## Q1. Power Pellets

Consider a Pacman game where Pacman can eat 3 types of pellets:

- Normal pellets (n-pellets), which are worth one point.
- Decaying pellets (d-pellets), which are worth $\max (0,5-t)$ points, where $t$ is time.
- Growing pellets (g-pellets), which are worth $t$ points, where $t$ is time.

The location and type of each pellet is fixed. The pellet's point value stops changing once eaten. For example, if Pacman eats one g-pellet at $t=1$ and one d-pellet at $t=2$, Pacman will have won $1+3=4$ points.

Pacman needs to find a path to win at least 10 points but he wants to minimize distance travelled. The cost between states is equal to distance travelled.
(a) Which of the following must be including for a minimum, sufficient state space?

## Pacman's location

Location and type of each pellet
How far Pacman has travelled
Current time
$\square$ How many pellets Pacman has eaten and the point value of each eaten pellet
Total points Pacman has won
Which pellets Pacman has eaten
(b) Which of the following are admissible heuristics? Let $x$ be the number of points won so far.Distance to closest pellet, except if in the goal state, in which case the heuristic value is 0 .
$\square$
Distance needed to win $10-x$ points, determining the value of all pellets as if they were n-pellets.
Distance needed to win $10-x$ points, determining the value of all pellets as if they were $g$-pellets
(i.e. all pellet values will be $t$.)
$\square$ Distance needed to win $10-x$ points, determining the value of all pellets as if they were d-pellets (i.e. all pellet values will be $\max (0,5-t)$.

Distance needed to win $10-x$ points assuming all pellets maintain current point value (g-pellets stop increasing in value and d-pellets stop decreasing in value)
$\square$ None of the above
(c) Instead of finding a path which minimizes distance, Pacman would like to find a path which minimizes the following:

$$
C_{\text {new }}=a * t+b * d
$$

where $t$ is the amount of time elapsed, $d$ is the distance travelled, and $a$ and $b$ are non-negative constants such that $a+b=1$. Pacman knows an admissible heuristic when he is trying to minimize time (i.e. when $a=1, b=0$ ), $h_{t}$, and when he is trying to minimize distance, $h_{d}$ (i.e. when $a=0, b=1$ ).
Which of the following heuristics is guaranteed to be admissible when minimizing $C_{\text {new }}$ ?$\operatorname{mean}\left(h_{t}, h_{d}\right)$$\min \left(h_{t}, h_{d}\right)$$\max \left(h_{t}, h_{d}\right)$$a * h_{t}+b * h_{d}$
None of the above

## Q2. Rubik's Search

A Rubik's cube has about $4.3 \times 10^{19}$ possible configurations, but any configuration can be solved in 20 moves or less. We pose the problem of solving a Rubik's cube as a search problem, where the states are the possible configurations, and there is an edge between two states if we can get from one state to another in a single move. Thus, we have $4.3 \times 10^{19}$ states. Each edge has cost 1 . Since we can make 27 moves from each state, the branching factor is 27 . Since any configuration can be solved in 20 moves or less, we have $h^{*}(n) \leq 20$.

For each of the following searches, estimate the approximate number of states expanded. Mark the option that is closest to the number of states expanded by the search. Assume that the shortest solution for our start state takes exactly 20 moves. Note that $27^{20}$ is much larger than $4.3 \times 10^{19}$.
(a) DFS Tree Search
(i) Best Case:20
$\bigcirc 4.3 \times 10^{19}$$27^{20}$
$\bigcirc \infty$ (never finishes)
(ii) Worst Case:
$\bigcirc 20$
$\bigcirc 4.3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc$ (never finishes)
(b) DFS graph search
(i) Best Case:
$\bigcirc 20$
$\bigcirc .3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(ii) Worst Case:
$\bigcirc \quad 20$
$\bigcirc 4.3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(c) BFS tree search
(i) Best Case:20
〇 $4.3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(ii) Worst Case:
$\bigcirc 20$
$\bigcirc 4.3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(d) BFS graph search
(i) Best Case:20$4.3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(ii) Worst Case:
$\bigcirc 20$
$\bigcirc 4.3 \times 10^{19}$$27^{20}$
$\bigcirc$ (never finishes)
(e) A* tree search with a perfect heuristic, $h^{*}(n)$, Best Case
$\bigcirc 20$
○ $4.3 \times 10^{19}$
$\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(f) $A^{*}$ tree search with a bad heuristic, $h(n)=20-h^{*}(n)$, Worst Case ○ 20 O $4.3 \times 10^{19}$ $\bigcirc 27^{20}$
$\bigcirc \infty$ (never finishes)
(g) A* graph search with a perfect heuristic, $h^{*}(n)$, Best Case
$\bigcirc \quad 20$
O $4.3 \times 10^{19}$
$27^{20}$$\infty$ (never finishes)
(h) A* graph search with a bad heuristic, $h(n)=20-h^{*}(n)$, Worst Case $\bigcirc 20$

O $4.3 \times 10^{19}$
$27^{20}$
$\bigcirc$ (never finishes)

## Q3. CSPs: Potluck Pandemonium

The potluck is coming up and the staff haven't figured out what to bring yet! They've pooled their resources and determined that they can bring some subset of the following items.

