# Simulated annealing

- Resembles the annealing process used to cool metals slowly to reach an ordered (low-energy) state
- Basic idea:
  - Allow "bad" moves occasionally, depending on "temperature"
  - High temperature => more bad moves allowed, shake the system out of its local minimum
  - Gradually reduce temperature according to some schedule
  - Sounds pretty flaky, doesn't it?

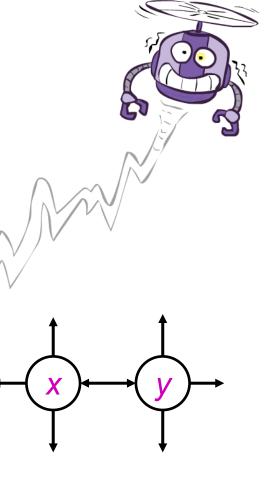
# Simulated annealing algorithm

- function SIMULATED-ANNEALING(problem, schedule) returns a state
- current ← problem.initial-state
- for t = 1 to  $\infty$  do
  - $T \leftarrow schedule(t)$
  - if T = 0 then return current
  - $\mathsf{next} \leftarrow \mathsf{a} \text{ randomly selected successor of } \mathsf{current}$
  - $\Delta E \leftarrow next.value current.value$
  - **if**  $\Delta E > 0$  **then** current  $\leftarrow$  next
    - else current  $\leftarrow$  next only with probability  $e^{\Delta E/T}$

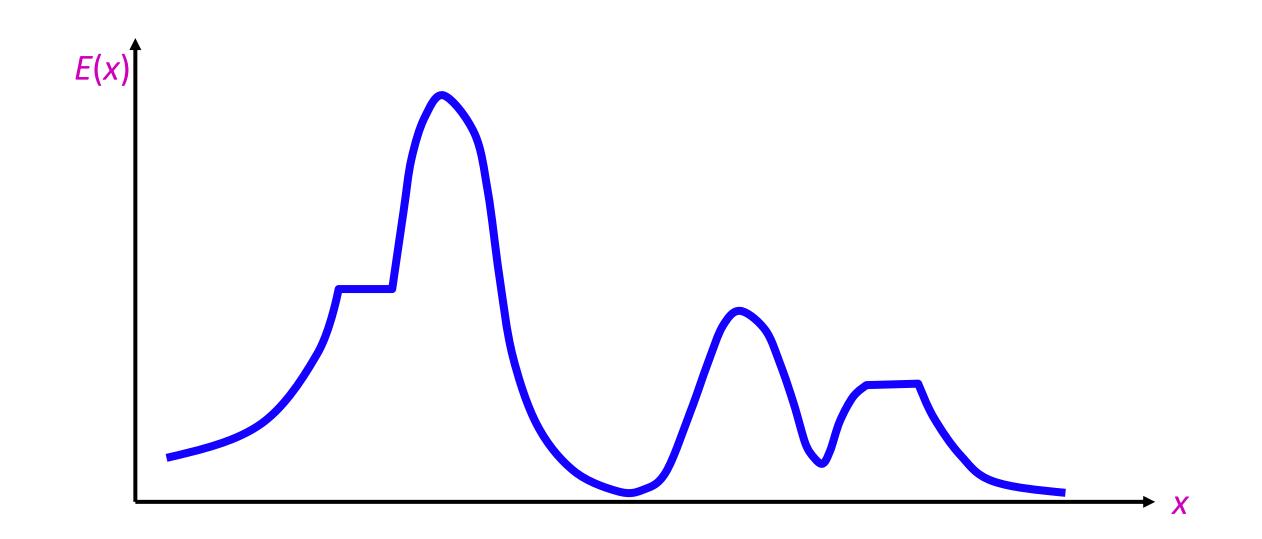


# **Simulated Annealing**

- Theoretical guarantee:
  - Stationary distribution (Boltzmann):  $P(x) \propto e^{E(x)/T}$
  - If T decreased slowly enough, will converge to optimal state!
- Proof sketch
  - Consider two adjacent states x, y with E(y) > E(x) [high is good]
  - Assume  $x \rightarrow y$  and  $y \rightarrow x$  and outdegrees D(x) = D(y) = D
  - Let P(x), P(y) be the equilibrium occupancy probabilities at T
  - Let  $P(x \rightarrow y)$  be the probability that state x transitions to state y



