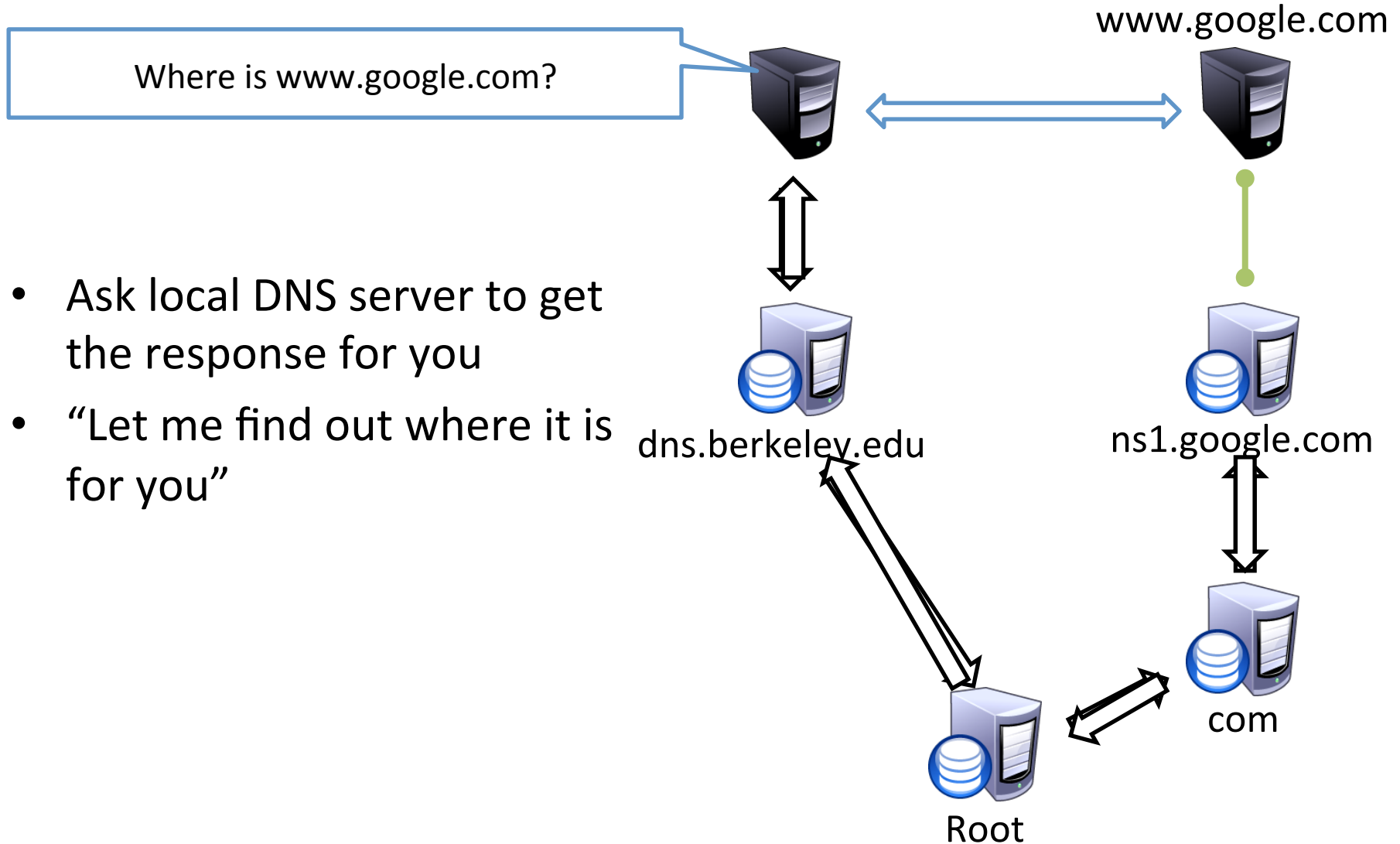
- Host addresses: e.g., 169.229.131.109
 - a number used by protocols
 - conforms to network structure (the “where”)
- Host names: e.g., instr.eecs.berkeley.edu
 - mnemonic name usable by humans
 - conforms to organizational structure (the “who”)
- The Domain Name System (DNS) is how we map from one to the other
 - a **directory service** for hosts on the Internet