1. Pho
2. Apricots
3. Frozen Yogurt
4. Fried Rice
5. Apple Pie
6. Animal Crackers

There are five people on the course staff: Taylor, Jonathan, Faraz, Brian, and Alvin. Each of them will only bring one item to the potluck.
i. If (F)araz brings the same item as someone else, it cannot be (B)rian.
ii. (A)lvin has pho-phobia so he won't bring Pho, but he'll be okay if someone else brings it.
iii. (B)rian is no longer allowed near a stove, so he can only bring items 2,3 , or 6 .
iv. (F)araz literally can't even; he won't bring items 2,4 , or 6 .
v. (J)onathan was busy, so he didn't see the last third of the list. Therefore, he will only bring item $1,2,3$, or 4.
vi. (T)aylor will only bring an item that is before an item that (J)onathan brings.
vii. (T)aylor is allergic to animal crackers, so he won't bring item 6. (If someone else brings it, he'll just stay away from that table.)
viii. (F)araz and (J)onathan will only bring items that have the same first letter (e.g. Frozen Yogurt and Fried Rice).
ix. (B)rian will only bring an item that is after an item that (A)lvin brings on the list.
x. (J)onathan and (T)aylor want to be unique; they won't bring the same item as anyone else.
(a) Which of the listed constraints are unary constraints?

(b) Rewrite implicit constraint viii. as an explicit constraint.
(c) How many edges are there in the constraint graph for this CSP?
(d) The table below shows the variable domains after all unary constraints have been enforced.

| $\mathbf{A}$ |  | 2 | 3 | 4 | 5 | 6 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{B}$ |  | 2 | 3 |  |  | 6 |
| $\mathbf{F}$ | 1 |  | 3 |  | 5 |  |
| $\mathbf{J}$ | 1 | 2 | 3 | 4 |  |  |
| $\mathbf{T}$ | 1 | 2 | 3 | 4 | 5 |  |

Following the MRV heuristic, which variable should we assign first? Break all ties alphabetically.A
$\bigcirc B$
F
○ J
T
(e) To decouple this from the previous question, assume that we choose to assign (F)araz first. In this question, we will choose which value to assign to using the Least Constraining Value method.
To determine the number of remaining values, enforce arc consistency to prune the domains. Then, count the total number of possible assignments (not the total number of remaining values). It may help you to enforce arc consistency twice, once before assigning values to (F)araz, and then again after assigning a value.
(i) Assigning $\mathrm{F}=$ $\qquad$ results in $\qquad$ possible assignments.
(ii) Assigning $\mathrm{F}=$ $\qquad$ results in $\qquad$ possible assignments.
(iii) Assigning $\mathrm{F}=$ $\qquad$ results in $\qquad$ possible assignments.
(iv) Using the LCV method, which value should we assign to F? If there is a tie, choose the lower number.
$\square$ 1 2

O 3
$3 \bigcirc 4$5
6
Extra tables for work:

| $\mathbf{A}$ |  | 2 | 3 | 4 | 5 | 6 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{B}$ |  | 2 | 3 |  |  | 6 |
| $\mathbf{F}$ | 1 |  | 3 |  | 5 |  |
| $\mathbf{J}$ | 1 | 2 | 3 | 4 |  |  |
| $\mathbf{T}$ | 1 | 2 | 3 | 4 | 5 |  |


| $\mathbf{A}$ |  | 2 | 3 | 4 | 5 | 6 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{B}$ |  | 2 | 3 |  |  | 6 |
| $\mathbf{F}$ | 1 |  | 3 |  | 5 |  |
| $\mathbf{J}$ | 1 | 2 | 3 | 4 |  |  |
| $\mathbf{T}$ | 1 | 2 | 3 | 4 | 5 |  |


| $\mathbf{A}$ |  | 2 | 3 | 4 | 5 | 6 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{B}$ |  | 2 | 3 |  |  | 6 |
| $\mathbf{F}$ | 1 |  | 3 |  | 5 |  |
| $\mathbf{J}$ | 1 | 2 | 3 | 4 |  |  |
| $\mathbf{T}$ | 1 | 2 | 3 | 4 | 5 |  |


| $\mathbf{A}$ |  | 2 | 3 | 4 | 5 | 6 |
| :---: | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{B}$ |  | 2 | 3 |  |  | 6 |
| $\mathbf{F}$ | 1 |  | 3 |  | 5 |  |
| $\mathbf{J}$ | 1 | 2 | 3 | 4 |  |  |
| $\mathbf{T}$ | 1 | 2 | 3 | 4 | 5 |  |