## Occupation probability as a function of *T*



# **Simulated Annealing**

- Is this convergence an interesting guarantee?
- Sounds like magic, but reality is reality:
  - The more downhill steps you need to escape a local optimum, the less likely you are to ever make them all in a row
  - "Slowly enough" may mean exponentially slowly
  - Random restart hillclimbing also converges to optimal state...
- Simulated annealing and its relatives are a key workhorse in VLSI layout and other optimal configuration problems

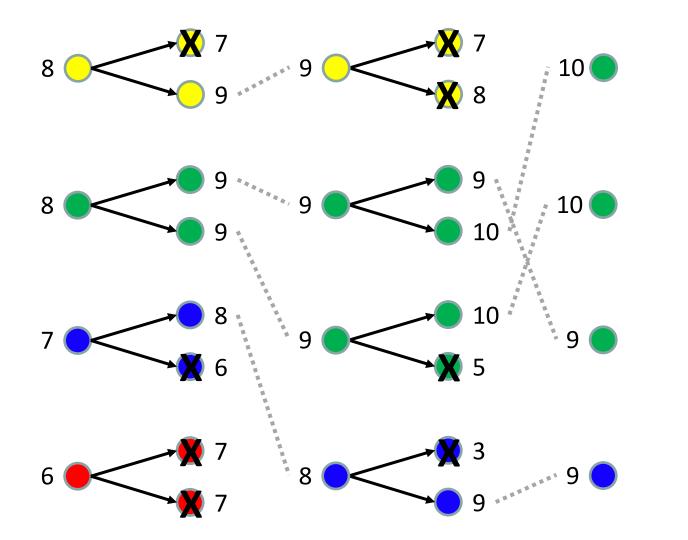


## Local beam search

Or, K chosen randomly with

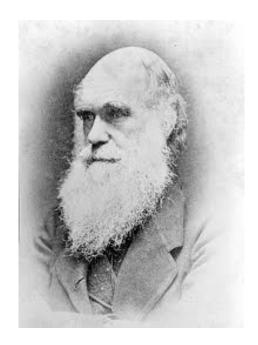
- Basic idea:
  - K copies of a local search algorithm, initialized randomly
  - For each iteration
    - a bias towards good ones
       Generate ALL successors from K current states
    - Choose best K of these to be the new current states

### Beam search example (K=4)

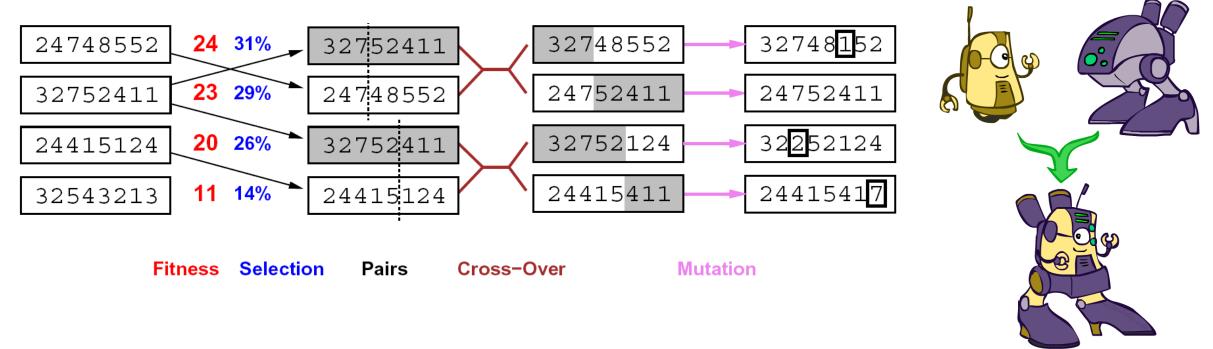


## Local beam search

- Why is this different from *K* local searches in parallel?
  - The searches communicate! "Come over here, the grass is greener!"
- What other well-known algorithm does this remind you of?
  - Evolution!