Hierarchical Namespace

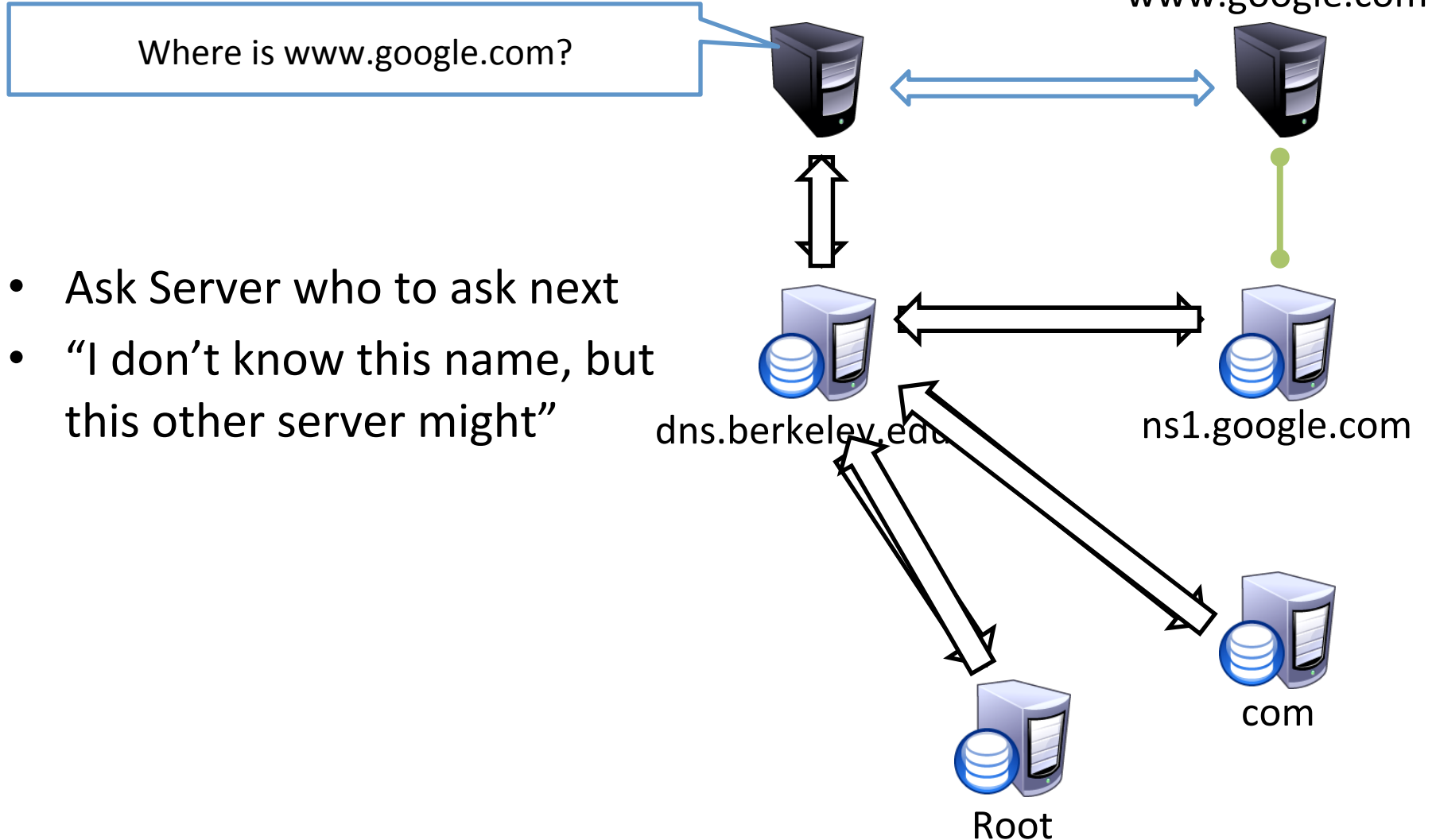


- “Top Level Domains” are at the top
- Domains are subtrees
 - e.g.: .edu, berkeley.edu, eecs.berkeley.edu
- Name is leaf-to-root path
 - instr.eecs.berkeley.edu
- Name collisions trivially avoided!
 - each domain’s responsibility

Recursive DNS Query



Iterative DNS query



DNS Records

- DNS info. stored as **resource records (RRs)**
 - RR is (name, value, type, TTL)
- Type = A: (-> Address)
 - name = hostname
 - value = IP address
- Type = NS: (-> Name Server)
 - name = domain
 - value = name of dns server for domain

DNS Records (contd.)