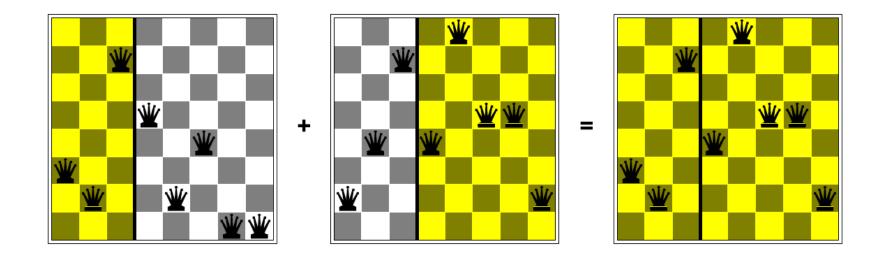


## Genetic algorithms



- Genetic algorithms use a natural selection metaphor
  - Resample K individuals at each step (selection) weighted by fitness function
  - Combine by pairwise crossover operators, plus mutation to give variety

### **Example: N-Queens**



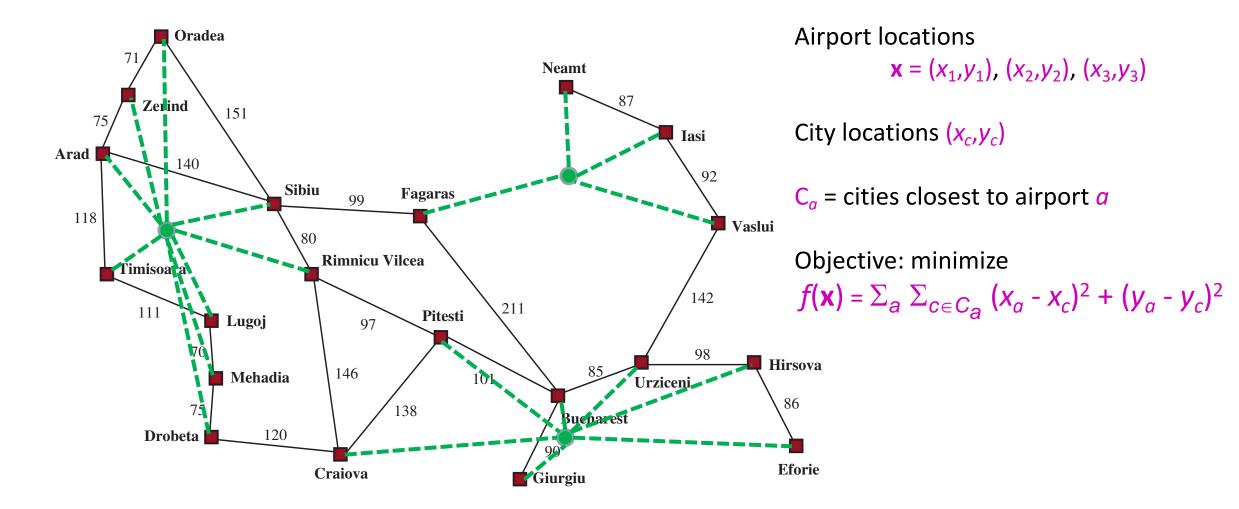
- Does crossover make sense here?
- What would mutation be?
- What would a good fitness function be?