- Type = CNAME: (-> Canonical NAME)
 - name = hostname
 - value = canonical name
- Type = MX: (-> Mail eXchanger)
 - name = domain in email address
 - value = canonical name(s) of mail server(s)

Fun with dig!

Hyper Text Transfer Protocol (HTTP)

- Client-server architecture
 - server is “always on” and “well known”
 - clients initiate contact to server
- Synchronous request/reply protocol
 - Runs over TCP, Port 80
- Stateless
- ASCII format

Client/Server communication

(method) (resource) (protocol version)
GET /somedir/page.html HTTP/1.1
Host: www.someschool.edu
User-agent: Mozilla/4.0
Connection: close
Accept-language: fr
(blank line)

(header) ← HTTP Request
(Client to Server)

(protocol version) (status code) (status phrase)
HTTP/1.1 200 OK
Connection close
Date: Thu, 06 Aug 2006 12:00:15 GMT
Server: Apache/1.3.0 (Unix)
Last-Modified: Mon, 22 Jun 2006 ...
Content-Length: 6821
Content-Type: text/html
(blank line)
(data) { data data data data data ...

HTTP Response (Server to Client)

HTTP's stateless-ness

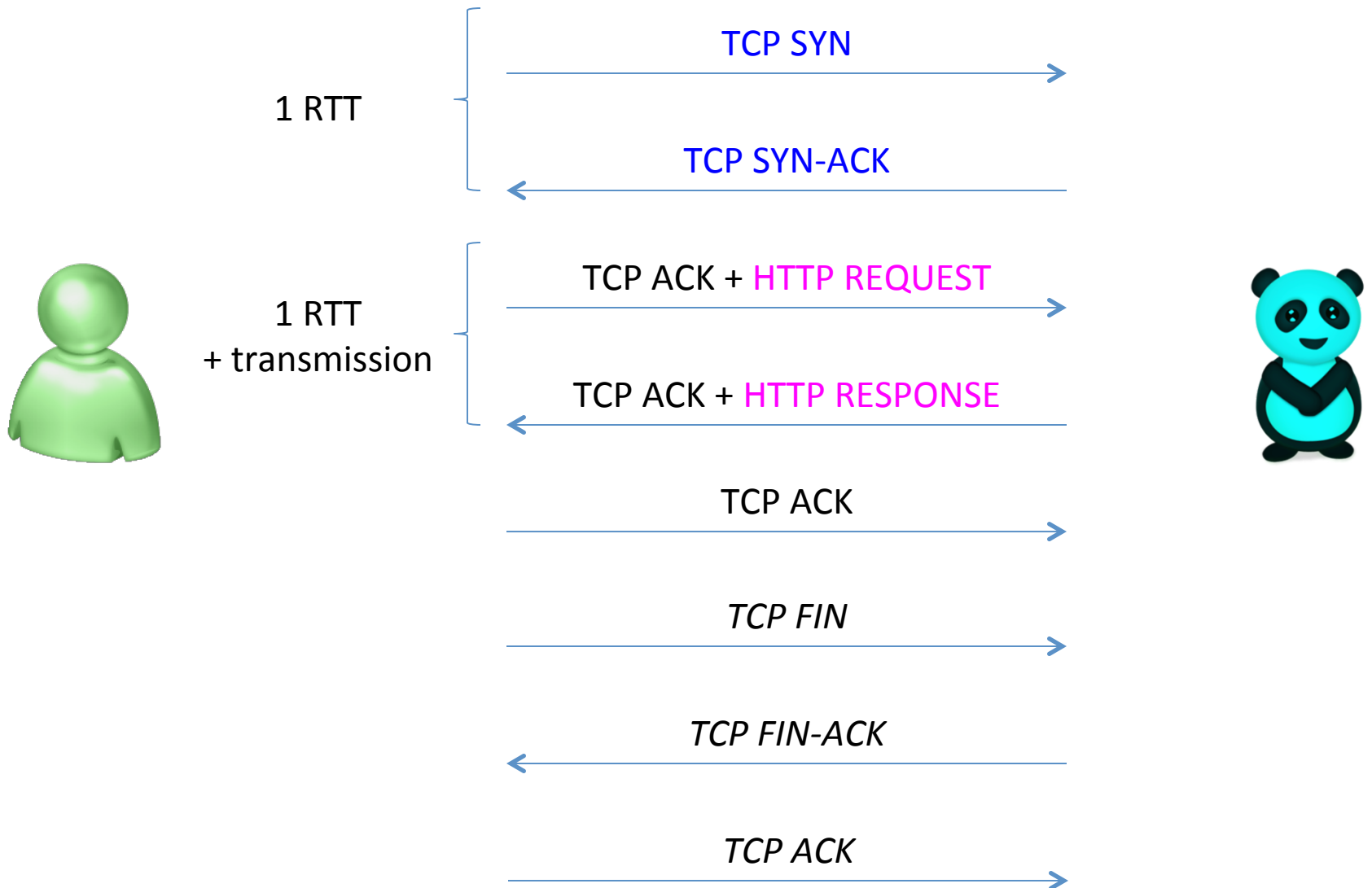
- Pros?
 - Scalable
 - Easier to k
 - Order of r
- Cons?
 - Can't keep
- Solution?
 - Client-side