### Local search in continuous spaces



## Example: Siting airports in Romania

Place 3 airports to minimize the sum of squared distances from each city to its nearest airport



# Handling a continuous state/action space

#### 1. Discretize it!

- Define a grid with increment  $\delta$ , use any of the discrete algorithms
- 2. Choose random perturbations to the state
  - a. First-choice hill-climbing: keep trying until something improves the state
  - b. Simulated annealing
- 3. Compute gradient of *f*(**x**) analytically

### Finding extrema in continuous space

- Gradient vector  $\nabla f(\mathbf{x}) = (\partial f / \partial x_1, \partial f / \partial y_1, \partial f / \partial x_2, ...)^{\mathsf{T}}$
- For the airports,  $f(\mathbf{x}) = \sum_a \sum_{c \in C_a} (x_a x_c)^2 + (y_a y_c)^2$
- $\partial f/\partial x_1 = \sum_{c \in C_1} 2(x_1 x_c)$
- At an extremum,  $\nabla f(\mathbf{x}) = 0$
- Can sometimes solve in closed form:  $x_1 = (\sum_{c \in C_1} x_c) / |C_1|$
- Is this a local or global minimum of *f*?
- Gradient descent:  $\mathbf{x} \leftarrow \mathbf{x} \alpha \nabla f(\mathbf{x})$ 
  - Huge range of algorithms for finding extrema using gradients

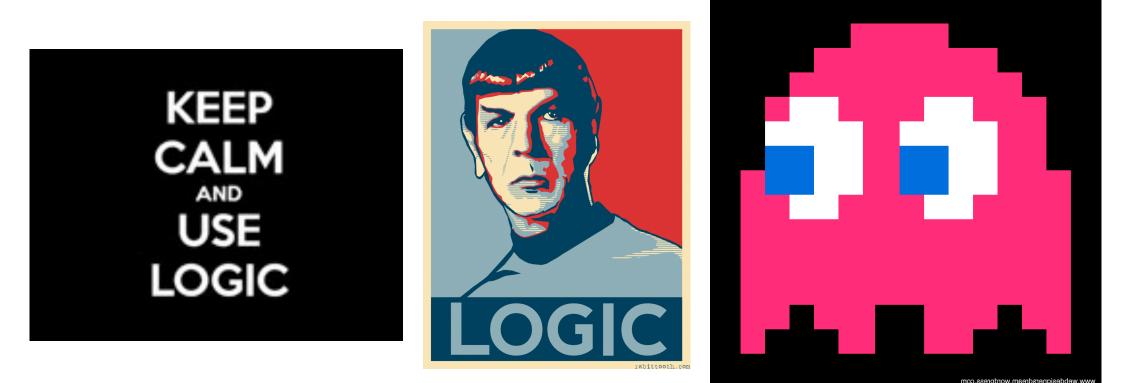
# Summary

- Many configuration and optimization problems can be formulated as local search
- General families of algorithms:
  - Hill-climbing, continuous optimization
  - Simulated annealing (and other stochastic methods)
  - Local beam search: multiple interaction searches
  - Genetic algorithms: break and recombine states

Many machine learning algorithms are local searches

# CS 188: Artificial Intelligence

#### **Propositional Logic I**



Instructors: Stuart Russell and Peyrin Kao

University of California, Berkeley

# Outline

#### 1. Propositional Logic I

- Basic concepts of knowledge, logic, reasoning
- Propositional logic: syntax and semantics, Pacworld example
- 2. Propositional logic II
  - Inference by theorem proving (briefly) and model checking
  - A Pac agent using propositional logic

# Agents that know things

- Agents acquire knowledge through perception, learning, language
  - Knowledge of the effects of actions ("transition model")
  - Knowledge of how the world affects sensors ("sensor model")
  - Knowledge of the current state of the world
- Can keep track of a partially observable world
- Can formulate plans to achieve goals
- Can design and build gravitational wave detectors.....

### LIGO



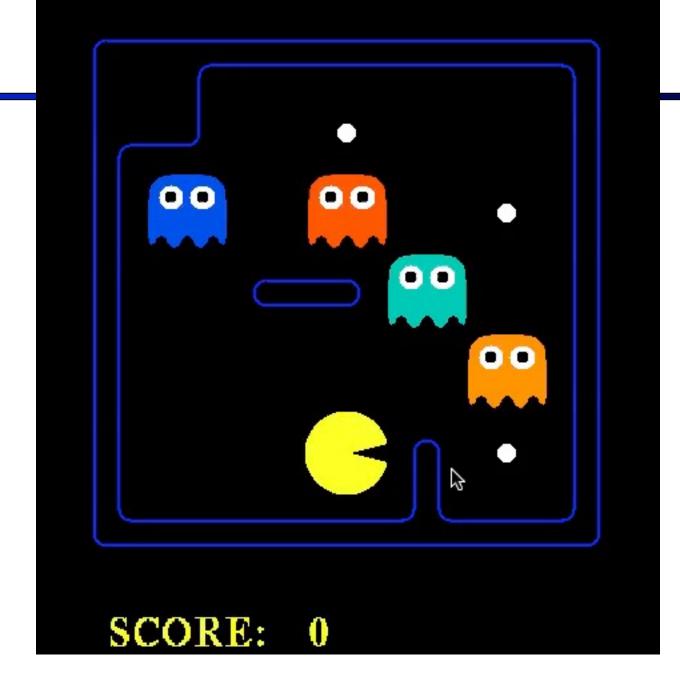
# Knowledge, contd.

- Knowledge base = set of sentences in a formal language
- Declarative approach to building an agent (or other system):
  - Tell it what it needs to know (or have it Learn the knowledge)
  - Then it can Ask itself what to do—answers should follow from the KB
- Agents can be viewed at the *knowledge level* i.e., what they *know*, regardless of how implemented
- A single inference algorithm can answer any answerable question

Knowledge base Inference engine

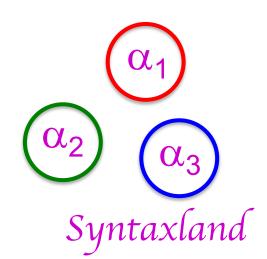
**Domain-specific facts** 

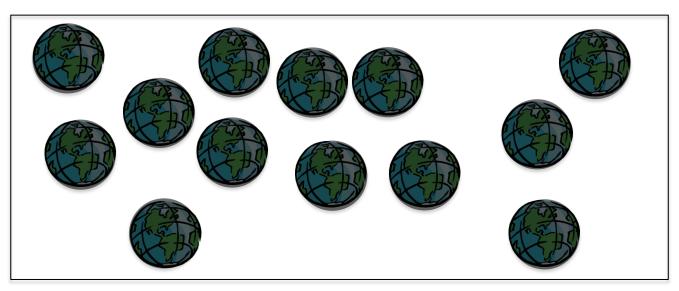
Generic code



# Logic

- *Syntax*: What sentences are allowed?
- Semantics:
  - What are the **possible worlds**?
  - Which sentences are true in which worlds? (i.e., definition of truth)





Semantícsland

# Different kinds of logic

#### Propositional logic

- Syntax:  $P \lor (\neg Q \land R)$ ;  $X_1 \Leftrightarrow$  (Raining  $\Rightarrow \neg$ Sunny)
- Possible world: {P=true,Q=true,R=false,S=true} or 1101
- Semantics:  $\alpha \land \beta$  is true in a world iff is  $\alpha$  true and  $\beta$  is true (etc.)
- First-order logic
  - Syntax:  $\forall x \exists y P(x,y) \land \neg Q(Joe,f(x)) \Rightarrow f(x)=f(y)$
  - Possible world: Objects o<sub>1</sub>, o<sub>2</sub>, o<sub>3</sub>; P holds for <o<sub>1</sub>, o<sub>2</sub>>; Q holds for <o<sub>3</sub>>; f(o<sub>1</sub>)=o<sub>1</sub>; Joe=o<sub>3</sub>; etc.
  - Semantics:  $\phi(\sigma)$  is true in a world if  $\sigma = o_j$  and  $\phi$  holds for  $o_j$ ; etc.