Cookies!



profiles...)

HTTP Performance: Non-persistent TCP Connection



Other options?

- **Concurrent** Requests and responses
 - Use multiple connections *in parallel*
- **Persistent** Connections
 - Maintain TCP connection across multiple requests
- **Pipelined** Requests and Responses
 - Batch requests and responses to reduce the number of packets



Easy ways to order!

1

1. Go to store
2. Order burger
3. Go to store
4. Order drink
5. Go to store
6. Order fries

2

1. Go to store
2. Order burger
3. Order drink
4. Order fries

3

1. Go to store with two friends
2. Each person orders one item (in parallel)

4

1. Go to store
2. Order burger, drink and fries



2

Cheeseburger, French Fries, and Medium Drink

Q2

c	Page	Media 1	Media 2	Media 3	Total
Sequential requests with non-persistent TCP connections	1 RTT (TCP) 1 RTT (HTTP) P/T	1 RTT (TCP) 1 RTT (HTTP) M/T	1 RTT (TCP) 1 RTT (HTTP) M/T	1 RTT (TCP) 1 RTT (HTTP) M/T	8 RTTs + P/T + 3M/T

Q3

c	Page	Media 1	Media 2	Media 3	Total
Concurrent with non-persistent TCP connections	<p>1 RTT (TCP)</p> <p>1 RTT (HTTP)</p> <p>P/T</p>	<p>1 RTT (TCP)</p> <p>1 RTT (HTTP)</p> <p>M/(T/3)</p>	<p>1 RTT (TCP)</p> <p>1 RTT (HTTP)</p> <p>M/(T/3)</p>	<p>1 RTT (TCP)</p> <p>1 RTT (HTTP)</p> <p>M/(T/3)</p>	<p>4 RTTs</p> <p>+ P/T</p> <p>+ 3M/T</p>

c	Page	Media 1	Media 2	Media 3	Total
Sequential with a single persistent TCP connection	1 RTT (TCP)				5 RTTs + P/T + 3M/T
	1 RTT (HTTP) P/T	1 RTT (HTTP) M/T	1 RTT (HTTP) M/T	1 RTT (HTTP) M/T	
Pipelined within a single persistent TCP connection	1 RTT (TCP)				3 RTTs + P/T + 3M/T
	1 RTT (HTTP) P/T	1 RTT (HTTP)			
		M/T	M/T	M/T	