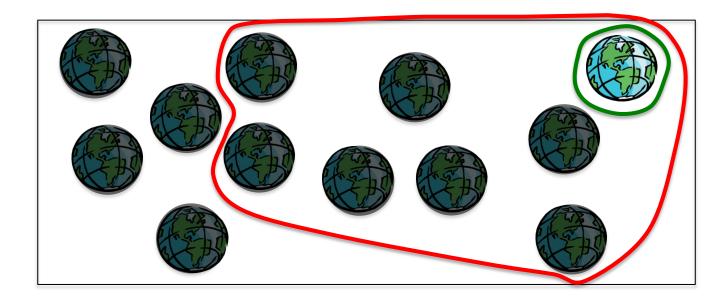
## Different kinds of logic, contd.

#### Relational databases:

- Syntax: ground relational sentences, e.g., Sibling(Ali,Bo)
- Possible worlds: (typed) objects and (typed) relations
- Semantics: sentences in the DB are true, everything else is false
  - Cannot express disjunction, implication, universals, etc.
  - Query language (SQL etc.) typically some variant of first-order logic
  - Often augmented by first-order rule languages, e.g., Datalog
- Knowledge graphs (roughly: relational DB + ontology of types and relations)
  - Google Knowledge Graph: 5 billion entities, 500 billion facts, >30% of queries
  - Facebook network: 2.93 billion people, trillions of posts, maybe quadrillions of facts

# Inference: entailment

- **Entailment**:  $\alpha \models \beta$  (" $\alpha$  entails  $\beta$ " or " $\beta$  follows from  $\alpha$ ") iff in every world where  $\alpha$  is true,  $\beta$  is also true
  - I.e., the  $\alpha$ -worlds are a <u>subset</u> of the  $\beta$ -worlds [models( $\alpha$ )  $\subseteq$  models( $\beta$ )]
- In the example,  $\alpha_2 \models \alpha_1$
- (Say  $\alpha_2$  is  $\neg Q \land R \land S \land W$  $\alpha_1$  is  $\neg Q$ )  $\alpha_1$



# Inference: proofs

- A proof is a *demonstration* of entailment between  $\alpha$  and  $\beta$
- **Sound** algorithm: everything it claims to prove is in fact entailed
- Complete algorithm: every that is entailed can be proved

# Inference: proofs

#### Method 1: model-checking

- For every possible world, if  $\alpha$  is true make sure that is  $\beta$  true too
- OK for propositional logic (finitely many worlds); not easy for first-order logic

#### Method 2: theorem-proving

- Search for a sequence of proof steps (applications of *inference rules*) leading from  $\alpha$  to  $\beta$
- E.g., from P and (P ⇒ Q), infer Q by Modus Ponens

## Propositional logic syntax

- Given: a set of proposition symbols {X<sub>1</sub>, X<sub>2</sub>,..., X<sub>n</sub>}
  - (we often add True and False for convenience)
- X<sub>i</sub> is a sentence
- If  $\alpha$  is a sentence then  $\neg \alpha$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \wedge \beta$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \lor \beta$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \Rightarrow \beta$  is a sentence
- If  $\alpha$  and  $\beta$  are sentences then  $\alpha \Leftrightarrow \beta$  is a sentence
- And p.s. there are no other sentences!

### **Propositional logic semantics**

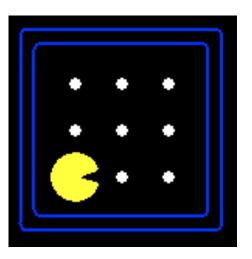
- Let m be a model assigning true or false to {X<sub>1</sub>, X<sub>2</sub>,..., X<sub>n</sub>}
- If  $\alpha$  is a symbol then its truth value is given in *m*
- $\neg \alpha$  is true in *m* iff  $\alpha$  is false in *m*
- $\alpha \wedge \beta$  is true in *m* iff  $\alpha$  is true in *m* and  $\beta$  is true in *m*
- $\alpha \lor \beta$  is true in *m* iff  $\alpha$  is true in *m* or  $\beta$  is true in *m*
- $\alpha \Rightarrow \beta$  is true in *m* iff  $\alpha$  is false in *m* or  $\beta$  is true in *m*
- $\alpha \Leftrightarrow \beta$  is true in *m* iff  $\alpha \Rightarrow \beta$  is true in *m* and  $\beta \Rightarrow \alpha$  is true in *m*

# Example: Partially observable Pacman

- Pacman knows the map but perceives just wall/gap to NSEW
- Formulation: what variables do we need?
  - Wall locations
    - Wall\_0,0 there is a wall at [0,0]
    - Wall\_0,1 there is a wall at [0,1], etc. (N symbols for N locations)
  - Percepts

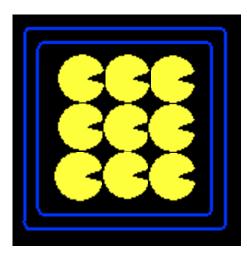
Diocked\_W (blocked by wall to my West) etc.

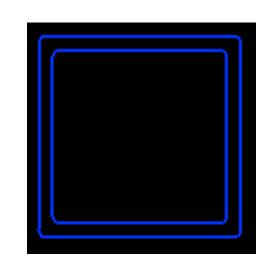
- Blocked\_W\_0 (blocked by wall to my West <u>at time 0</u>) etc. (47 symbols for T time steps)
- Actions
  - W\_0 (Pacman moves West at time 0), E\_0 etc. (4T symbols)
- Pacman's location
  - At\_0,0\_0 (Pacman is at [0,0] at time 0), At\_0,1\_0 etc. (NT symbols)

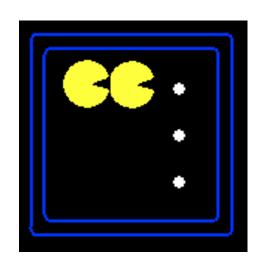


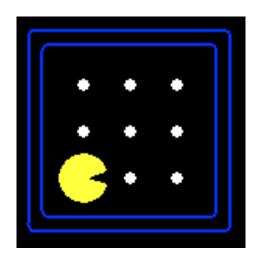
## How many possible worlds?

- N locations, T time steps => N + 4T + 4T + NT = O(NT) variables
- *O*(2<sup>*NT*</sup>) possible worlds!
- N=200, T=400 => ~10<sup>24000</sup> worlds
- Each world is a complete "history"
  - But most of them are pretty weird!



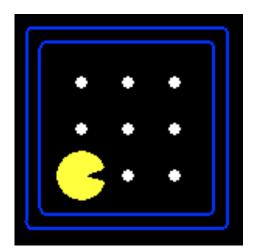






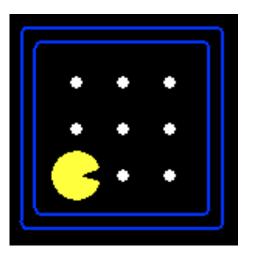
### Pacman's knowledge base: Map

- Pacman knows where the walls are:
  - Wall\_0,0 ^ Wall\_0,1 ^ Wall\_0,2 ^ Wall\_0,3 ^ Wall\_0,4 ^ Wall\_1,4 ^ ...
- Pacman knows where the walls aren't!
  - $\neg$ Wall\_1,1  $\land \neg$ Wall\_1,2  $\land \neg$ Wall\_1,3  $\land \neg$ Wall\_2,1  $\land \neg$ Wall\_2,2  $\land \dots$



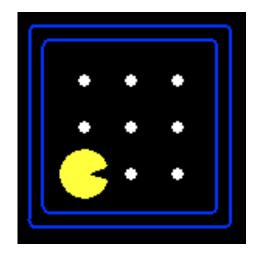
### Pacman's knowledge base: Initial state

- Pacman doesn't know where he is
- But he knows he's somewhere!
  - At\_1,1\_0  $\lor$  At\_1,2\_0  $\lor$  At\_1,3\_0  $\lor$  At\_2,1\_0  $\lor$  ...



# Pacman's knowledge base: Sensor model

- State facts about how Pacman's percepts arise...
  - Percept variable at t> \leftarrow <some condition on world at t>
- Pacman perceives a wall to the West at time t if and only if he is in x, y and there is a wall at x-1, y
  - Blocked\_W\_0  $\Leftrightarrow$  ((At\_1,1\_0  $\land$  Wall\_0,1) v (At\_1,2\_0  $\land$  Wall\_0,2) v
    - (At\_1,3\_0 ^ Wall\_0,3) v .... )
  - 4T sentences, each of size O(N)
  - Note: these are valid for any map



### Pacman's knowledge base: Transition model

- How does each state variable at each time gets its value?
  - Here we care about location variables, e.g., At\_3,3\_17
- A state variable X gets its value according to a successor-state axiom
  - X\_t ⇔ [X\_t-1 ∧ ¬(some action\_t-1 made it false)] v
    [¬X t-1 ∧ (some action t-1 made it true)]
- For Pacman location:
  - At\_3,3\_17 ⇔ [At\_3,3\_16 ∧ ¬((¬Wall\_3,4 ∧ N\_16) ∨ (¬Wall\_4,3 ∧ E\_16) ∨ ...)]
     v [¬At\_3,3\_16 ∧ ((At\_3,2\_16 ∧ ¬Wall\_3,3 ∧ N\_16) ∨ ...)]
     (At\_2,3\_16 ∧ ¬Wall\_3,3 ∧ N\_16) ∨ ...)]

### How many sentences?

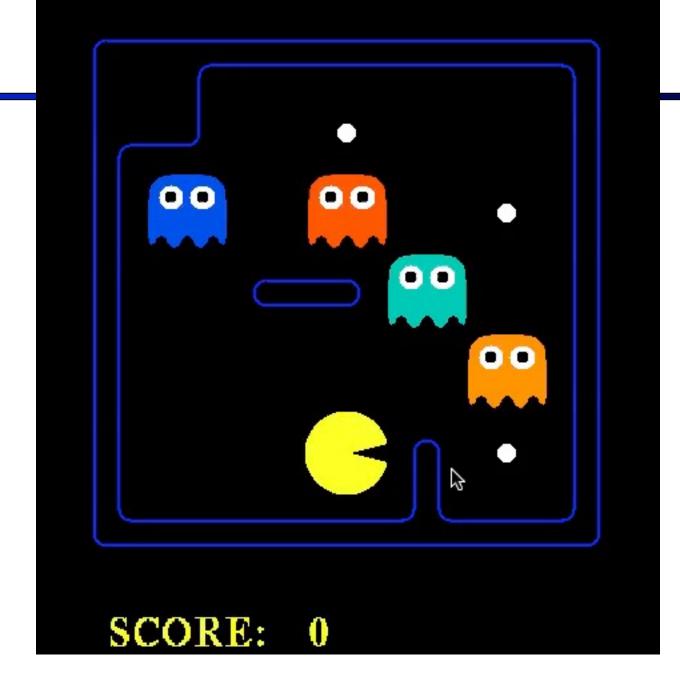
- Vast majority of KB occupied by O(NT) transition model sentences
  - Each about 10 lines of text
  - N=200, T=400 => ~800,000 lines of text, or 20,000 pages
- This is because propositional logic has limited expressive power
- Are we really going to write 20,000 pages of logic sentences???
- No, but your code will generate all those sentences!
- In first-order logic, we need O(1) transition model sentences
- (State-space search uses atomic states: how do we keep the transition model representation small???)

## Some reasoning tasks

#### Localization with a map and local sensing:

- Given an initial KB, plus a sequence of percepts and actions, where am I?
- Mapping with a location sensor:
  - Given an initial KB, plus a sequence of percepts and actions, what is the map?
- Simultaneous localization and mapping:
  - Given ..., where am I and what is the map?
- Planning:
  - Given ..., what action sequence is guaranteed to reach the goal?

#### ALL OF THESE USE THE SAME KB AND THE SAME ALGORITHM!!



# Summary

- One possible agent architecture: knowledge + inference
- Logics provide a formal way to encode knowledge
  - A logic is defined by: syntax, set of possible worlds, truth condition
- A simple KB for Pacman covers the initial state, sensor model, and transition model
- Logical inference computes entailment relations among sentences, enabling a wide range of tasks to be solved